

Technical Analysis of Code-Signed “Blister” Malware Campaign (Part 2)

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The blister is a code-signed malware that drops a malicious DLL file on the victim’s system, which is then executed by the loader via `rundll32.exe`, resulting in the deployment of a RAT/C2 beacon, thus allowing unauthorized access to the target system over the internet. Blister Malware campaigns have been active since 15 September 2021.

Part I of CloudSEK’s analysis provides a detailed understanding of how the loader functions. Part 2 will delve into the details of this campaign’s second stage, which is the `.dll` payload, and its internal working.

Dissecting the Malicious DLL – Blister Malware

As discussed in Part 1, the Blister dropper drops the malicious `.dll` file in the `Temp directory` of the user, inside a newly created folder. This malicious `.dll` then carries out the second stage of the campaign, in which a RAT/ agent is deployed on the system to gain unauthorized access and steal data.

The Blister dropper calls the function `LaunchColorCpl`, which is one of the functions exported by the `.dll`, via `rundll32.exe`.

Ordinal	Function RVA	Name Ordinal	Name RVA	Name
(nFunctions)	Dword	Word	Dword	szAnsi
00000001	00003A6C	0000	00011E20	LaunchColorCpl
00000002	000037AA	0001	00011E2F	DllCanUnloadNow
00000003	00005D62	0002	00011E3F	DllGetClassObject
00000004	000039F5	0003	00011E51	DllMain
00000005	000037D1	0004	00011E59	DllRegisterServer
00000006	000037DB	0005	00011E6B	DllUnregisterServer

Functions exported by the malicious DLL

Staging

The exported function `LaunchColorCpl` retrieves the staging code from the resource section of the PE file. This staging code is protected by a simple XOR encoding scheme.

```

1717390E  FFD7          call edi
17173910  8BC6          mov eax,esi
17173912  83E0 03       and eax,3
17173915  8A4405 E8     mov al,byte ptr ss:[ebp+eax-18]
EIP -> 17173919  30041E       xor byte ptr ds:[esi+ebx],al
1717391C  46           inc esi
1717391D  81FE E0890100 cmp esi,189E0
17173923  72 EB       jb holorui.17173910
17173925  8D45 DC     lea eax,dword ptr ss:[ebp-24]

```

Code responsible for decoding the staging code

Address	Hex	ASCII
172040E6	50 00 00 E8 35 24 E8 8F 1A D4 F8 66 72 8A E8 8F	P...è5\$è..òofr.è.
172040F6	C9 F9 56 BA 98 10 BA 66 69 2B E9 8F C9 F8 F3 AB	ÈuV°...°fi+é.Èóó«
17204106	99 10 68 4B 89 F9 68 CF 98 10 00 81 BD 10 E8 0C	..kk.ùhI...%.è.
17204116	5D 00 01 72 05 10 E8 06 DC C8 01 57 B6 10 E8 DE]..r...è.ÛÈ.w].èp
17204126	70 92 C7 8F 99 93 95 63 99 1F 6D 52 C8 11 E8 04	p.Ç...c...mRE.è.
17204136	1C 8C 1F 70 66 F9 EA D6 98 10 2B 0C E4 F4 E8 66	..pFüèÖ...+..ãòèf
17204146	26 1C E8 8F 5E 55 2C 8F 99 10 E8 0F 62 A5 6B F2	&.è.ÀU...è.b≠kò
17204156	5D 10 01 A0 89 10 E8 67 2B 8D E8 8F 1A D4 D0 06]...èg+.è..òð.

Encoded staging code in the resource section of the PE file

- After the iterative decoding of the staging code, the control is transferred to decoded code in the memory.
- The control flow is transferred to the staging code by calling the address in the EAX register.

```

17173910 8BC6 mov eax,esi
17173912 83E0 03 and eax,3
17173915 8A4405 E8 mov al,byte ptr ss:[ebp+eax-18]
17173919 30041E xor byte ptr ds:[esi+ebx],al
1717391C 46 inc esi
1717391D 81FE E0890100 cmp esi,189E0
17173923 72 EB jb holorui.17173910
17173925 8D45 DC lea eax,dword ptr ss:[ebp-24]
17173928 50 push eax
17173929 6A 20 push 20
1717392B 8D45 F0 lea eax,dword ptr ss:[ebp-10]
1717392E 50 push eax
1717392F 8D45 EC lea eax,dword ptr ss:[ebp-14]
17173932 50 push eax
17173933 6A FF push FFFFFFFF
17173935 FFD7 call edi
17173937 8D45 E8 lea eax,dword ptr ss:[ebp-18]
1717393A 50 push eax
1717393B 8D83 905A0000 lea eax,dword ptr ds:[ebx+5A90]
17173941 FFD0 call eax
17173943 5F pop edi
17173944 5B pop ebx
17173945 33C0 xor eax,eax

```

Calling the address in the EAX register

Anti-Analysis

- The staging code is heavily obfuscated, and has a logic similar to a spaghetti code, to hinder analysis. All the calls to Windows APIs are obscured and dynamically resolved.
- The first thing that the staging code does is to make the malware go to sleep by calling the Sleep Windows API. This is a typical strategy used by most malicious codes to bypass security sandboxes and dynamic testing of security products.
- The hex value “927Co” is passed to `kernel32.759F9010` i.e the Sleep function. This value (927Co) translates to “600000” in decimal. Since the Sleep API takes arguments in milliseconds (ms), the 600000 ms get converted to 10 minutes.
- When the malware resumes from sleep, it fetches the final payload from the resource section of the PE file.

```

009FEA18 759F9010 kernel32.759F9010
009FEA1C 000927C0
009FEA20 00000000
009FEA24 009FEA5C
009FEA28 00000000
009FEA2C E20400E6

```

Stackframe before the malware calls the Sleep Windows API

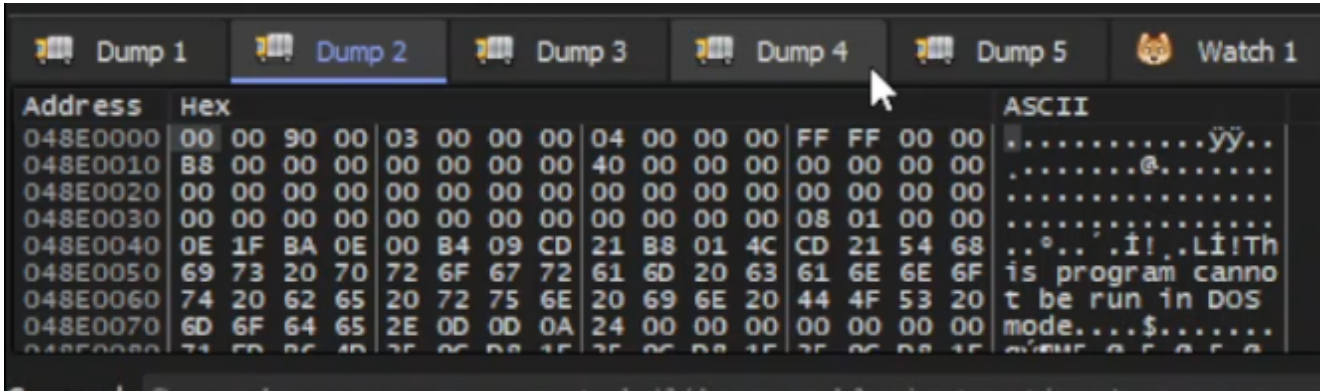
```

Address  Hex  ASCII
04880000  C8 33 11 25 32 0D 8A E2 53 01 C0 1E A5 CC EB 88  È3.%2..âs.A.viè.
04880010  88 7E CE 2C 8B 14 01 23 5E 53 96 23 80 9D A7 69  .~I,...#^S.#.;$i
04880020  66 D1 89 8D 86 DB 6B C8 A8 F2 09 22 19 60 A0 10  fN...ÙkÈ"ò.". .
04880030  64 B7 56 78 5F AF 8A 42 E7 E7 1B A0 3D 72 09 97  d.Vx_ .Bcc. =r..
04880040  AA DC 25 FF 21 3C EB CA FB BF 31 EE F7 F0 B3 C7  *Ùxy!<èÈÙ;1i÷ð*C
04880050  4B E0 FA A0 65 96 CA 01 E7 A3 41 15 32 0B 7E F0  Kàu e.È.cíA.2.~ð
04880060  EF F0 79 74 A9 FA F5 A5 D3 3B 16 62 82 84 D8 98  iðyt@úð%0;.b='ò.
04880070  C7 CF 3D 18 DD D7 9B 28 FF 89 BE 3F 91 F6 2A D0  CI=.ÿx.(ÿ.%?.ò*D

```

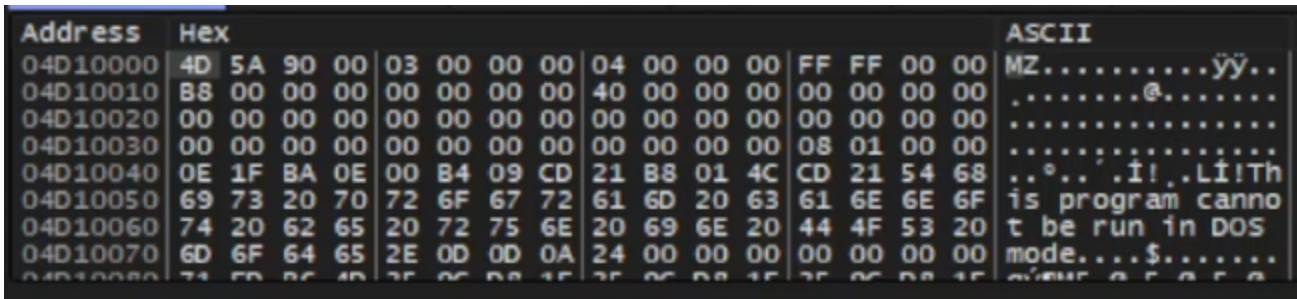
Snippet of the protected payload stored in the memory

In the memory, the protected payload is decoded. The presence of a DOS header, in the payload bytes, confirms that the payload is in PE format and not a shellcode.



Decrypted payload stored in the memory

An interesting observation from this analysis, is the addition of MZ byte after the decryption process. In the above image, the initial byte is not MZ, rather the MZ byte is later added at the beginning of the payload separately. This behavior is primarily for operational security.

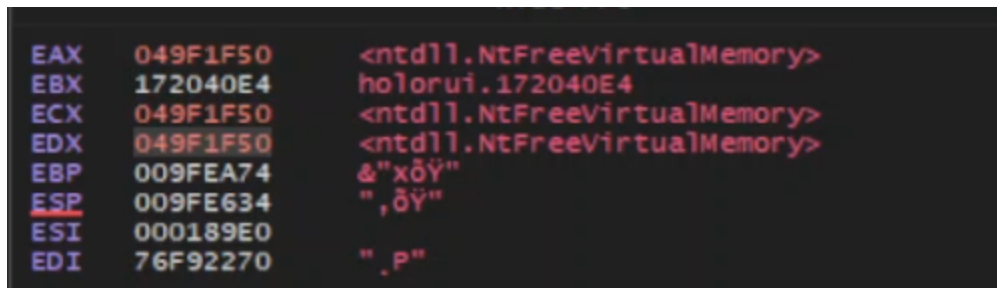


Addition of the MZ byte after the decryption process

Process Hollowing

In general, process hollowing allows an attacker to change the content of a legitimate process from genuine code to malicious code before it is executed by carving out the code logic within the target process.

- After decrypting the final payload, the malware prepares for execution.
- This is done by creating a new process to deploy the extracted code and then performing process hollowing to execute the payload in the remote process. The staging code retrieves the *Rundll32.exe* location from the compromised system.



Retrieval of the location of rundll32.exe

A new process of *Rundll32.exe* is created via the *CreateProcessInternalW* API in the suspended state.

```
EAX 049F35C0 <ntdll.ZwSetContextThread>
EBX 172040E4 holorui.172040E4
ECX 049F35C0 <ntdll.ZwSetContextThread>
EDX 049F35C0 <ntdll.ZwSetContextThread>
EBP 009FEA74 &"xδY"
ESP 009FE63C ",δY"
ESI 000189E0
EDI 76F92270 ".P"
EIP 172166D8 holorui.172166D8
```

Creation of the new rendll32.exe

- The malware uses the following Win32 APIs for process hollowing:
 - *ZwUnmapViewOfSection*
 - *ZwReadVirtualMemory*
 - *ZwWriteVirtualMemory*
 - *ZwGetContextThread*
 - *ZwSetContextThread*
 - *NtResumeThread*
- *ZwWriteVirtualMemory* is used to write malicious code into the target process.
- To make the thread of the new process point to newly written code, the attacker alters the entry point of the current thread via *ZwGetContextThread* and *ZwSetContextThread*.
- These functions are used to perform processor housekeeping activities on the data structure that stores the current context of the running thread. Process hollowing takes advantage of these features to make the process thread run the attacker code.

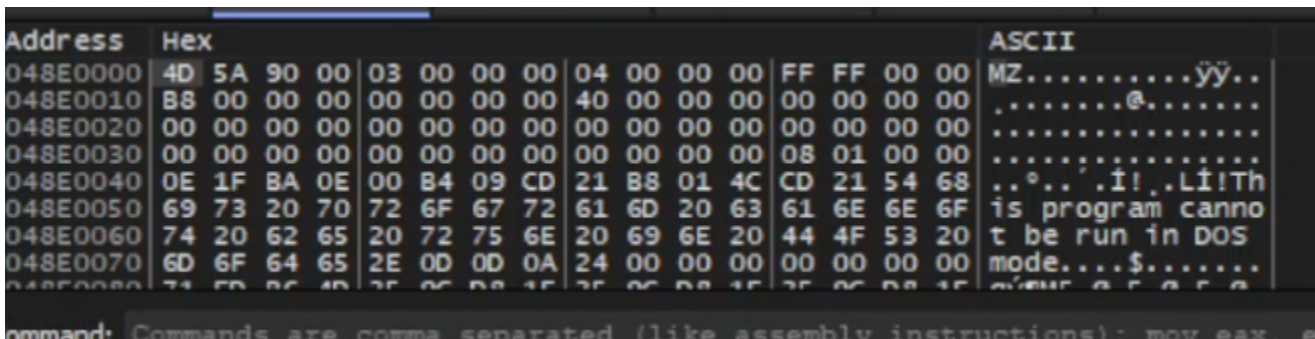
Step by Step Working of the DLL

The staging code allocates a new memory via *ZwAllocateVirtualMemory* to transfer the previously decrypted final payload.

```
EAX 049F1ED0 <ntdll.ZwAllocateVirtualMemory>
EBX 172040E4 holorui.172040E4
ECX 049F1ED0 <ntdll.ZwAllocateVirtualMemory>
EDX 049F1ED0 <ntdll.ZwAllocateVirtualMemory>
EBP 009FEA74 &"xδY"
ESP 009FE62C ",δY"
ESI 000189E0
EDI 76F92270 ".P"
EIP 17207598 holorui.17207598
```

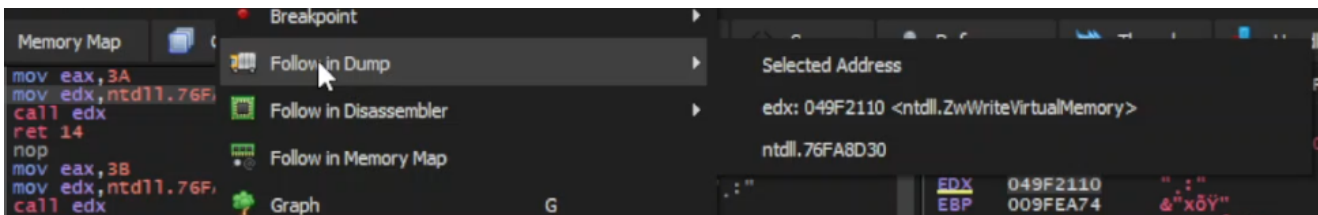
Allocation of new memory via ZwAllocateVirtualMemory

The payload is then copied to a newly created buffer.. Based on CloudSEK's testing on the extracted payload, one of the analyzed samples contained the *Raccoon stealer* as the final stage payload. However, other samples used *Cobalt Strike beacon* and *BitRAT* to compromise the target and gain unauthorized access.



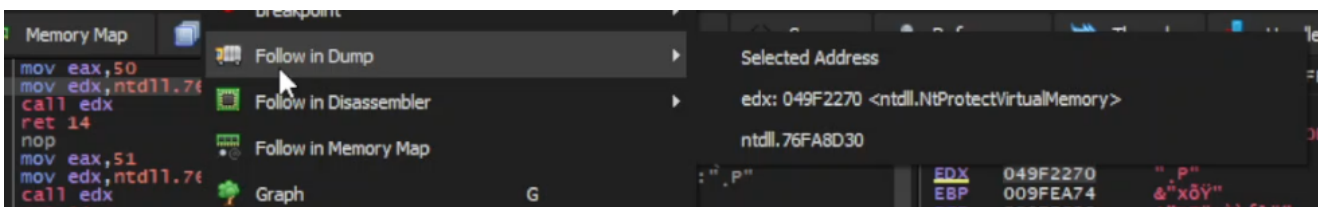
Moving the payload to a newly created buffer

The staging code then injects the code into the newly created remote process i.e *Rundll32.exe*.



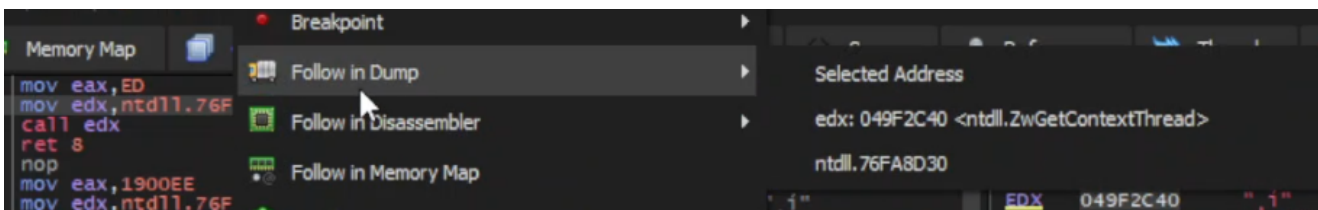
Code injections into the newly created rendll32.exe

Later, the memory protections are changed to appropriate ones for the execution of the residing code via *NTPProtectVirtualMemory*.



Alteration of the memory protections

The thread context is retrieved via *ZwGetContextThread* API to change the entry point of the thread to execute the payload injected into the remote process.



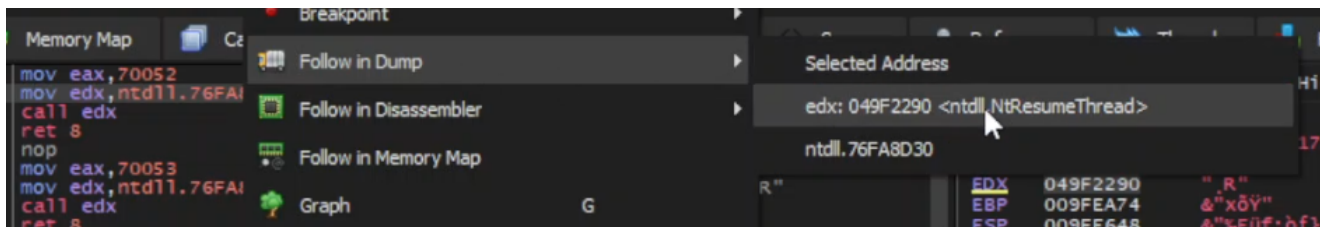
Addition of the MZ byte after the decryption process

The `ZwSetContextThread` is used to modify the thread entry point to that of the newly copied PE file.

```
EAX 049F35C0 <ntdll.ZwSetContextThread>
EBX 172040E4 holorui.172040E4
ECX 049F35C0 <ntdll.ZwSetContextThread>
EDX 049F35C0 <ntdll.ZwSetContextThread>
EBP 009FEA74 &"xØY"
ESP 009FE63C ",ØY"
ESI 000189E0
EDI 76F92270 ".P"
EIP 172166D8 holorui.172166D8
```

Modification of the thread entry point to the copied PE file

At the final stage of process hollowing, the suspended thread of the `Rundll32.exe` is resumed via `NtResumeThread`. Then the `Rundll32.exe` process starts executing the malicious code hollowed into it by the malware.



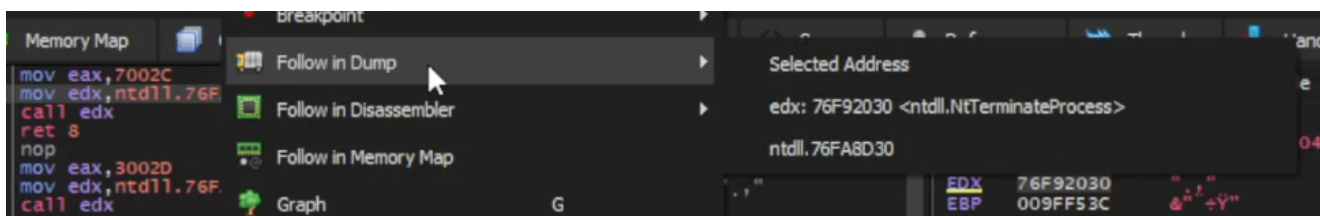
Resuming the suspended thread

In the clean-up process, the staging code uses `NtFreeVirtualMemory` to release the allocated memory, which holds the payload assembly, one by one.

```
EAX 049F1F50 <ntdll.NtFreeVirtualMemory>
EBX 172040E4 holorui.172040E4
ECX 049F1F50 <ntdll.NtFreeVirtualMemory>
EDX 049F1F50 <ntdll.NtFreeVirtualMemory>
EBP 009FEA74 &"xØY"
ESP 009FE634 ",ØY"
ESI 000189E0
EDI 76F92270 ".P"
```

Clean-up process releasing the allocated memory



The current process used for staging is terminated via the `NtTerminateProcess`.



Termination of the current process

Blister Malware – Maintaining Persistence





- The Blister malware achieves persistence on the target system by creating an “lnk” file named `proamingsGames` in the `C:\Users\
<username>\AppData\Roaming\Microsoft\Windows\Start Menu\Startup` directory.
- Whenever the user logs in, `explorer.exe` executes any file in the `Startup` folder. As a result, when the user signs into the account, following the boot process, the malware runs as a child process of `explorer.exe`.

Name	Date modified	Type	Size
 <code>proamingsGames</code>		Shortcut	1 KB

Ink file produced in the Startup directory

The target for the lnk file is set as

`C:\ProgramData\proamingsGames\proamingsGames.dll,LaunchColorCpl`. Here, the malware copies the `Rundll32.exe` as `proamingsGames.exe` and the malicious .dll (initially into `C:\ProgramData\proamingsGames` directory) is dropped in the `Temp` folder.

Name	Date modified	Type	Size
 <code>proamingsGames.dll</code>		Application exten...	1,114 KB
 <code>proamingsGames</code>		Application	61 KB

Contents of the proamingsGames.dll file

Every time that the system powers up and the user logs in, the lnk file runs a malicious `.dll` through a renamed instance of `Rundll32.exe`.

Conclusion

Given that threat actors are actively using valid code-signing certificates in Windows systems, to avoid detection by antivirus software, it is essential for network and endpoint security products to be updated with the malwares' latest Indicators of Compromise (IoCs). The latest IoCs for the Blister Malware are enumerated in [Part 1 of the technical analysis](#).

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