# **BAZARLOADER: Unpacking an ISO File Infection**

() Offset.net/reverse-engineering/bazarloader-iso-file-infection/

April 19, 2022

```
qmemcpy(stack_str, v16, 0xDui64);
v2 = 0i64;
stack_str[0xD] = 0;
do
 stack_str[v2] = ((0xFFFFFFC2 * (stack_str[v2] - 0x33)) % 0x7F + 0x7F) % 0x7F;
  ++v2;
}
while ( v2 \neq 0xD );
kernel32_dll_handle = ::GetModuleHandleA(stack_str);// "Kernel32.dll"
output-kernel32_handle = kernel32_dll_handle;
if ( kernel32_dll_handle )
£
 *&v19[8] = 0x4C0C7878;
 *v19 = 0x925110C6F334C7Ei64;
 qmemcpy(&v19[0xC], "__K", 3);
 qmemcpy(stack_str, v19, 0xFui64);
 v4 = 0i64;
 stack_str[0xF] = 0;
 do
 Ł
   stack_str[v4] = ((0xFFFFFE6 * (stack_str[v4] - 0x4B)) % 0x7F + 0x7F) % 0x7F;
   ++v4;
 }
 while ( v4 \neq 0xF );
 GetProcAddress = ::GetProcAddress(kernel32_dll_handle, stack_str);// "GetProcAddress"
 *&v18[8] = 0x42680E7A;
 v6 = 0i64;

    <u>Chuong Dong</u>
```

- 19th April 2022
- 1 Comment

BAZARLOADER (aka BAZARBACKDOOR) is a Windows-based loader that spreads through attachments in phishing emails. During an infection, the final loader payload typically downloads and executes a Cobalt Strike beacon to provide remote access for the threat actors, which, in a lot of cases, leads to ransomware being deployed to the victim's machine.

In this initial post, we will unpack the different stages of a BAZARLOADER infection that comes in the form of an optical disk image (ISO) file. We will also dive into the obfuscation methods used by the main BAZARLOADER payload.

To follow along, you can grab the sample as well as the PCAP files for it on <u>Malware-Traffic-Analysis.net</u>.

SHA256: 0900b4eb02bdcaefd21df169d21794c8c70bfbc68b2f0612861fcabc82f28149

### Step 1: Mounting ISO File & Extracting Stage 1 Executable

Recent BAZARLOADER samples arrive in emails containing OneDrive links to download an ISO file to avoid detection since most AVs tend to ignore this particular file type. With Windows 7 and above integrating the mounting functionality into Windows Explorer, we can mount any ISO file as a virtual drive by double-clicking on it.

When we mount the malicious ISO file, we see that a drive is mounted on the system that contains a shortcut file named "**Attachments.Ink**" and a hidden file named "**documents.log**".



The shortcut file has to be run by the victim to begin the chain of infection. We can quickly extract the actual command being executed by this shortcut from its **Properties** window.

C:\Windows\System32\rundll32.exe documents.log,vspa

📕 Attachments	Properties	×
General Shorton	ut Details	
At	ttachments	
Target type:	Application	
Target location:	System32	
Target:	ws\System32\rundll32.exe documents.log,v	spa
Start in:		
Shortcut key:	None	
Run:	Normal window	$\sim$
Comment:		
Open File Lo	Change Icon Advanced.	
	OK Cancel App	ply

Once the victim double-clicks on the shortcut file, the command executes the Windows **rundll32.exe** program to launch the **"documents.log"** file. This lets us know that the file being launched is a DLL file, and the entry point is its export function **vspa**.

## Step 2: Extracting Second Stage Shellcode

Taking a quick look in IDA, we can somewhat tell that the extracted DLL is packed since it has only a few functions and a really suspicious looking buffer of ASCII characters in its custom **.odata** section.

odata:000000018001D000	odata	segment para public 'DATA' use64
odata:000000018001D000		assume cs:odata
odata:000000018001D000		;org 18001D000h
odata:000000018001D000	a46a52995f9a819	db '46a52995f9a8193ba0efa7b0724a8a0120d51ff08d25efd057bd2e82ddb7cf3b7'
odata:000000018001D000		; DATA XREF: sub_180003DDA+16B↑o
odata:000000018001D000		db \8b10a67cce05ae7ba5165f1056a55e16018b4ed9209b7abab438990ad1265733d
odata:000000018001D000		db '9f8c63da54c29ddca1b52ce9a76bb242593b05d64fbb5a3fb69597d3d585e816a'
odata:000000018001D000		db 'fb8b4732d7ce7a8e6295859b3500849a3c414a2513d8a47f5f526fc2134bcd790'
odata:000000018001D000		db '2ca04ec5f10f6ee393e94f784db16e720860b7fbdbc14627c340d97da2ec61f64'
odata:000000018001D000		db '0a76dda73907e734fc639336b9d877aea3cd4ea1e6f06da14899a0cc00be9348b'
odata:000000018001D000		db 'b8fcc24344aeae404958ac307254dde2d46c622acf094f3cf91d1749b1ec6650c'
odata:000000018001D000		db '34248e0cf8d5d9be621dce73c5485fae5284e1dc89d11a92bbada73f85452fe7e'
odata:000000018001D000		db '0f80b1949ce5b5a77bae11b971dc7cda9a4289c4f41a0e75ff7688f201ac670e5'
odata:000000018001D000		db '822582d189e6a404b6e1a4306fec1ceea8f4197751d09e9a3c7ae2a5a5777d785'
odata:000000018001D000		db '75736804b1eeb59541b14e0b2da06fbca87a8d7f182cc2bec9c2216fb5d008ec3'
odata:000000018001D000		db 'a709bdd81f443819ef59b2b76fb27b62f3a65e601d58a7ac034b06d004981cee0'
odata:000000018001D000		db '9fce980c429bfb99f57ce5aa327f3345e90025c8b4d86465e02128fb43ae3c993'
odata:000000018001D000		db '05ef4d217d4d59c9a55180957348d81dcd6d1f23228813751267c17bbdcff7cb9'
odata:000000018001D000		db '9094910fe27d00bfa451089831f6e3ba570e0d2f6acf622fccb0ab7852d59b326'
odata:000000018001D000		db 'c27b490b46adfad18756317257078fa0bcd7dc652976af11db8d80e390ff58cf4'
odata:000000018001D000		db '80767409e573309171b2e67ddc514344dc5ae51935f276acfd95b00bc98ce1376'
odata:000000018001D000		db '3bdc2580c907eaa727bedc1142133309a4e6ba1df780983d3611b1e228181587d'
odata:000000018001D000		db '1c3cb4972d35037fb8c864608baae3defd64b5be63d19bf3b969793721002cb35'
odata:000000018001D000		db '4bc685d2806e402edd6b2b0342be26b9e5867953537e136f61c79c81523c57c09'
odata:000000018001D000		db 'e7b293cb93446930f52ad4759b3db809084cf0e810bfe90398b25060023f838c0'
odata:000000018001D000		db
odata:000000018001D000		db '3523582c41233ebdb56159416efbd5df899b9cac591e93b8db3c6fc8b329ac81a'
odata:000000018001D000		db '627bd11bbc66e7ebc49d01abd0ed6e734f8995df02e3205823e99aa5efdc75025'
odata:000000018001D000		db '985ef694758f157b5d51eea17e973af119da097b54f9def2201e94bbb1523dff6'
odata:000000018001D000		db '03405df7916cf4fda8d6fc9059a2a7fb1470460cf286ac93e3928fe1d85125be2'

With that in mind, we will just perform some quick static analysis to determine where we can dump the next stage.

In the first function of the **vspa** export, we see **sub\_1800045D6** takes a DWORD in as the parameter. This function returns a variable that contains an address to a function that is later called in the code.

```
API_func = (__int64 (__fastcall *)(void *, __int64))sub_1800045D6(0x67CC0818i64);// API hashing?
v9 = a46a52995f9a819;
for ( i = 0; !i; i = 1 )
DWORD2(v6) += 147086;
if ( DWORD2(v6) )
{
 v7 = DWORD2(v6);
ProcessEnvironmentBlock = NtCurrentTeb() -> ProcessEnvironmentBlock;
*(_QWORD *)&v6 = API_func(ProcessEnvironmentBlock -> ProcessHeap, 8i64);
v5 = 0;
for ( j = 0; !j; j = 1 )
{
 v10 = v5 + (_QWORD)v6;
 sub_180004E6D(v10);
v5 += 147086;
}
```

At this point, we can safely guess that **sub\_1800045D6** is an API resolving function, and the parameter it takes is the hash of the API's name. Because this is still the unpacking phase, we won't dive too deep into analyzing this function.

Instead, I'll just use **OALabs's** <u>HashDB</u> IDA plugin to quickly reverse-lookup the hashing algorithm used from the hash. The result shows that the hash corresponds to an API name hashed with <u>Metasploit's hashing algorithm **ROR13**</u>.

<pre> •)(void </pre>	<pre>*. int64))sub 3</pre>	1800045D6(0x67CC0818i64):// API hashing?									
	Matched Algorithms		×								
	The following algorithms contain a matching hash. Select an algorithm to set as the default for HashDB.										
	Algorithm	Size (Bits)									
rrentTe essEnvi	metasploit	32									
	Line 1 of 1										
		OK Cancel									

After determining the hashing algorithm, we can use **HashDB** to quickly look up the APIs being resolved by this function. It becomes clear that this function resolves the

**RtIAllocateHeap** API, calls that to allocate a heap buffer and writes the encoded ASCII data to it.



From this point onward, we can guess that the packer will decode this buffer and launch it somewhere later in the code. If we skip toward the end of the **vspa** export, we see a **call** instructionon a variable that is not returned from the API resolving function, so it can potentially be our tail jump.

```
if ( (unsigned int)sub_180003FE6(&v19, (__int64)&v25, 0x40u) )
{
    return 0i64;
}
else
{
    qword_18001A9E0 = 0i64;
    v15 = 0i64;
    strcpy(v13, "vh5;");
    v13[1] += 11;
    v13[2] += 59;
    v13[3] += 38;
    v18 = v13;
    qword_18001A9E0 = v19(v26[4], qword_18001A9E8);// v19 is not returned from API_resolve?
```

The last function to modify that **v19** variable is **sub\_180003FE6**, so we can quickly take a look at that.

It turns out the **sub\_180003FE6** function just resolves and calls **NtMapViewOfSection** to map a view of a section into the virtual address space and writes the base address of the view into the **v19** variable. Then, it just executes **qmemcpy** to copy the data in the second variable to the returned virtual base address.

```
_int64 __fastcall w_map_view_of_section(__int64 BaseAddress, __int64 ViewSize, int Win32Protect)
ş
 // [COLLAPSED LOCAL DECLARATIONS. PRESS KEYPAD CTRL-"+" TO EXPAND]
 NtCreateSection = API_hashing(NtCreateSection_0);
 NtMapViewOfSection = API_hashing(NtMapViewOfSection_0);
 NtClose = API_hashing(NtClose_0);
 section_handle = 0i64;
 v8 = 0i64;
 ViewSize_1 = ViewSize;
 v4 = (NtCreateSection)(&section_handle, 0xF001Fi64);
 if (v4 \ge 0)
 {
   v4 = (NtMapViewOfSection)(
          section_handle,
          0xFFFFFFFFFFFFFFFfui64,
          BaseAddress,
          0i64,
          0i64,
          0i64.
          &v8,
          0,
          Win32Protect);
    (NtClose)(section_handle);
 }
 return v4;
```

```
__int64 __fastcall sub_180003FE6(_QWORD *a1, __int64 a2, int a3)
{
    int v4; // [rsp+0h] [rbp-18h]
    v4 = w_map_view_of_section(a1, *(a2 + 8), a3);
    if ( v4 ≥ 0 )
        w_qmemcpy(*a1, *a2, *(a2 + 8));
    return v4;
}
```

This tells us two things. First, our guess that the **v19** variable will contain the address to executable code is correct. Second, we know that the executable code is shellcode since the data is mapped and executed directly at offset 0 from where it is written.

From here, we can set up **x64dbg**, execute the DLL file at the **vspa** export, and break at the call instruction. After stepping into the function, we will be at the head of the shellcode.

💥 de	ocume	nts.log ·	PID: 800	8 - TI	hread: N	Aain Thr	ead 798	30 - x64d	lbg																		
File	View	Debug	Tracing	) P	lugins	Favouri	tes O	ptions	Help				ngine)														
_	-											-															
	9		→ "		ร ล	F   🦛	' ÷	1		2		=	1	-				📕   I		2							
2025																											
Sevel .	CPU	1	Log	•	Notes		вгеакро	oints		Memory	/ мар		Calls	stack	<u> </u>	SEH	5	Script		Symbo	DIS	Source	<ul> <li>References</li> </ul>		Inreads		Handles
RIP						000002	361128	0000	4	8:8BC	4			mov r	rax,rs	р											
						000002	36112E	0003	4	8:895	8 08			mov (	qword	ptr c	ls:[r	ax+8],	rbx								
						000002	361128	0007	- 1	C:894	8 20			mov (	gword	ptr c	is:[r	ax+20]	,19								
						000002	361126	OOOF	4	8:895	0 10			mov d	word	otr d	is: fr	ax+101	rdx			o. vapa					
							361128	0013		5					rbp												
							361128			6																	
						000002	361126	0015	5	7					rdi							di:&"46a529	95f9a8193ba0efa		4a8a01200	151110	8d25efd0
						000002	361128	0016	4	1:54				push	r12							12:EntryPol					
						000002	361126	0014		1.55				push	r14							14.0470					
						000002	361128	0010	4	1:57				push	r15												
							361128			8:88E				mov r	rbp,rs												
							361128			8:81E	C 8000	00000			rsp,80												
						000002	361126	0028	4	C:8BF	9				r15,rc							cx:"MZ"					
						000002	361128	002B	B	9 407	72607			mov	ecx,72	67740											
						000002	361120	0030	5	8 BFU	20278			call	23611	02574	10										
						000002	361126	0034	4	8:894	5 AR			mov e	word	ntr s	s:[r	bo-581	.rax								
						000002	361128	003E	4	8:8BD	8			mov r	rbx.ra	x			,. a.		l ra	ax:"vspa"					
							361128	0041	Ē	8 AE0	50000				23611	2E05F											
							361128	0046	В	9 58A	453E5				ecx,E5	53A45											
							361128	004B	4	C:8BE	8				r13,ra							ax:"vspa"					
						000002	361128	004E	E	8 A10	50000				23611	2E05F											
						000002	361128	0053	B	9 10E	18AC 3			mov e	ecx,C	8AE11											
						000002	361126	0058	2	C:88F	00000			mov r	22611	X 2EOEE						14:"MZ", ra					
						000002	361126	0060	8	9 AER	15094			mov	72011	SCR14											
						000002	361128	0065	4	8:894	5 88			mov e	word	otr s	s: fr	bo-481	.rax								
						000002	361128	0069	4	8:88F	0			mov r	rsi.ra	x					r:	ax:"vspa"					
						000002	361128	006C		8 830	50000				23611												
							361128		в	9 330	09E95				ecx,95	9E003											
							361128			8:894	5 CO				qword	ptr s	;s:[r	bp-40]									
						000002	361128	007A	4	8:8BF	8			mov r	rdi,ra	x						di:&"46a529	95f9a8193ba0efa	7b072	4a8a01200	<b>151ff</b> 0	8d25efd0
						000002	361128	007D	E	8 720	50000				23611	2E05F											
						000002	361126	0082	4	5:33E	4			XOF I	120,r	120					-	de l'accosti de	avt "vena"				
						000002	361126	0088		8+85D	8			test	rhy r	hy						<del>o: vsp</del> a", ri	axi vspa				
						000002	361126	0088		F84 4	305000	00		ie 23	361126	05D4											
						000002	361128	0091	4	D:85E	D			test	r13.r	13											
							361128	0094	¥ 0	F84 3	A05 000	00			361128	05D4											
							361128		- 4	D:85F	6			test													
						100003	261126	mnen	<b>••</b> ••	E 0 4 9	105000	101	1	14 71	261125	nen a					1						

We can now dump this virtual memory buffer to retrieve the second stage shellcode for the next unpacking step.

## Step 3: Extracting The Final BAZARLOADER Payload

When we examine the shellcode in IDA, we can quickly use the same trick with **HashDB** above to see that the shellcode also performs API hashing with Metasploit's **ROR13**.

At the entry point above, the shellcode resolves a set of functions that it will call, most notably **VirtualAlloc** and **VirtualProtect**. These two functions are typically used by packers to allocate virtual memory to decode and write the next stage executable in before launching it.

With this in mind, our next step should be debugging the shellcode and setting breakpoints at these two API calls. We can pick up where we are after dumping in **x64dbg** during **Step 2**, or we can launch the shellcode directly in our debugger using OALabs's <u>BlobRunner</u> or similar shellcode launcher.

Our first hit with **VirtualAlloc** is a call to allocate a virtual memory buffer at virtual address 0x204140000 with the size of 0x2A000 bytes.



We can run until **VirtualAlloc** returns and start monitoring the memory at address 0x204140000. After running until the next **VirtualProtect** call, we see that a valid PE executable has been written to this memory region.

qword ptr ds:[000	007FFB8ADE28	808 "ðLÑ^û∖×	k7F"]= <kernel< th=""><th>lbase.VirtualProtect&gt;</th></kernel<>	lbase.VirtualProtect>
.text:00007FFB8A	07BC70 kerne	el32.dll:\$18	3C70 #1B070 <	<virtualprotect></virtualprotect>
🕮 Dump 1 🛄	Dump 2	Dump 3	🛄 Dump 4	🕮 Dump 5 🚳 Watch 1 🖂 Locals 💙 Struct
Address	Hex			ASCII
000000204140000 000000204140020 000000204140020 000000204140030 000000204140050 000000204140050 000000204140050 000000204140070 000000204140050 0000000204140050 0000002041400D0 0000002041400D0 0000002041400D0 0000002041400F0 0000002041400F0 0000002041400F0 00000020414010	4D         5A         90         00           B8         00         00         00           00         00         00         00           00         00         00         00           00         00         00         00           00         00         00         00           01         1F         RA         0           02         73         20         7           74         20         62         6           60         67         64         6           50         45         00         0           00         00         14         0           00         00         20         0           00         00         20         0           00         00         00         00           00         00         00         00           00         00         00         00	$ \begin{array}{cccccccccccccccccccccccccccccccccccc$	00         04         00         00           00         40         00         00           00         00         00         00           00         00         00         00         00           00         00         00         00         00           00         00         00         00         00           00         12         18         01         -           72         61         6D         20         6           62         20         69         6E         2           00         14         00         00         00           00         13         00         0         00           00         00         14         03         0           00         00         10         00         00         00           00         00         10         00         00         00         00           00         00         00         00         00         00         00         00	00       FF       FF       00 <td< td=""></td<>
000000204140130 000000204140140 0000000204140150 0000000204140160 0000000204140170 0000000204140180	00 90 02 0 00 00 00 0 C0 F2 01 0 00 00 00 0 00 00 00 0 00 00 00 0	0       C4       00       00       00         0       00       00       00       00       00         0       28       00       00       00       00         0       00       00       00       00       00       00         0       00       00       00       00       00       00         0       00       00       00       00       00       00	00         00         00         00         00         00           00         00         00         00         00         00         00           00         00         00         00         00         00         00         00           00         14         62         02         0         00	00 00 00 00 00 00 00A 00 00 00 00 00 00 00 00 00 00 00 Ab 00 80 01 00 00 00 80 01 00 00 78 74 00 00 00 Text

Finally, we can dump this memory region into a file to extract the BAZARLOADER payload.

## Step 4: BAZARLOADER's String Obfuscation

As we begin performing static analysis on BAZARLOADER, it is crucial that we identify obfuscation methods that the malware uses.

One of those methods is string obfuscation, where the malware uses encoded stack strings to hide them from static analysis.



As shown, a typical encoded string is pushed on the stack and decoded dynamically using some multiplication, subtraction, and modulus operations.

There are different ways to resolve these stack strings, such as writing IDAPython scripts, <u>emulation</u>, or just running the program in a debugger and dumping the stack strings when they are resolved.



## Step 5: BAZARLOADER's API Obfuscation

BAZARLOADER obfuscates most of its API calls through a few structures that it constructs in the **DIIEntryPoint** function.

First, the malware populates the following structure that contains a handle to **Kernel32.dll** and addresses to API required to load libraries and get their API addresses.

```
struct API_IMPORT_STRUCT {
    HANDLE kernel32_handle;
    FARPROC mw_GetProcAddress;
    FARPROC mw_LoadLibraryW;
    FARPROC mw_LoadLibraryA;
    FARPROC mw_LoadLibraryA2;
    FARPROC mw_FreeLibrary;
    FARPROC mw_GetModuleHandleW;
    FARPROC mw_GetModuleHandleA;
};
```

It calls **GetModuleHandle** to retrieve the handle to **Kernel32.dll**, calls **GetProcAddress** to retrieve the address of the **GetProcAddress** API, and writes those in the structure.

```
qmemcpy(stack_str, v16, 0xDui64);
v2 = 0i64;
stack_str[0xD] = 0;
do
Ł
  stack_str[v2] = ((0xFFFFFFC2 * (stack_str[v2] - 0x33)) % 0x7F + 0x7F) % 0x7F;
  ++v2;
}
while ( v2 \neq 0xD );
kernel32_dll_handle = ::GetModuleHandleA(stack_str); // "Kernel32.dll"
output-kernel32_handle = kernel32_dll_handle;
if ( kernel32_dll_handle )
{
 *&v19[8] = 0x4C0C7878;
 *v19 = 0x925110C6F334C7Ei64;
 qmemcpy(&v19[0xC], "__K", 3);
 qmemcpy(stack_str, v19, 0xFui64);
 v4 = 0i64;
 stack_str[0xF] = 0;
 do
  £
   stack_str[v4] = ((0xFFFFFE6 * (stack_str[v4] - 0x4B)) % 0x7F + 0x7F) % 0x7F;
    ++v4;
 while ( v4 \neq 0xF );
 GetProcAddress = ::GetProcAddress(kernel32_dll_handle, stack_str); // "GetProcAddress"
 *&v18[8] = 0x42680E7A;
 v6 = 0i64;
```

Using the structure's **GetProcAddress** API field, BAZARLOADER retrieves the rest of the required APIs to populate other fields in the structure. This **API\_IMPORT\_STRUCT** structure will later be used to import other libraries' APIs.

```
GetModuleHandleW = get_proc_addr(output, output-kernel32_handle, stack_str); // "GetModuleHandleW"
v21 = 6;
v14 = 0i64;
output -> mw_GetModuleHandleW = GetModuleHandleW;
*&v20 = 0x5E3B6F28114D5D7Di64;
*(&v20 + 1) = 0x6A5D5E6F3A266B5Di64;
v26 = 0;
*stack_str = v20;
v25 = 6;
do
Ł
 stack_str[v14] = (7 * (stack_str[v14] - 6) % 0x7F + 0x7F) % 0x7F;
  ++v14;
}
while ( v14 \neq 0x11 );
GetModuleHandleA = get_proc_addr(output, output-kernel32_handle, stack_str);// "GetModuleHandleA"
```

Next, for each library to be imported, BAZARLOADER populates the following

**LIBRARY\_STRUCT** structure that contains a set of functions to interact with the library and the library handle.

```
struct LIB_FUNCS
{
    FARPROC free_lib;
    FARPROC w_free_lib;
    __int64 (__fastcall *get_API_addr)(API_IMPORT_STRUCT*, HANDLE, char*);
};
struct LIBRARY_STRUCT
{
    LIB_FUNCS *lib_funcs;
    HANDLE lib_handle;
};
```

The first 2 functions in the **LIB\_FUNCS** structure just call the **FreeLibrary** API from the global **API\_IMPORT\_STRUCT** to free the library module.

The third function calls the **GetProcAddress** from the **API\_IMPORT\_STRUCT**'s field to retrieve the address of an API exported from that specific library.

```
int __fastcall free_lib(LIBRARY_STRUCT *lib_struct)
{
    char *v1; // rax
    HANDLE lib_handle; // rcx
    v1 = &unk_20415F090 + 0x10;
    lib_struct→lib_funcs = (&unk_20415F090 + 0x10);
    lib_handle = lib_struct→lib_handle;
    if ( lib_handle )
      LODWORD(v1) = (API_IMPORT_STRUCT→mw_FreeLibrary)(lib_handle);
    return v1;
}
```

\_\_int64 \_\_fastcall get\_proc\_addr(API\_IMPORT\_STRUCT \*API\_IMPORT\_STRUCT, \_\_int64 library\_handle, \_\_int64 API\_to\_find)
{
 return (API\_IMPORT\_STRUCT->mw\_GetProcAddress)(library\_handle, API\_to\_find);

To begin populating each **LIBRARY\_STRUCT** structure, BAZARLOADER decodes the library name from a stack string and populates it with the corresponding set of functions and the library handle retrieved from calling **LoadLibraryA**.

```
int __fastcall set_up_lib_struct(LIBRARY_STRUCT *output, __int64 library_name)
ş
  bool v2; // zf
  API_IMPORT_STRUCT *v5; // rsi
  void *library_handle; // rax
  v2 = API_IMPORT_STRUCT == 0i64;
  output \rightarrow lib_funcs = (\&unk_20415F090 + 0x10);
  output \rightarrow lib_handle = 0i64;
  if ( v2 )
  Ł
    v5 = w_HeapAlloc(0x48ui64);
    populate_kernel32_funcs_maybe(v5);
    API_IMPORT_STRUCT = v5;
  }
  library_handle = (API_IMPORT_STRUCT -> mw_LoadLibraryA2)(library_name);
  output→lib_handle = library_handle;
  return test_exporting_library(library_handle);
```

Below is the list of all libraries used by the malware.

```
kernel32.dll, wininet.dll, advapi32.dll, ole32.dll, rpcrt4.dll, shell32.dll,
bcrypt.dll, crypt32.dll, dnsapi.dll, netapi32.dll, shlwapi.dll, user32.dll,
ktmw32.dll
```

The **LIBRARY\_STRUCT** structures corresponding to these are pushed into a global list in the order below.

```
struct LIBRARY_STRUCT_LIST
{
    LIBRARY_STRUCT *lib_struct_kernel32;
    LIBRARY_STRUCT *lib_struct_wininet;
    LIBRARY_STRUCT *lib_struct_advapi32;
    LIBRARY_STRUCT *lib_struct_ole32;
    LIBRARY_STRUCT *lib_struct_shell32;
    LIBRARY_STRUCT *lib_struct_shell32;
    LIBRARY_STRUCT *lib_struct_crypt32;
    LIBRARY_STRUCT *lib_struct_dnsapi;
    LIBRARY_STRUCT *lib_struct_netapi32;
    LIBRARY_STRUCT *lib_struct_netapi32;
    LIBRARY_STRUCT *lib_struct_netapi32;
    LIBRARY_STRUCT *lib_struct_netapi32;
    LIBRARY_STRUCT *lib_struct_netapi32;
    LIBRARY_STRUCT *lib_struct_shlwapi;
    LIBRARY_STRUCT *lib_struct_user32;
    LIBRARY_STRUCT *lib_struct_ktmw32;
```

```
v2 = w_HeapAlloc_16_bytes();
set_up_lib_struct_kernel32(v2);
LIBRARY_STRUCT_LIST \rightarrow lib_struct_kernel32 = v2;
v3 = w_HeapAlloc_16_bytes();
set_up_lib_struct_wininet(v3);
LIBRARY_STRUCT_LIST→lib_struct_wininet = v3;
v4 = w_HeapAlloc_16_bytes();
set_up_lib_struct_advapi32(v4);
LIBRARY_STRUCT_LIST→lib_struct_advapi32 = v4;
v5 = w_HeapAlloc_16_bytes();
set_up_lib_struct_ole32(v5);
v6 = w_HeapAlloc_16_bytes();
set_up_lib_struct_rpcrt4(v6);
LIBRARY_STRUCT_LIST \rightarrow lib_struct_rpcrt4 = v6;
v7 = w_HeapAlloc_16_bytes();
set_up_lib_struct_shell32(v7);
LIBRARY_STRUCT_LIST→lib_struct_shell32 = v7;
v8 = w_HeapAlloc_16_bytes();
set_up_lib_struct_bcrypt(v8);
v9 = w_HeapAlloc_16_bytes();
set_up_lib_struct_crypt32(v9);
LIBRARY_STRUCT_LIST→lib_struct_crypt32 = v9;
v10 = w_HeapAlloc_16_bytes();
set_up_lib_struct_dnsapi(v10);
LIBRARY_STRUCT_LIST \rightarrow lib_struct_dnsapi = v10;
v11 = w_HeapAlloc_16_bytes();
set_up_lib_struct_netapi32(v11);
LIBRARY_STRUCT_LIST→lib_struct_netapi32 = v11;
v12 = w_HeapAlloc_16_bytes();
set_up_lib_struct_shlwapi(v12);
```

After this global list of **LIBRARY\_STRUCT** is populated, an API can be called from a function taking in its corresponding library's **LIBRARY\_STRUCT** structure and its parameters.

This function resolves the API name from a stack string, retrieves the API's address using the **get\_API\_addr** function from the library structure, and calls the API with its parameters.

:	<pre>int64fastcall w_Sleep(LIBRARY_STRUCT *kernel32_lib_struct, unsigned int dwMilliseconds)</pre>
{	
	int64 i; // rcx
	int64 (fastcall *Sleep)(_QWORD); // rax
	char Sleep_str[16]; // [rsp+28h] [rbp-10h] BYREF
:	strcpy(Sleep_str, "r-,,dT");
-	for ( $i = 0.64; i \neq 6; ++i$ )
	<pre>Sleep_str[i] = (7 * (Sleep_str[i] - 0x54) % 0x7F + 0x7F) % 0x7F; // "Sleep"</pre>
	<pre>Sleep = w_get_API_addr(kernel32_lib_struct, Sleep_str);</pre>
1	return Sleep(dwMilliseconds);
3	

```
v96 = 0i64;
do
{
  v97 = lpString[v96++];
  lpMemc = v97;
  v98 = GetProcessHeap();
  HeapFree(v98, 0, lpMemc);
}
while ( v96 ≠ 0x11 );
w_Sleep(LIB_STRUCT_ARR→lib_struct_kernel32, 3000u);
```

The way the wrapper function is setup to call the actual API is really intuitive, making the code simple to understand through static analysis. However, it's a bit more difficult to automate the process since there is no API hashing involved.

For my analysis, I just manually decode the stack strings in my debugger and rename the wrapper function accordingly.

At this point, we have fully unpacked BAZARLOADER and understood how the malware obfuscates its strings and APIs to make analysis harder.

In the next blog post, we will fully analyze how the loader downloads and launches a Cobalt Strike beacon from its C2 servers!

## 1 Comment

Comments are closed.