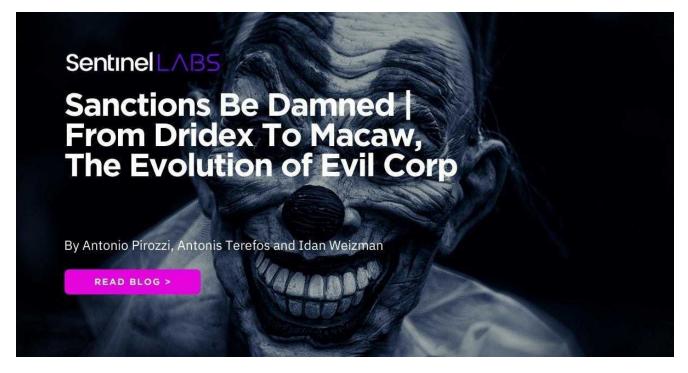
# Sanctions Be Damned | From Dridex to Macaw, The Evolution of Evil Corp

(ii) sentinelone.com/labs/sanctions-be-damned-from-dridex-to-macaw-the-evolution-of-evil-corp/

Antonio Pirozzi



By Antonio Pirozzi, Antonis Terefos and Idan Weizman

#### **Executive Summary**

- Since OFAC sanctions in 2020, the global intelligence community has been split into different camps as to how Evil Corp is operating.
- SentinelLabs assesses with high confidence that WastedLocker, Hades, Phoenix Locker, PayloadBIN belong to the same cluster. There are strong overlaps in terms of code similarities, packers, TTPs and configurations.
- SentinelLabs assesses with high confidence that the Macaw ransomware variant is derived from the same codebase as Hades.
- Our analysis indicates that Evil Corp became a customer of the CryptOne packer-as-aservice from March 2020. We created a static unpacker, <u>de-CryptOne</u> for CryptOne and identified different versions of this cryptor which have never previously been reported.

Read the Full Report

#### Introduction

Evil Corp (EC) is an advanced cybercrime operations cluster originating from Russia that has been active since 2007. The UK National Crime Agency <u>called</u> it "the world's most harmful cyber crime group." In December 2019, the U.S. Treasury Department's Office of Foreign Assets Control (OFAC) issued a <u>sanction</u> against 17 individuals and seven entities related to EC cyber operations for causing financial losses of more than 100 million dollars with Dridex.

After the indictments, the global intelligence community was split into different camps as to how Evil Corp was operating. Some assessed that there was a voluntary transition of EC operations to another 'trusted' partner while the core group remained the controller of operations. Some had theories that Evil Corp had stopped operating and that another advanced actor operated Hades, trying to mimic the same *modus operandi* as Evil Corp to mislead attribution. Others <u>claimed</u> possible attribution to the HAFNIUM activity cluster.

SentinelLabs has conducted an in-depth review and technical analysis of Evil Corp activity, malware and TTPs. Our <u>full report</u> has a number of important findings for the research community. We relied heavily on our analysis of a crypter tool dubbed "CryptOne", which supports our wider clustering of Evil Corp activity. Our research also argues that the original operators continue to be active despite the sanctions, continuously changing their TTPs in order to stay under the radar.

In this post, we summarize some key observations from our technical analysis on the evolution of Evil Corp from Dridex through to Macaw Locker and, for the first time, publicly describe CryptOne and the role it plays in Evil Corp malware development. For the full technical analysis, comprehensive IOCs and YARA hunting rules, please see the <u>full report</u>.

## **Overview of Recent Evil Corp Activity**

After the OFAC indictment, we witnessed a change in Evil Corp TTPs: from 2020, they started to frequently change their payload signatures, using different exploitation tools and methods of initial access. They switched from Dridex to the SocGholish framework to confuse attribution and distance themselves from both Dridex and Bitpaymer, which fell within the scope of the sanctions. During this period, they started relying more heavily on Cobalt Strike to gain an initial foothold and perform lateral movement, rather than PowerShell Empire.

In May 2020, a new ransomware variant appeared in the wild dubbed <u>WastedLocker</u>. WastedLocker (<u>S0612</u>) employed techniques to obfuscate its code and perform tasks similar to those already seen in BitPaymer and Dridex. Those similarities allowed the threat intelligence community to <u>identify</u> the connections between the malware families.

In December 2020, a new ransomware variant named Hades was first seen in the wild and publicly <u>reported</u>. Hades is a 64-bit compiled version of WastedLocker that displays important code and functionality <u>overlaps</u>. A few months later, in March 2021, a new variant

Phoenix Locker appeared in the wild. Our analysis suggests this is a rebranded version of Hades with little to no changes. Later, a new variant named PayloadBIN <u>appeared</u> in the wild, a continuation from Phoenix Locker.

## A Unique Cluster: BitPaymer, WastedLocker, Hades, Phoenix Locker, PayloadBIN

From our analysis, we discovered evidence of code overlaps, as well as shared configurations, packers and TTPs leading us to assess with high confidence that Bitpaymer, WastedLocker, Hades, PhoenixLocker and PayloadBIN share a common codebase. Our <u>full report</u> goes into the evidence in fine detail. The following section presents a brief summary.

#### From BitPaymer to WastedLocker

Previous <u>research</u> shows a sort of knowledge reuse between BitPaymer and WastedLocker. SentinelLabs analysis shows that Hades and WastedLocker share the same codebase.

Among other similarities, detailed in the <u>full report</u>, we observe that the RSA functions – responsible for asymmetrically encrypting the keys which were used in the AES phase to encrypt files – are identical in both ransomware variants, hinting that the same utility library was used.

#### From WastedLocker to Hades

Previous <u>research</u> assessed the main similarities and differences between the two ransomware families. SentinelLabs analysis shows that Hades and WestedLocker share the same codebase.

Again we see the same RSA functions in both families. Both also implement file and directory enumeration logic identically. Comparing the logic and the Control Flow Graph of both routines, we conclude that both ransomware use the same code for file and directory enumeration. We also found similarities between the functions responsible for drive enumeration.

#### From Hades to Phoenix Locker

In the samples we analyzed, we discovered that Phoenix Locker was a reused and newlypacked Hades payload. Hades and Phoenix samples were compiled at the same time. We confirmed that they reused a 'clean' Hades version each time, statically introducing junk code with the help of a script in order to alter the signature. The compiler and linker versions are also the same. This technique of payload reuse was also seen in BitPaymer in order to make the ransomware polymorphic and more evasive.

#### From Phoenix Locker to PayloadBIN

We observed that the majority of PayloadBIN functions overlap with PhoenixLocker. File enumerating functions are practically identical.

We conducted further similarity analysis by analyzing the TTPs of the different variants. We did this by extracting the main command lines from all the ransomwares and comparing them. We distinguished two distinct clusters.

From Hades onwards, we found a unique self-delete implementation including the waitfor command.

```
cmd /c waitfor /t 10 pause /d y & attrib -h
"C:\Users\Admin\AppData\Roaming\CenterLibrary\Tip" & del
"C:\Users\Admin\AppData\Roaming\CenterLibrary\Tip" & rd
"C:\Users\Admin\AppData\Roaming\CenterLibrary\"
```

This command is not present in WastedLocker, where the **choice** command is used instead:

```
cmd /c choice /t 10 /d y & attrib -h "C:\Users\Admin\AppData\Roaming\Wmi" & del
"C:\Users\Admin\AppData\Roaming\Wmi"
```

Whilst syntax difference may seem like a significant difference, these two implementations are very similar: the logic is the same, only the signature changes.

All ransomwares have the same implementation of Shadows copy deletion:

C:\Windows\system32\vssadmin.exe Delete Shadows /All /Quiet

The evidence of this code reuse supports the assessment that it is almost certain these ransomware families are related to the same 'factory'.

## Analysis of the Cypherpunk Variant

A new, possibly experimental, variant dubbed "Cypherpunk" – first <u>reported</u> in June 2021was analyzed and linked to the same lineage.

C:\Users\Lucas\Documents\OneNote Notebooks\Personal\General.one.cypherpunk

C:\Users\Lucas\Documents\OneNote Notebooks\Personal\CONTACT-TO-DECRYPT.txt

```
C:\Users\Lucas\Documents\awards.xls.cypherpunk
```

```
C:\Users\Lucas\Desktop\ZoneMap.dwf.cypherpunk
```

```
C:\Users\Administrator\Searches\Everywhere.search-ms.cypherpunk
```

C:\Users\Lucas\Desktop\th (2).jpg.cypherpunk

C:\Users\Lucas\Documents\pexels-photo-46710.jpeg.cypherpunk

C:\Users\Lucas\Desktop\ppt\_ch10.ppt.cypherpunk

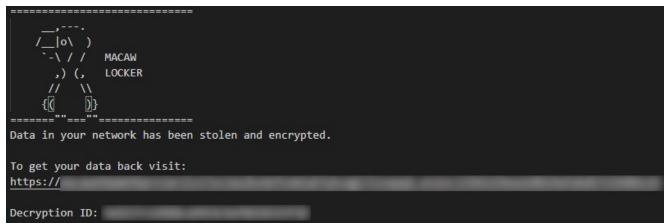
```
C:\Users\Lucas\Desktop\WEF_Future_of_Jobs.pdf.cypherpunk
```

Code similarity analysis shows that the Cypherpunk version (SHA1 <u>e8d485259e64fd375e03844c03775eda40862e1c</u>) is the same as the previous PayloadBIN variant. It was compiled on 2021-04-01 17:15:24, 20 days after the PayLoadBIN sample. It is possible that this is another attempt at rebranding. Although this variant was reported, it was <u>improperly flagged</u> as Hades.

SentinelLabs assesses this new finding is likely an indication that Evil Corp is still working on updating their tradecraft in order to change their signature and stay under the radar.

## **Evil Corp Pivots to Macaw Locker Ransomware**

In October 2021, a new ransomware variant named 'Macaw Locker' appeared in the wild, in an attack that began on October 10th against <u>Olympus</u>. A few days later <u>Sinclair Broadcast</u> <u>Group</u> was also attacked, causing widespread disruption. Some researchers <u>claimed</u> a possible connection with WastedLocker, but to date no further details have emerged.



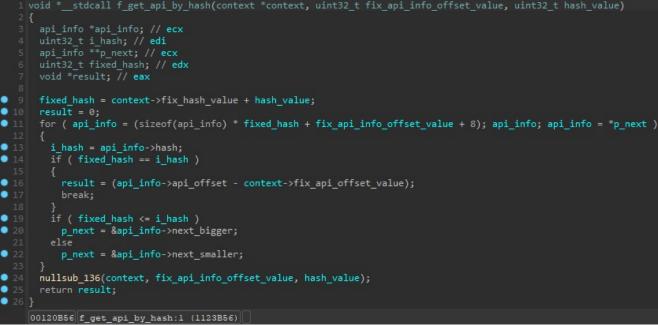
#### Macaw ransom note

The ransomware presents anti-analysis features like API hashing and indirect API calls with the intention of evading analysis. One aspect that immediately sets Macaw apart is that it requires a custom token, provided from the command line, which appears to be specific to each victim; without it, the ransomware won't execute.

macaw\_sample.exe -k

The use of a custom token is also seen in <u>Egregor</u> and <u>BlackCat</u> ransomware families, and is a technique used to aid anti-analysis (<u>T1497.002</u>).

Another new addition to Macaw is a special function that acquires the imports for APIs at runtime, instead of when the executable is started via the PE import section. Below, we can see the function that is used before each API call to get its address prior to the call itself.



Macaw function to dynamically fetch addresses

The function gets a 32-bit value that uniquely represents the required API and searches for it through a data structure created beforehand. The data structure can be described as an array with small binary search trees in each of its entries.

We assessed the similarity of two core functions between Hades and Macaw. In both strains, the implementation is the same. The only minor differences are from the imports fetched at runtime.

## CryptOne: One Packer To Rule Them All

CryptOne (also known as HellowinPacker) was a special packer used by Evil Corp up until mid-2021.

CryptOne appears to have first been noticed in <u>2015</u>. Early versions were used by an assortment of different malware families such as <u>NetWalker</u>, Gozi, Dridex, <u>Hancitor</u> and Zloader. In 2019, Bromium analyzed and reported it as in use by <u>Emotet</u>. In June 2020, <u>NCC</u> <u>Group</u> reported that CryptOne was used to pack WastedLocker. In 2021, <u>researchers</u> observed CryptOne being advertised as a Packer-as-a-Service on various crime-oriented forums.

CryptOne has the following characteristics and features:

- Sandbox evasion with getInputState() or GetKeyState() API;
- Anti-emulation with UCOMIEnumConnections and the IActiveScriptParseProcedure32 interface;
- Code-flow obfuscation;

We created a static unpacker, <u>de-CryptOne</u>, which unpacks both x86 and x64 samples. It outputs two files:

- 1. the shellcode responsible for unpacking
- 2. the unpacked sample.

We collected CryptOne packed samples, and with the use of the above tool, <u>unpacked and</u> <u>categorized them at scale</u>.

### Unpacking CryptOne

CryptOne unpacking method consists of two stages:

- 1. Decrypts and executes embedded shellcode.
- 2. Shellcode decrypts and executes embedded executables.

CryptOne gets chunks of the encrypted data, which are separated by junk.

Address	He																ASCII
			AF	00	54	64	74	50	61		62	41	77		72	65	
0000000004011c6																	.1TAtPalcAwAre
00000000004011D6															61	6C	0;÷irG0al
00000000004011E6			6C											-		69	R1100;;;Vi
00000000004011F6			75	61		E6			EE				13		00	00	!Ouaÿæreîi
000000000401206			00	55							_		CC	-		69	.i.UyEap%Cew1ÆFi
000000000401216				_		_	_		DD				C6	BF	_	50	çĂįÝÇrtÆ¿lP
000000000401226			74			D4			13				13			61	Afte 0jioa
000000000401236			69											A1			Tib.AryNÙA;MM
000000000401246	_		4D		_	_	_	_	_	_							MMMMMMMMMMMMMMMM
000000000401256			00											CD		48	MM;Gegiod^1eH
000000000401266															61		rīdlnāį.CAĀat
000000000401276												00	13	A1		00	6çil.àjj
000000000401286			74						FC				EE	D2		00	@ÂtF*ÍePüÇntîÔ
000000000401296															00	00	.iöriçÅFi÷Å
00000000004012A6								00	13	A1	43	6C	B4	D1	65	48	.iiiCl´ÑeH ≛1dl®ii.G
00000000004012B6		CE	64	6C	AE	A0	00	00	13	A1	00	00	13	A1	00	47	*1dl@jj.G
00000000004012c6	<b>B</b> 6	D4	54	65	A6	D0	50	61	87	C9	41	00	4D	4D	4D	4D	OTE DPa. ÉA. MMMM
0000000004012D6	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	4D	MMMMMMMMMMMMMMM
00000000004012E6	13	A1	00	00	13	A1	00	00	67	<b>D</b> ()	74	77	7F	C6	6E	41	. j jgDtr. AnA
00000000004012F6	13	A1	00	00	13	A1	00	00	13	A1	00	00	13	CD	73	74	. j j j İst
000000000401306	41	C0	61	74	52	A1	00	00	13	A1	00	00	13	A1	00	00	AAatRjjj
000000000401316	13	A1	52	74	2F	E2	64	64	4D	D6	6E	63	07	CA	6F	6E	. ¡Rt/âddMônc.Éon
000000000401326	DF	C1	62	6C	F6	A0	00	4D	EE	CF	73	61	F4	C5	42	6F	BÁblö .MîïsaôABo
000000000401336	EB	E1	00	00	13	A1	00	00	13	A1	00	00	D4	C5	74	50	ëájjÔÁtP
000000000401346	C1	CB				C4		65	C0	CF	00	00	DF	CB	61	64	AEcA AreA1. BEad
000000000401356	D7	C9	62	72	<b>B2</b>	CE	79	45	9B	E1	4D	4D	4D	4D	4D	4D	×Ébr²ÎyE. áмммммм
000000000401366			4D														MMMMMMMMMMMMMM.
000000000401376		C6	72	6E	6E	CD	33	32	2D	C5	6C	6C	13	AL	00	00	x&rnnI32-All.i
000000000401386			73									00	13	A1	00	00	AÆssrÄeBlùA
0000000000401396														AI			FDerd/1

CryptOne junk data Example Memory Dump:

- 0x5EE00, Encrypted size
- 0x4011CA, Address of encrypted data
- 0x4D/"M", Junk data
- 0x14, Junk size
- 0x7A, Chunk Size

After removal of the junk data, the decryption starts with a simple XOR-Key which increases by 0x4 in each round. The initial XOR-Key is **0xA113**.



#### CryptOne XOR Key

Once the shellcode is decrypted, we can partially observe the string "This program cannot be run in DOS mode" where this data contains an executable which requires a second decryption.

47		74			6F	63	41	64	64	72	65	73			00	GetProcAddress
00					56		72		75	61	6C		6C		6F	VirtualAllo
63	00		00			00			00	56			74	75	61	cVirtua
GC			65		00	00	00	00	00	00	00	00	00	00	55	1FreeU
6E	GD	61			69	65	77	4F	66	46	69	6C	65	00	00	nmapViewOfFile
00	00	00	00	56	69	72	74	75	61	6C	50	72	6F	74	65	VirtualProte
63	74	00	00	00	00	00	00	00	4C	6F	61	64	4C	69	62	ctLoadLib
72	61	72	79	45	78	41	00	00	00	00	00	00	00	47	65	raryExAGe
74	4D	6F	64	75	6C	65	48		6E	64	6C	65	41	00	00	tModuleHandleA
00	00	00	43	72	65	61	74	65	46	69	6C	65	41	00	00	CreateFileA
00	00	00	00	00	00		00		65	74	46		6C	65	50	SetFileP
6F	69	6E	74		72		00		00	00	00		57	72	69	ointerWri
74			69	6C		00			00	00	00	00	00	00	00	teFile
00	00		6C		73		48	61	6E	64	6C		00	00	00	CloseHandle
00	00	00	00	00	00	00	47	65	74	54	65	6D	70	50	61	GetTempPa
74	68	41	00	00	00	00	00		00	00	00		73	74	72	thAlstr
6C	65	6E	41	00	00	00	00	00	00	00	00		00	00	00	1enA
00	6C	73	74	72	63	61	74	41	00	00	00		00	00	00	.1strcatA CryptOpe portiolly
00	00	00	00	00	00	52	74	GC	41	64	64	46	75	6E	63	CryptOne partially
74	69	6F	6E	54	61	62	6C	65	00	00	4D	65	73	73	61	tionTableMessa
67	65	42	6F	78	41	00	00	00	00	00	00	00	00	00	00	geBoxA
47	65	74	50	72	6F	63	41	64	64	72	65	73	73	00	00	GetProcAddress
4C	6F	61	64	4C	69	62	72	61		79		78	41	00	00	LoadLibraryExA
6B	65	72	6E	65	6C	33	32	2E	64	6C	6C	00	00		00	kernel32.dll
4D			73	61	67	65	42	<u>6</u> F	78	41	00	_		00		MessageBoxA
75					32		64	6C	6C	00				00		user32.d11
01	00					00		02	00	00	00			00	00	
10	00	00	00	80	00	00	00	20	00	00		40	00	00	00	·····
48	83	C4	40	FF	E1	00	DC	05	00	A4	59	90	00	EA	03	H.A@ÿa.U¤Yê.
00			03	00	00	FE	FB	00	00	31	03	00	00	E9	03	iþü1é.
00	00		04	00	00	E9	03	00	00	E9	03	00	00	E9	03	)éé.
00	00		03	00	00	E9	03	00	00	E9		00		E9	03	eee.
00	00		03	00	00	F1	04	00	00	E7	1A	-	0E	E9	AF	enç.º.e
09	CD	C8	BB	01	4C	AC	25	54	68	00	77	20	70	FB	6A	.İÈ».L¬%Th.w_pûj
67	72	C8	68	20	63	C8	69	6E	6F	DD	23	62	65	09	76	grÈh cÈinoY#be.v
75	6E	09	GD	6E	20	A5	4A	53	20	C4	6A	64	65	FF	08	un.mn ¥JS Aidey.
00	~ *		0.0	00			0.0	00				22	20	2.2		

#### decrypted shellcode

Similar to previous decryption, this time the shellcode decrypts the embedded binary.

int64fastcall shellcode_xor(int64 exe {	c, unsigned int exec_size)
<pre>int64 result; // rax unsigned int i; // [rsp+0h] [rbp-18h]</pre>	
for $(i = 0; ; i += 4)$	
<pre>{     result = exec_size;     if ( i &gt;= exec_size )         break; }</pre>	Fastcall
*(_DWORD *)(exec + i) += i;	
<pre>*(_DWORD *)(exec + i) ^= i + 0x3E9; } return result; }</pre>	// Initial XOR-Key 0x3E9

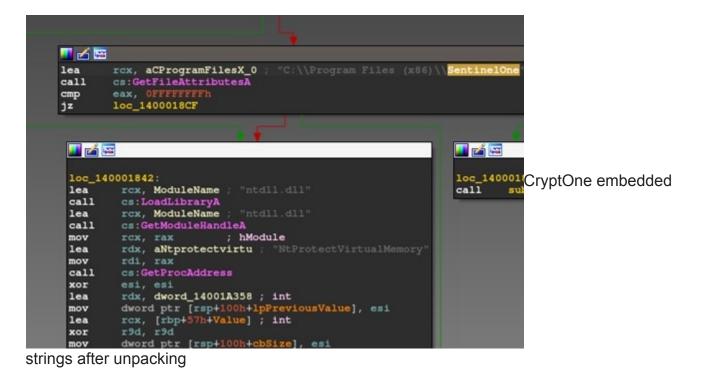
#### Shellcode XOR

The shellcode allocates and copies the encrypted executable and starts the decryption loop; once it finishes, it jumps to the EntryPoint and executes the unpacked sample.

<pre></pre>	[rsp+20]:"MZ" [rsp+20]:"MZ"	
		CryptOne
Dump 2 💭 Dump 3 📜	Dump 4 🛄 Dump 5 👹 Watch 1 📧 Loca	
Hex	ASCII	l l
B8         00         00         00         00         00         00           00         00         00         00         00         00         00         00           00         00         00         00         00         00         00         00         00           00         00         00         00         00         00         00         00         00           00         00         00         00         00         00         00         00         00           0E         1F         BA         0E         00         B4         09           69         73         20         70         72         6F         67           74         20         62         65         20         72         75           6D         6F         64         65         2E         0D         0D           9E         C6         32         8D         DA         A7         5C           81         CF         58         DF         DA         A7         5C           DA         A7         5D         DE         BA         A7         5C </td <td>00 04 00 00 00 FF FF 00 00 MZÿÿ 00 40 00 CD 21 88 01 4C CD 21 54 68°Í!.LİIT 72 61 6D 20 63 61 6E 6E 6F is program canno 6E 20 69 6E 20 44 4F 53 20 t be run in DOS 0A 24 00 00 00 00 00 00 00 mode\$ DE DA A7 5C DE DA A7 5C DEĨXBD§\DI_BB§\D DE 81 CF 5F DF DF A7 5C DEĨYB]§\DI_BB§\D DE 81 CF 5D DF D9 A7 5C DEÏYB]§\D.Ï_BB§\D DE D1 C8 59 DF FF A7 5C DEÏYB]§\D.Ï_BB§\D DE D1 C8 5F DF D2 A7 5C DEÏYB]§\D.Ï_BD§\D</td> <td></td>	00 04 00 00 00 FF FF 00 00 MZÿÿ 00 40 00 CD 21 88 01 4C CD 21 54 68°Í!.LİIT 72 61 6D 20 63 61 6E 6E 6F is program canno 6E 20 69 6E 20 44 4F 53 20 t be run in DOS 0A 24 00 00 00 00 00 00 00 mode\$ DE DA A7 5C DE DA A7 5C DEĨXBD§\DI_BB§\D DE 81 CF 5F DF DF A7 5C DEĨYB]§\DI_BB§\D DE 81 CF 5D DF D9 A7 5C DEÏYB]§\D.Ï_BB§\D DE D1 C8 59 DF FF A7 5C DEÏYB]§\D.Ï_BB§\D DE D1 C8 5F DF D2 A7 5C DEÏYB]§\D.Ï_BD§\D	

executing the unpacked sample

At this stage we can observe strings related to the unpacked sample.



## A Unique Factory

Hunting for CryptOne led us to identify different implementations of the stub, some of which have never been reported previously. Each version is identified by a certain signature, listed below:

- 11111111\\{aa5b6a80-b834-11d0-932f-00a0c90dcaa9}
- 1nterfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- 444erfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- 555erfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- 5nterfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- 987erfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- Interfac4\\{b196b287-bab4-101a-b69c-00aa00341d07}
- InterfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- aaaerfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- interfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
- rrrerfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}

The first part of the string is composed of a custom string (11111111, 1nterfacE, 444erfacE,...) which is replaced at runtime by the 'interface' keyword, creating the following registry key:

HKEY\_CLASSES\_ROOT\interface\{b196b287-bab4-101a-b69c-00aa00341d07}

The registry keys are related to the UCOMIEnumConnections and IActiveScriptParseProcedure32 interfaces respectively.

Once executed, the cryptor checks for the presence of those keys before loading the next stage payload. If it does not find the keys, then the malware goes into an endless loop without doing anything as an anti-emulation technique. This works because some emulators do not implement the full Windows registry.

In reviewing two different versions of CryptOne:

aaerfacE\\{b196b287-bab4-101a-b69c-00aa00341d07}
11111111\\{aa5b6a80-b834-11d0-932f-00a0c90dcaa9}

we noticed that in order to update the signature, the actor needs to re-compile the cryptor as the cryptor implementation changes.

## CryptOne Timeline

Our analysis shows that it is likely Evil Corp started being a customer of the CryptOne service from March 2020. From March to May 2020 we found WastedLocker, gozi\_rm3 (version:3.00 build:854) and Dridex (10121) samples were all packed and compiled in the same timeframe using the same CryptOne stub signature(InterfacE).

For a limited period of time between May 2020 and August 2020, we observed different versions of CryptOne overlaps.



CryptOne overlaps between May 2020 and August 2020

It seems that from a specific point in time, around September 2020, Hades, PhoenixLocker and PayloadBIN started adopting a specific CryptOne stub identified by the signature:

From December 2020, the CryptOne version '111111111' appeared in the wild without any overlap.

## Conclusion

Clustering Evil Corp activity is demonstrably difficult considering that the group has changed TTPs several times in order to bypass sanctions and stay under the radar. This is in addition to the overall trend of actors receding back into secrecy. In this research, we connect the dots in the Evil Corp ecosystem, cluster Evil Corp malware, document the group's activities and provide insight into their TTPs.

SentinelLabs assesses with high confidence that WastedLocker, Hades, PhoenixLocker, Macaw Locker and PayloadBIN belong to the same cluster. Our assessment is based on code similarity and reuse, timeline consistency and nearly identical TTPs across the ransomware families indicating there is a consistent *modus operandi* for the cluster. In addition, we assess that there is a likely evolutionary link between WastedLocker and BitPaymer, and suggest that it can be attributed to the same Evil Corp activity cluster.

We fully expect that Evil Corp will continue to evolve and target organizations. In addition, we assess it is likely they will also continue to advance their tradecraft, finding new methods of evading detection and misleading attribution. SentinelLabs will continue tracking this activity cluster to provide insight into its evolution.

In-depth technical analysis, Indicators of Compromise and further technical references are available in the full report.

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