Buer Loader Analysis, a Rusted malware program

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January 20, 2022

Malware analysis is part of the CTI team's daily routine. This article presents the analysis of a Rust strain of Buer Loader from the reception of the samples to the writing of a stage2* extraction script. Despite several protection mechanisms, it was possible to extract all the samples in different ways. TEHTRIS provides the code for such an extraction.

Summary

Introduction

The Rust language [1] is more and more used [2] by developers. Indeed, the philosophy of this language is to maximize security while offering performances close to C++ in a pleasant syntax. Many ambitious projects like Kerla [3] aim at offering an operating system that is compatible with any Linux binary, without modification. This gives Rust a lot of credibility and explains its adoption by developers. From a low-level point of view, this language generates a more complex binary code to sign and analyze compared to C by offering more instructions and adding obfuscation at low cost, which is of interest to malware authors.

If still few adopt it, it is however not surprising to note that this language is starting to make a name for itself in the malware bestiary, as shown by Buer Loader [3] which includes variants developed in this language.

If this malware program has already been analyzed [3], it remains interesting to study its "stage1*" to identify the protection mechanisms and the evolutions concerning its obfuscation over time. Indeed, sandbox bypass techniques [5] and memory loading of offended binaries [6] by an unknown mechanism have been observed. This was enough to trigger the interest of a retro-analysis of the loader in question.

The goal is to characterize and extract the main load of the malware that was specifically spotted in computing environments between July and August 2021.

An emulation-based extraction method was presented at the Hack-It-N conference in December 2021. A replay of the presentation is available on Youtube [16].

Analysis

The retrieved samples work on the same model, i.e. a base in Rust loading in memory a stage2, itself written in Rust. Here is the list of these samples:

The compilation dates seem consistent with each other. Indeed, they match approximately the submission dates on Virus Total [15] (with a small delay). The binaries are stripped (i.e. symbols unnecessary for the binary to work have been removed, especially function names and debugging symbols) but still contain information about the used source code:

virtual-machine:/dev/shm<mark>\$ strings 88689636f4b2287701b63f42c12e7e2387bf4c3ecc45eeb8a61ea707126bad9b.exe | grep '^src.*rs\$</mark>

Unfortunately, BinDiff [7] is not able to compare binaries. The control flow is too complex for it. Therefore, the code comparisons were done manually by our experts.

Our analysis determined that there seem to be only 2 source files outside of the open source (mainly cryptographic) libraries used. A first anti-sandbox technique consists in wasting time by performing an unnecessary loop. The order of magnitude of the time needed (and lost) is about one minute with a 100% occupied CPU.

Anti-sandbox by waste of time

After passing through this loop, we arrive at the decryption function of stage2. This one contains a large number of successive calls to useless functions, among which we find:

- **GetCommandLineA;**
- **GetCurrentProcess;**
- **GetEnvironmentStrings;**
- **GetLastError;**
- **GetProcessHeap;**
- **GetTickCount;**

It is probably a saturation mechanism using calls (call to a Windows library) and not taking any argument to override possible sandboxes, which complicates code emulation. It seems that a script generates all the calls. This technique also has the advantage of adding a credible call/instruction ratio compared to a legitimate program. Counting the calls sometimes allows to find decryption or unpacking routines:

 \sim \sim

- GetTickCount $call$
- $call$ GetTickCount
- $call$ GetEnvironmentStrings
- call GetProcessHeap
- $call$ GetCurrentProcess
- call GetCurrentProcess


```
push
        esi
sub
        esp, 0AACh
        esi, ecx
mov
        exc, 1mov
        RE_DECOY
call
        exc, 3
mov
        RE DECOY
call
mov
        exc, 2
        RE_DECOY
ca11exc, 2mov
        RE_DECOY
call
        ecx, 7
mov
        RE_DECOY
call
mov
        exc, 2
cal1RE DECOY
mov
        exc, 7RE_DECOY
call
        exc, 2mov
        RE DECOY
call
        ecx, 4
mov
        RE_DECOY
call
        ecx, 6
mov
call
        RE_DECOY
```


Decoys variant 2

These calls are repeated throughout the desobfuscation routine. The data are successively pushed on the stack, then in the heap with the help of kernel32 !HeapAlloc, in the form of DWORD:

Stage2 buffer reconstruction

The call graph is very different across samples, reinforcing our hypothesis that the code is generated from a script. The mix of calls and desobfuscation involves a special effort to blend in a behavior that appears to be legitimate.

Control flow Variant1

Control flow Variant2

This technique allows to include binary data by limiting the entropy between 0.5 and 0.8, very close to what one would expect from legitimate instructions.

Entropy of the malware program

This data, once reconstituted, is placed in a buffer which is decrypted. We then recognize a Key Schedule type mechanism, followed by a generator reminiscent of RC4:

RC4 Random generator

The Rust paradigm makes these decryption routines difficult to identify. Tools such as findcrypt [8] do not detect them. The following check identifies the entire desobfuscation chain:

>>> magic.from_buffer(ARC4.new(b"DQOQHMLGYU").decrypt(binascii.unhexlify('31 18 14 11 30 A&
E5 D5 04 73 4D EA 62 29 20 F6 EF 2E B9 70 42 7D 0E 15 89 B6 71 7B 4D'.replace(' ', ''))))
'LZMA compressed data, non-streamed, siz 30 A8 02 4B D2 06 E4 69 7C CE A0 BE 4C FE 1D DB FE D2 34 21

Decryption in python

The decompression is performed by the version 0.1.3 of the lzma-rs library [9], and the decryption keys are easily visible in the code. The following script allows its extraction:

Key extraction script

The list of RC4 keys for the samples is as follows:

- **YDVHHCYTCH ;**
- **DQOQHMLGYU ;**
- **BWSCVQZXOB ;**
- **SIMHIDVSCR ;**
- **RGZPMAAQRP ;**
- **YVMOOVSIOF ;**
- **KSQKGUUTXZ ;**
- **EUWPUQYDTT ;**
- KHBXNNHKNN;

Unfortunately, the number of keys is insufficient to evaluate the quality of the PRNG [10]. Note that the charset is very small, which does not matter since the key is in plain text in the ".rdata" section. Given the extremely variable and complex call graph, writing a generic data extraction script that works for all samples is tedious. Now that we know the key (cf. previous extraction script), we need to find the encrypted data. These, once desobfuscated, have a header with invariant data:

 $\mathbf{1}$ 04066000 00 E0 0C 06 9C 00 89 00 00 00 00 00 00 00 00 00 .à. $. @...^\circ ...$ $090...$ $\overline{2}$ 04066010 00 40 0B 00 00 B0 0B 00 2E F5 39 D8 00 00 00 04 ¿ 1...0¨.KÒ.äi|Ï $\overline{\mathbf{3}}$ 04066020 31 18 14 11 30 A8 02 4B D2 06 E4 69 7C CF A0 BF $\overline{4}$ 5 041D3000 9C 00 A8 00 9C 00 A8 00 00 00 00 00 00 00 00 00 6 041D3010 00 B0 0A 00 00 F0 0A 00 0A 16 96 88 00 00 00 04 $^{\circ}$. . . \eth µ..~\.ýVëH0Ñr½Ü¥ 041D3020 B5 8D 1E 7E 5C 88 FD 56 EB 48 30 D1 72 BD DC A5 8 07301000 9C 00 C1 00 9C 00 C1 00 00 00 00 00 00 00 00 00 \ldots A \ldots A \ldots 10 07301010 00 80 0F 00 00 A0 0F 00 06 27 1B 75 00 00 00 04 $. _u$ RÆó.s.¹.ðECxûûj 52 C6 F3 9E 73 9F B4 1A F0 45 43 78 FB FB A1 9B 11 07301020

3 blocks of encrypted buffer header

The following script then allows to find the extracted data in memory, provided that the program has performed the decryption phase by carrying out an Egg-hunting search:

Encrypted data recovery script

The only thing missing to automate the extraction is to find an address or to set a breakpoint, which is done in a fairly logical way by finding the references to the RC4 key:

Malware launching script

And the script works on the first try for all samples. The source code of the tool is available on the TEHTRIS github [19].

An alternative method was presented at the Hack-it-N conference [16]. It consists in extracting the stage2 by emulation using the unicorn lib [17]. The source code is available on the TEHTRIS github [18].

The extraction of stage2 from the studied samples gives the following list:

The timestamps seem coherent between themselves and we note that the compilation dates correspond with a gap of 5 to 10 minutes, suggesting a manual packing step.

We note that stage2 verifies the presence of a virtual machine by testing the presence of the following executables:

- **> vboxservice.exe ;**
- **> vboxtray.exe ;**
- **> vmtoolsd.exe ;**
- **> vmwaretray.exe ;**
- **> vmwareuser.exe ;**
- **> vmacthlp.exe ;**
- **> vmsrvc.exe ;**
- **> vmusrvc.exe ;**
- **> prl_cc.exe ; > prl_tools.exe ; > xenservice.exe ; > qemu-ga.exe ;**
- **> windanr.exe.**

Once the desobfuscation step is done, the binary is no longer obfuscated and easily delivers its secrets. For example C2 [11]:

List of C2 URLs

Or encrypts list of files constituting the source code:

Source code metadata The analysis of the stage2 has already been done, however, and we will not describe it here.

Conclusion

Evading automated malware detection systems is a balancing act between binary entropy, hiding encryption functions, slowing down the analysis time, obtaining a consistent call/instruction ratio… The use of Rust adds a machine code overlay which is however much less than what a Go [12], delphi [13], cython [14], etc. compiler would do, making a manual analysis quite reasonable.

However, these techniques are no match for a reverser or a well-configured sandbox. The manual analysis of the most known and least traceable threats is a plus in the continuous improvement of our sandbox and EDR products. It is very likely that in the future, more and more malware programs will be developed in Rust.

* The use of stageN consists in separating the malware program into several specialized sub-parts in order to escape antivirus detection. In this case, stage1 is the malware program as sent to victims and stage2 is the payload encapsulated in stage1.

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