Aria-Body Loader? Is that you?

medium.com/insomniacs/aria-body-loader-is-that-you-53bdd630f8a1

BlueMonkey September 29, 2021

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9 min read

Hello! This is my first time writing a blog on a loader which I had gotten hold of. So, I am a new analyst in the Malware Analysis field and I am trying to do a research into cobalt strike. Recently, I ran a modified public YARA rule for cobalt strike on Virus Total and from the scan, I found two samples which I find interesting.

> 1e56c3f05bb53d2dfa60bc016e8509b12fd3beb5f567d274a184bb67af1eb19c
>
> c5696e660f3cfa9232756418e40ad18729cfe32fb284bba2314dd523ba527258

These two samples have a relative large size (17.05 MB) as compared to other files and their size is exactly the same. Additionally, their upload timing is quite close to each other, around 3 minutes apart. I started to analyze the files and from my finding, it doesn't look anything like a cobalt strike loader. Turning to my friend who have much more experience in malware analysis, I told him my findings and he told me that from what I had described, it sounds a little bit like Aria-Body instead. So I did some read up on Aria-Body and here are what I had found….

In 2020, Check Point Research release a write up (http://research[.]checkpoint[.]com/2020/nikon-apt-cyber-espionage-reloaded) describing how Naikon APT group is using Aria-body. In the report, they summarized the loader to have these capabilities:

1. Establish persistence via the Startup folder or the Run registry key [some variants].

2. Inject itself to another process such as rundll32.exe and dllhost.exe [some variants].

3. Decrypt two blobs: Import Table and the loader configuration.

4. Utilize a DGA algorithm if required

5. Contact the embedded / calculated C&C address in order to retrieve the next stage payload.

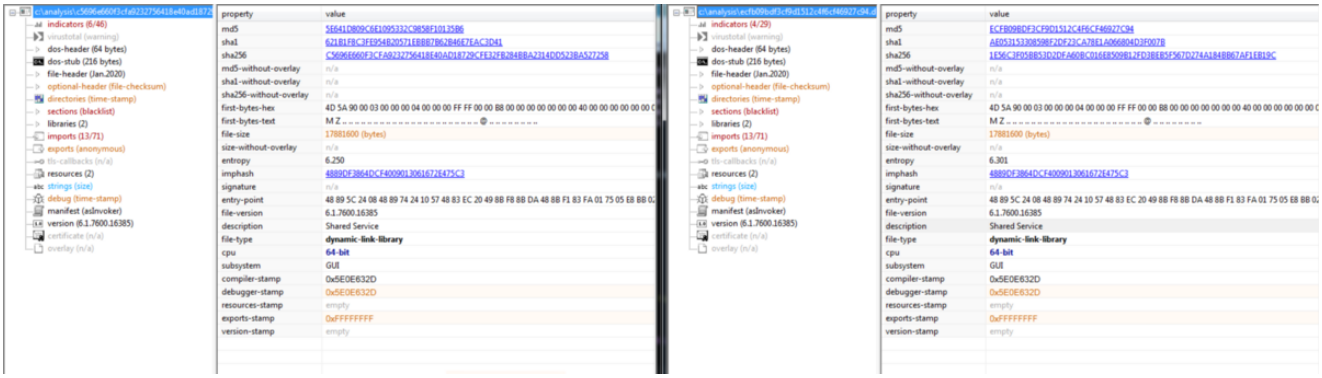
6. Decrypt the received payload DLL (Aria-Body backdoor).

7. Load and execute an exported function of the DLL — calculated using djb2 hashing algorithm.

Take note on these points as I will be mentioning some of them in this post.

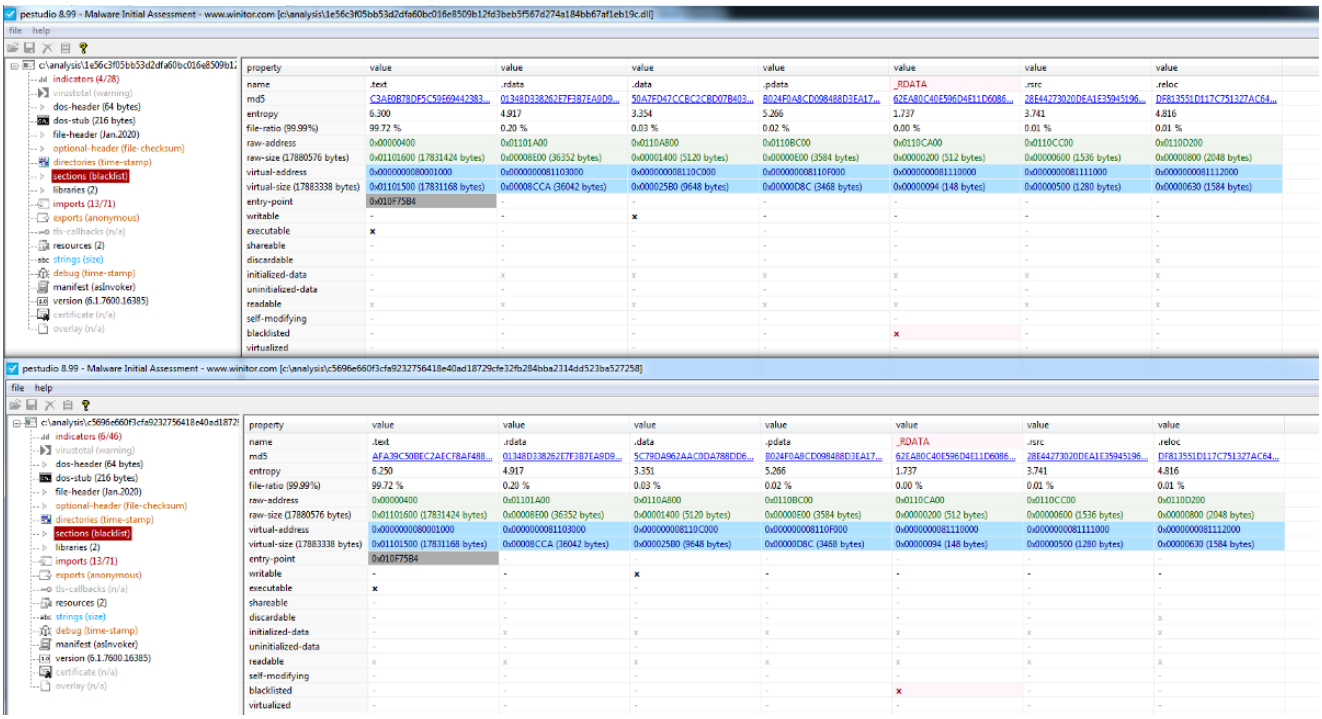
Analysis of the samples

As I had mentioned in the intro, the first thing that I noticed about the two sample is that both of them have exactly the same size. When I loaded them into PE studio, I noticed that they share the same compiler, debugger and exports timestamp.



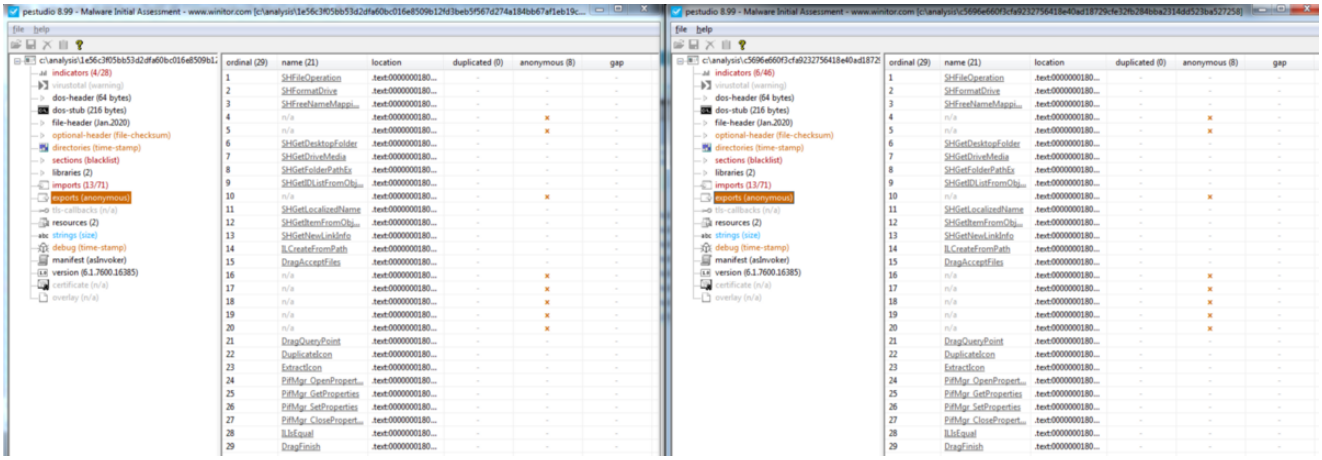
Similar Timestamp

Now looking at the section's metadata, we can see that all but two of them have the same hash. The two sections with the different hash is the *.text* and the *.data* section. Although they have a different hash, we can see that the raw-size, virtual-size, raw-address and virtual-address are the same.



Similar Section

Looking into the import and export section, this two samples also have a same import and export.



Similar Export

Looking at the static properties analysis, these two files seems like twins. They could have the same origins or they might be built using a builder. But this is just my speculation at this point as there are not enough information to support the claim.

Looking for the “action”

Now that I had done the analysis on the files properties and confirmed that it is a 64 bit DLL, it’s time to throw the sample 1e56c3f05bb53d2dfa60bc016e8509b12fd3beb5f567d274a184bb67af1eb19c into IDA for analysis. After IDA has finished loading, the first thing that it displayed is this:

```
; BOOL __stdcall DllMain(HINSTANCE hinstDLL, DWORD fdwReason, LPVOID lpvReserved)
DllMain proc near
mov     eax, 1
retn
DllMain endp
```

DllMain

That's all ._. not helpful at all. Seems like I have to find the "action" through other means~~

Since this is a DLL, we can try to look for the "action" by checking the export functions.

SHFileOperation	0000000180001210	1
SHFormatDrive	**0000000180001260**	**2**
SHFreeNameMappings	**0000000180001290**	**3**
	00000001800015F0	4
	00000001800015F0	5
SHGetDesktopFolder	**00000001800012C0**	**6**
SHGetDriveMedia	**00000001800012F0**	**7**
SHGetFolderPathEx	0000000180001320	8
SHGetIDListFromObject	**0000000180001350**	**9**
	00000001800015F0	10
SHGetLocalizedName	**0000000180001380**	**11**
SHGetItemFromObject	**00000001800013B0**	**12**
SHGetNewLinkInfo	00000001800013E0	13
ILCreateFromPath	0000000180001980	14
DragAcceptFiles	**0000000180001410**	**15**
	00000001800015F0	16
	00000001800015F0	17
	00000001800015F0	18
	00000001800015F0	19
	00000001800015F0	20
DragQueryPoint	**0000000180001470**	**21**
DuplicateIcon	**00000001800014A0**	**22**
ExtractIcon	00000001800014D0	23
PifMgr_OpenProperties	**0000000180001500**	**24**
PifMgr_GetProperties	**0000000180001530**	**25**
PifMgr_SetProperties	**0000000180001560**	**26**
PifMgr_CloseProperties	**0000000180001590**	**27**
ILIsEqual	**00000001800015C0**	**28**
DragFinish	**0000000180001440**	**29**
DllEntryPoint	**00000001810F75B4**	**[main entry]**

Export Functions

From this list of 30 export functions, two of them *DllEntryPoint* and *ILCreateFromPath* caught my attention. After looking through the two functions, I had determined that the *ILCreateFromPath* function contains the "actions" that we are interested in.

Obfuscation

While scrolling through the *ILCreateFromPath* function, I noticed a pattern:

```
push    rbx
sub     rsp, 30h
lea     rcx, aPrtPP0tnpnWsly ; "]prt-.p,^p,0tnpN.,wSlyowp,b"
mov     rbx, rdx
call    sub_180001010
mov     rcx, [rbx]
lea     rdx, kernel_20
call    rax
mov     cs:qword_18110E598, rax
test    rax, rax
jz      loc_180001AD2
```

Encoded String

Noticed that the value "*]prt-.p,^p,0tnpN.,wSlyowp,b*" in the variable *aPrtPP0tnpnWsly* was lea into *rcx* followed by calling the sub function *sub_180001010* then followed by a *call rax*. From my analysis, the function sub_18001010 consist of two parts.

1 — Decode the string

First it will decode a string that is passed in as argument which in this case is the value in the variable *aPrtPP0tnpnWsly*.

```
loc_180001100:
movsx   ecx, byte ptr [r9+r11]
lea     r11, [r11+1]
add     ecx, 1Ch
mov     eax, 1948B0FDh
imul    ecx
inc     edi
sar     edx, 3
mov     eax, edx
shr     eax, 1Fh
add     edx, eax
imul    eax, edx, 51h ; 'Q'
sub     ecx, eax
movsxd  rax, edi
add     cl, 2Ah ; '*'
mov     [r11-1], cl
cmp     rax, r8
jb      short loc_180001100
```

Decode Section

This function decodes the characters by applying the concept of Substitution cipher where it takes the ASCII value of each character, add 28 follow by mod 81 and finally add 42. This is the formula for the substitution cipher that I had just describe: *plain_text = (cipher_text + 28) % 81 + 42*. Thus, the value of variable *aPrtPP0tnpnWsly* decodes into *RegisterServiceCtrlHandlerW* which is a Win32 API.

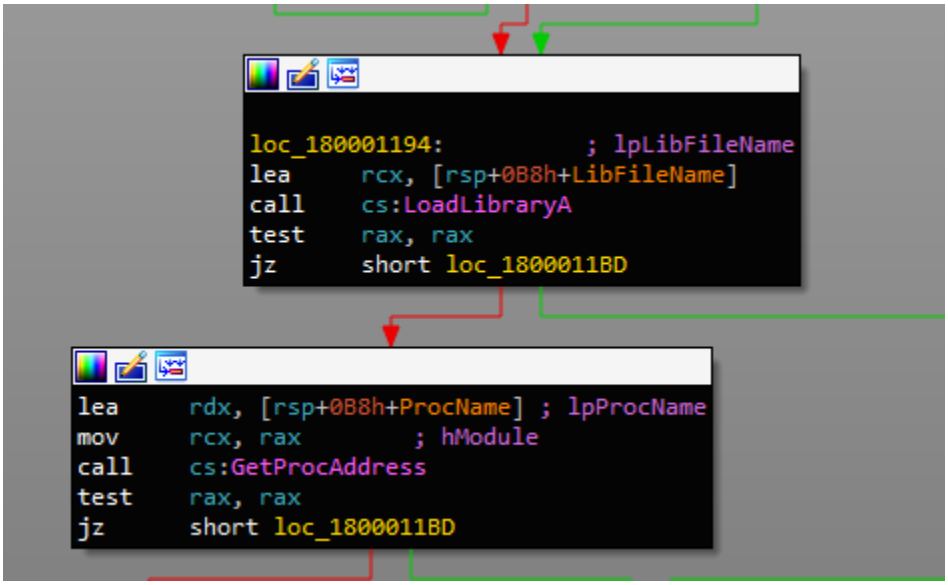
Using this formula, I wrote a simple python script to do the decryption and ran it on all the strings found in the same data section as *aPrtPP0tnpnWsly* and here are the decrypted strings:

Obfuscated	De-obfuscated
vp,ypw>=9oww	kernel32.dll
Rp.^3-.pxOt,pn.z,3b	GetSystemDirectoryW
N,pl.p[,znp--b	CreateProcessW
Rp._s,ploNzy.p2.	GetThreadContext
^p._s,ploNzy.p2.	SetThreadContext
at,./lwLwwznP2	VirtualAllocEx
b,t.p[,znp--Xpxz,3	WriteProcessMemory
]p-/xp_s,plo	ResumeThread
blt.Qz,^tyrwpZmupn.	WaitForSingleObject
lo0l*t>=9oww	advapi32.dll
]prt-.p,^p,0tnpN.,wSlyowp,b	RegisterServiceCtrlHandlerW
^p.^p,0tnp^.l./-	SetServiceStatus
zwp>=9oww	ole32.dll
NzN,pl.pR/to	CoCreateGuid
,*n,.?9oww	rpcrt4.dll
`/to_z^.,tyrb	UuidToStringW

Decoded String

Looking at the de-obfuscated strings, it seems like they are hiding function calls in strings and decode them during runtime so that we cannot most of its capabilities just from looking at imports table. From the list of the De-obfuscated strings, we can see that some of the capabilities of this malware includes creating thread and writing into memory.

2 — GetProcAddress



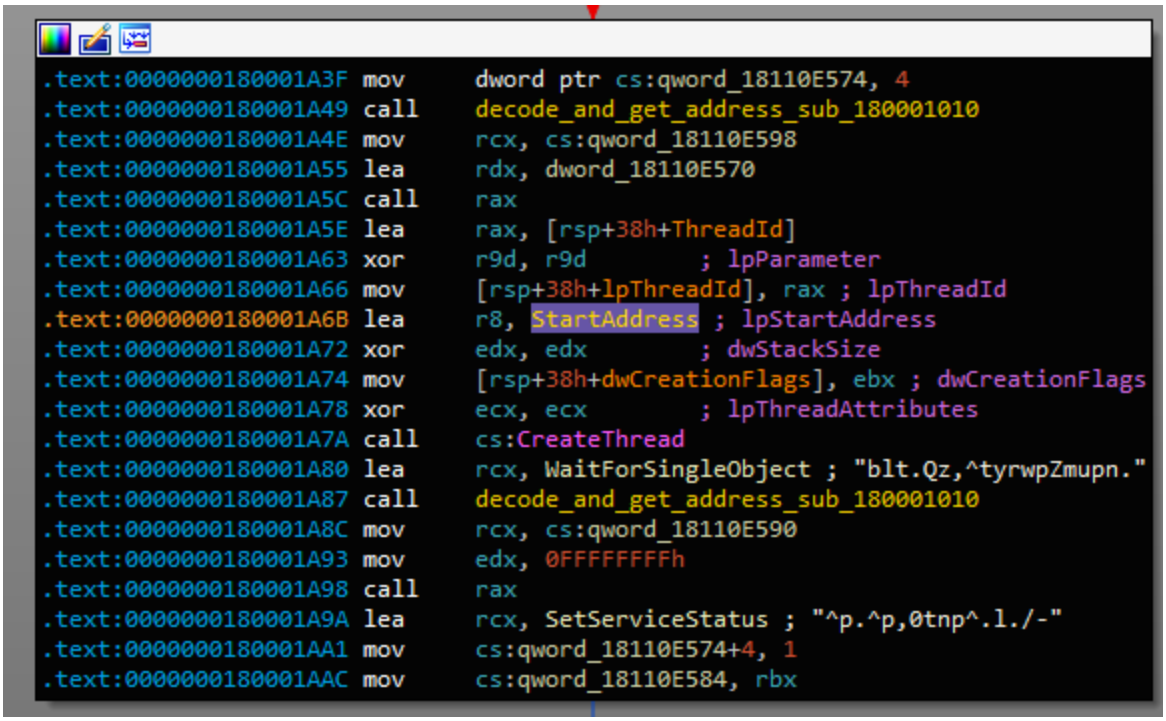
getProcAddress

Once the string is decoded, the function will then call *LoadLibraryA* and *GetProcAddress* before returning the address of the call.

Alright, now that we know that this malware has the ability to hide its function calls what's next?

Threading and New Processes

Following down the *ILCreateFromPath*, I saw that the malware creates a new thread to execute the function called *StartAddress*.



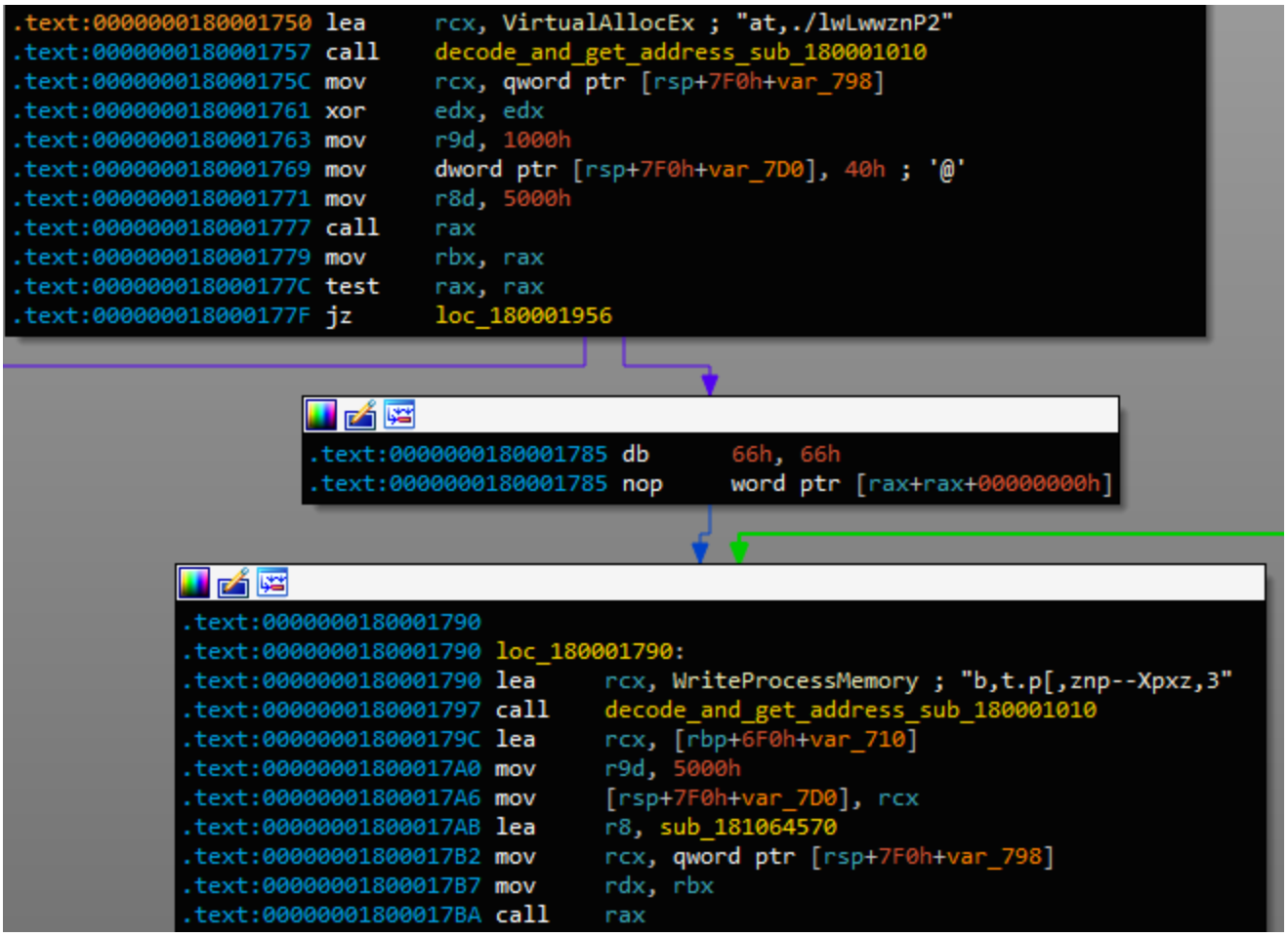
Create New Thread
So the only thing to do is to follow and look at what does the *StartAddress* function do.

Looking through the *StartAddress* function, the first thing that caught my eyes is a variable containing the string "*%s\\dllhost.exe /Processid:{%s}*" followed by *CreateProcessW*. It looks like the malware is trying to create a *dllhost* process.

```
.text:0000000180001698 lea     rdx, aSDllhostExePro ; "%s\\dllhost.exe /Processid:{%s}"
.text:000000018000169F lea     rcx, [rbp+6F0h+var_220] ; LPWSTR
.text:00000001800016A6 call    cs:wsprintfW
.text:00000001800016AC xorps   xmm0, xmm0
.text:00000001800016AF lea     rcx, [rbp+6F0h+var_6F0] ; void *
.text:00000001800016B3 xor     eax, eax
.text:00000001800016B5 xor     edx, edx        ; Val
.text:00000001800016B7 mov     r8d, 4D0h       ; Size
.text:00000001800016BD mov     [rbp+6F0h+var_720], rax
.text:00000001800016C1 movups  [rsp+7F0h+var_780], xmm0
.text:00000001800016C6 mov     [rsp+7F0h+var_788], rax
.text:00000001800016CB movups  [rbp+6F0h+var_770], xmm0
.text:00000001800016CF movups  [rbp+6F0h+var_760], xmm0
.text:00000001800016D3 movups  [rbp+6F0h+var_750], xmm0
.text:00000001800016D7 movups  [rbp+6F0h+var_740], xmm0
.text:00000001800016DB movups  [rbp+6F0h+var_730], xmm0
.text:00000001800016DF movups  [rsp+7F0h+var_798], xmm0
.text:00000001800016E4 call    memset
.text:00000001800016E9 lea     rcx, CreateProcessW ; "N,pl.p[,znp--b"
.text:00000001800016F0 call    decode_and_get_address_sub_180001010
.text:00000001800016F5 xor     edi, edi
.text:00000001800016F7 lea     rcx, [rsp+7F0h+var_798]
.text:00000001800016FC mov     [rsp+7F0h+var_7A8], rcx
.text:0000000180001701 lea     rdx, [rbp+6F0h+var_220]
.text:0000000180001708 lea     rcx, [rsp+7F0h+var_780]
.text:000000018000170D xor     r9d, r9d
.text:0000000180001710 mov     [rsp+7F0h+var_7B0], rcx
.text:0000000180001715 xor     r8d, r8d
.text:0000000180001718 mov     [rsp+7F0h+var_7B8], rdi
.text:000000018000171D xor     ecx, ecx
.text:000000018000171F mov     [rsp+7F0h+var_7C0], rdi
.text:0000000180001724 mov     [rsp+7F0h+var_7C8], 4
.text:000000018000172C mov     dword ptr [rsp+7F0h+var_7D0], edi
.text:0000000180001730 call    rax
```

Create New Process

After calling *CreateProcessW*, the malware then proceeds to call *VirtualAllocEx* followed by *WriteProcessMemory*.



Allocate And Write To Memory

From the above code, we can see that the malware used the *WriteProcessMemory* function to write the function *sub_181064570* into *dllhost* process created earlier. Although this seems to be a common process injection, it matches the checkpoint's report, where it mentions that it injects itself to another process such as rundll32.exe and dllhost.exe.

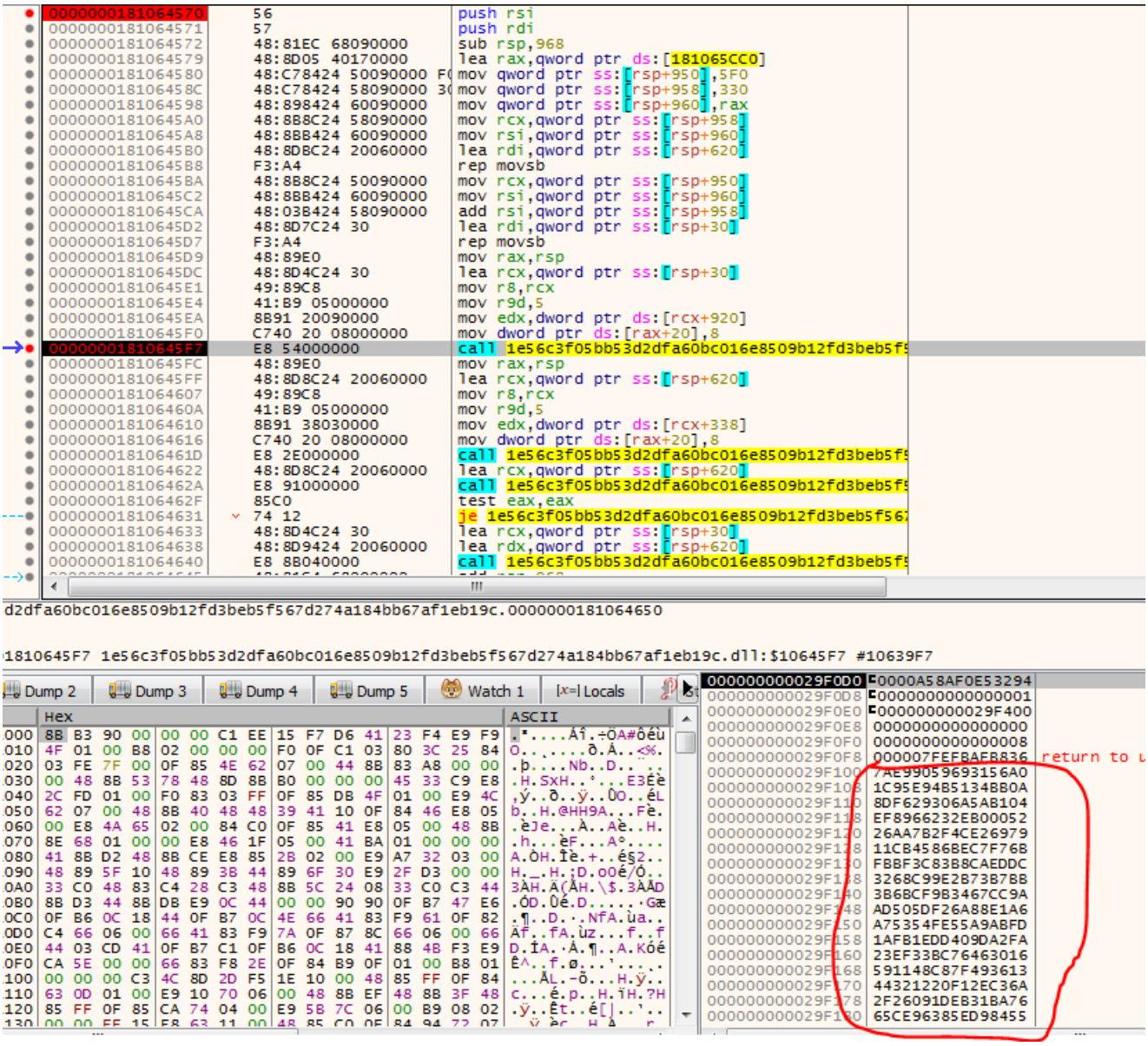
Decoding Embedded Data

Finally, we have reached the part where we can see what this malware actually wants to do! This is how the first few lines of the function written into the process's memory looks like:

```
.text:0000000181064570
.text:0000000181064570 push    rsi
.text:0000000181064571 push    rdi
.text:0000000181064572 sub     rsp, 968h
.text:0000000181064579 lea     rax, qword_181065CC0
.text:0000000181064580 mov     [rsp+978h+var_28], 5F0h
.text:000000018106458C mov     [rsp+978h+var_20], 330h
.text:0000000181064598 mov     [rsp+978h+var_18], rax
.text:00000001810645A0 mov     rcx, [rsp+978h+var_20]
.text:00000001810645A8 mov     rsi, [rsp+978h+var_18]
.text:00000001810645B0 lea     rdi, [rsp+978h+var_358]
.text:00000001810645B8 rep movsb
.text:00000001810645BA mov     rcx, [rsp+978h+var_28]
.text:00000001810645C2 mov     rsi, [rsp+978h+var_18]
.text:00000001810645CA add     rsi, [rsp+978h+var_20]
.text:00000001810645D2 lea     rdi, [rsp+978h+var_948]
.text:00000001810645D7 rep movsb
.text:00000001810645D9 mov     rax, rsp
.text:00000001810645DC lea     rcx, [rsp+978h+var_948]
```

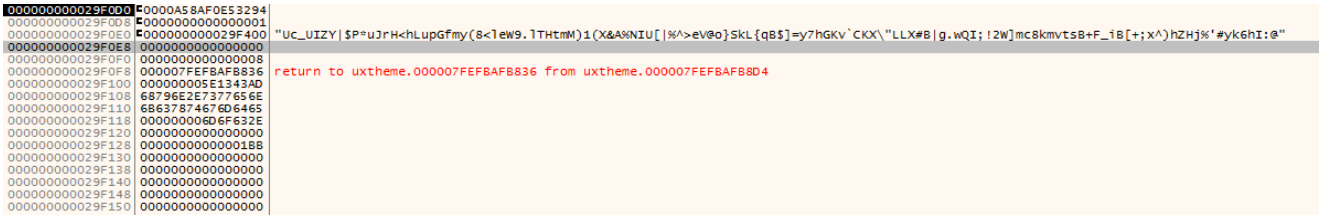
Write to Memory

It looks like the malware copying two sets of data located at *qword_181065CC0* into the memory.



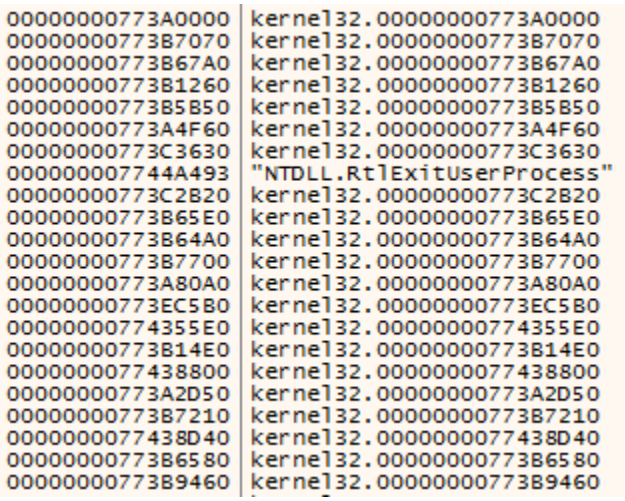
Encoded Blob in Memory

After copying the data into the memory, the malware calls a function which will decode the data.



First Decoded Blob

The first blob of data contains the URL of the C2 server "*news.nyhedmgtxck.com*" and a string of characters which doesn't seems to be used in any part of the execution.



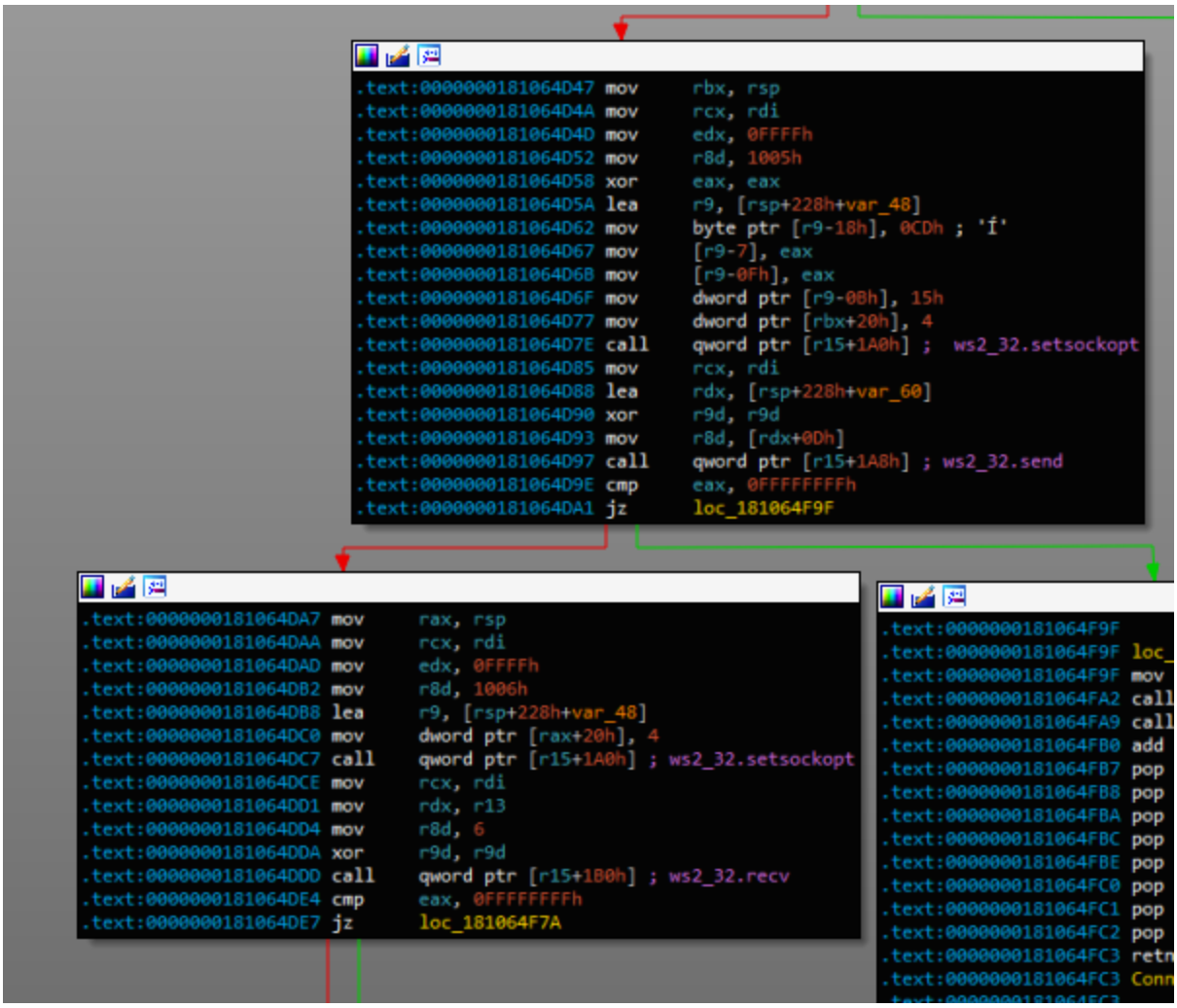
Second Decoded Blob

And the second blob of data contains the imports table which the malware will use in the next phase of its activity. Wait a minute… does the 2 blobs of data sounds familiar?

From checkpoint's report on aria-body loader, they mentioned that one of the functionality of the loader is to decrypt two blobs of data into an Import Table and a loader configuration.

Download and execute payload

By using the decoded import tables, the malware attempts to connects to the C2 URL to download a file.



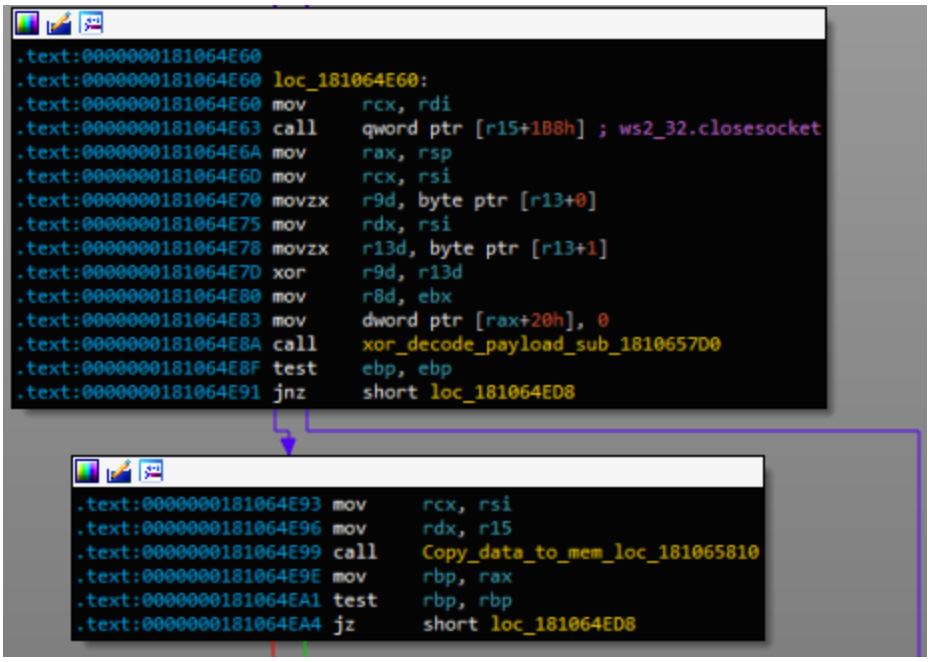
Download Payload

To this point, it actually matches the points mentioned in Check Point's report where Aria-body contact the embedded / calculated C&C address in order to download retrieve the next stage payload.

Too bad for us, the URL has already been sinkhole. Therefore, I am not be able download the sample for analysis):

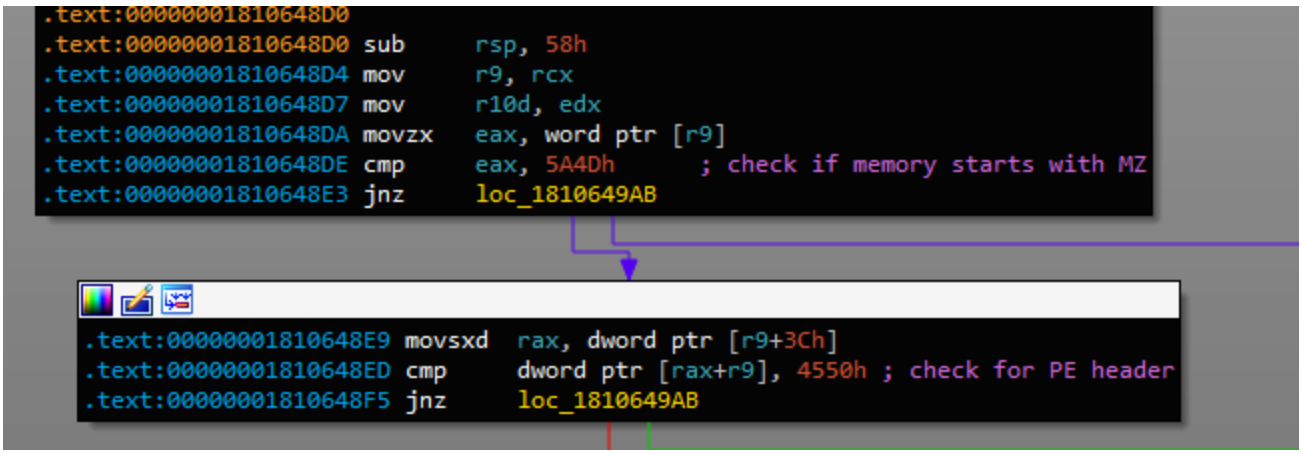
It's not the end yet! Although I am not able to analyze the next stage payload, I am still able to see what this loader does before passing control to the next stage payload :D

Once the payload is downloaded, the malware will first decode the payload with a XOR function. The decoded payload will then reside only in the memory. Which suggest that it could be a file-less malware.



Decode Payload and Copy to Memory

Next, the malware then calls a function which checks if the payload contains the magic number “MZ” and “PE”.



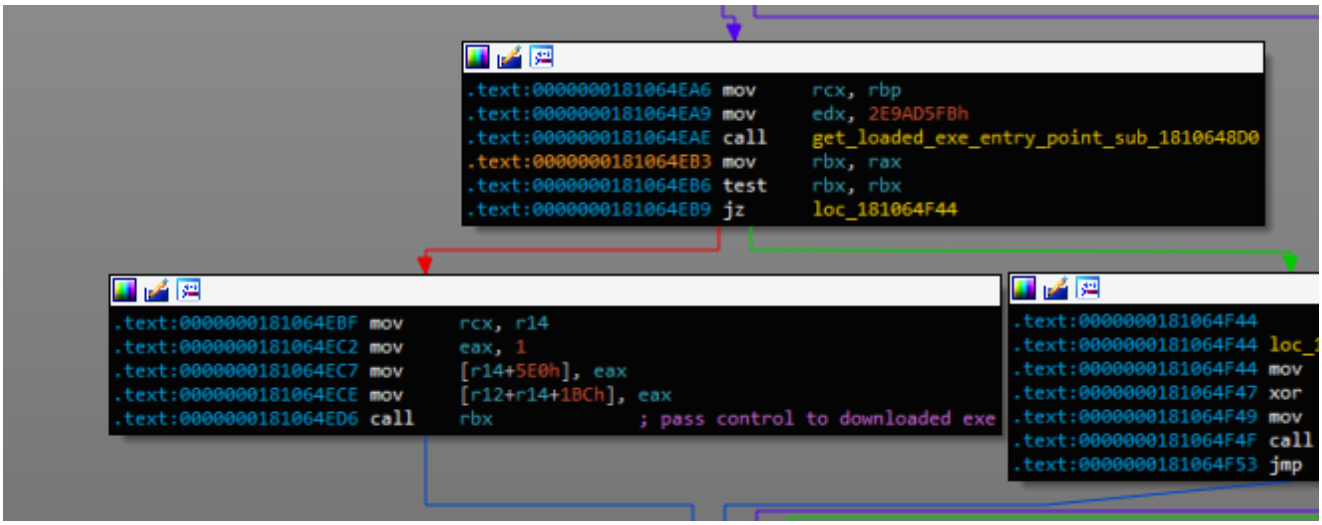
Check for PE and Section Header

Once verified, the malware will finally get the entry point to the payload by calculating the djb2 hash of the payload’s export and comparing it with 0x2E9AD5FB. Without the second stage payload, I am unable to determine what is the export name based on that hash.

```
result = 5381i64;
for ( i = (char)*a1; *a1; i = (char)*a1 )
{
  if ( a3 && i >= 0x61 )
    i -= 32;
  result = i + 33 * (_DWORD)result;
```

DJB2 Hash

Finally, the malware then passes the execution control to the payload.



The way this malware get the entry point also matches what Check Point had described where the loader loads and execute an exported function of the DLL — calculated using djb2 hashing algorithm.

Conclusion

Phew… Finally! We've reached the conclusion~~ v^^v

From the analysis, this malware looks like a loader which will download a payload from the C2 and execute the payload on the memory. The capabilities of this sample is very similar to the Aria-body loader that is described by Check Point where 5 out of the 7 points matching the analysis. I am unable to determine if this sample "establishes persistence via startup folder or run registry" and the "utilization of the DGA algorithm". Putting the capabilities aside, I had look through the sample with the hash "40c49ecbe1b7bdodbb935138661b6ca4" mentioned in Check Point's report and code wise, it looks vastly different from this sample.

Noticed that up to this point, I have only talked about the analysis of one of the samples. Well, I had done the analysis on both of the sample and in regards of the code executions, they are the same. The only difference between the two sample in regards to what is relevant to the execution and its function, is that the C2 string and the string of character in the first blob of data is different. Instead of going to "*news[.]nyhedmgtxck[.]com*", the C2 of the other sample is "*www[.]etnwtmrkh[.]com*" both of which are sinkholed.

Therefore, based on the capabilities, am I right to say that this could be a variant of Aria-Body loader?

Hashes Analyzed:

> 1e56c3f05bb53d2dfa60bc016e8509b12fd3beb5f567d274a184bb67af1eb19c
>
> c5696e660f3cfa9232756418e40ad18729cfe32fb284bba2314dd523ba527258