Inside STL: Smart pointers

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The C++ standard library comes with a few smart pointer types.

The simplest one is unique_ptr: This class babysits a raw pointer and remembers to delete it at destruction (in an appropriate manner). Dumping the contents of a unique_ptr is just looking at the raw pointer inside.

The complication is that there is also a deleter object in the unique_ptr. This deleter object is usually an empty class, so it is stored as part of a compressed pair.

The Visual Studio debugger has a visualizer that understands unique_ptr, but here's what it looks like at a low level in the Microsoft implementation:

```
0:000< ?? p
class std::unique_ptr<S,std::default_delete<S> >
  +0x000 _Mypair : std::_Compressed_pair<std::default_delete<S>,S *,1>
0:000< ?? p._Mypair
class std::_Compressed_pair<std::default_delete<S>,S *,1>
  +0x000 _Myval2 : 0x0000020a`11f08490 S
0:000< ?? p._Mypair._Myval2
struct S * 0x0000020a`11f08490 ← here is the unique object
  +0x000 a : 0n42
```
Next up is are the buddies shared_ptr and weak_ptr. These guys are just alternate policies around the same design, which centers around a "control block".

```
struct control_block
{
    virtual void Dispose() = 0;
    virtual void Delete() = 0;
    std::atomic<unsigned long> shareds;
    std::atomic<unsigned long> weaks;
};
template<typename T>
struct shared_ptr
{
    T* object;
    control_block* control;
};
template<typename T>
struct weak_ptr
{
    T* object;
    control_block* control;
};
```
The control block keeps track of how many shared and weak pointers exist, but in a somewhat unusual way. The naïve approach would be to have shareds and weaks count the number of shared and weak pointers, respectively. However, this ends up being more complicated because detecting when the last pointer to the control block has been destroyed would require us to perform atomic checks on *two* member variables simultaneously.

You can solve this by having weaks be the number of shared pointers *plus* the number of weak pointers. This is the total number of things keeping the control block alive. A test of a single atomic variable will tell you when the last pointer has been destroyed.

Here's how the fundamental operations play out with this design:

```
struct control_block
{
    virtual void Dispose() = 0;
    virtual void Delete() = 0;
    std::atomic<unsigned long> shareds;
    std::atomic<unsigned long> refs;
};
```
Create the first shared pointer: Create a new control block with shareds and refs both initialized to 1.

Copy a nonempty weak pointer: Atomically increment the refs. (We know that the refs is at least one, because it includes the weak pointer we are copying from.)

Copy a nonempty shared pointer: Atomically increment both shareds and refs. (We know that both are already at least one, because they includes the shared pointer we are copying from.)

Convert a nonempty shared pointer to a weak pointer: Atomically increment the refs. (We know that the refs is at least one, because it includes the weak pointer we are copying from.)

Convert a nonempty weak pointer to a shared pointer: Atomically increment the shareds, but only if it was originally nonzero.¹ If successful, then also increment the refs. If shareds was zero, then produce an empty shared pointer instead.

Destruct a nonempty weak pointer: Atomically decrement the refs. If this decrements to zero, then call Delete() to delete the control block.

Destruct a nonempty shared pointer: Atomically decrement the shareds. If this decrements to zero, then call Dispose() to destruct the managed object. Next, atomically decrement the refs. If this also decrements to zero, then call Delete() to delete the control block.

This design works, but it turns out you can be more clever about it: Instead of counting all shared pointers toward the refs, count only one shared pointer toward the refs, and update refs only when the number of shareds transitions between 1 and 0.

So here's our optimized version:

```
struct control_block
{
    virtual void Dispose() = 0;
    virtual void Delete() = 0;
    std::atomic<unsigned long> shareds;
    std::atomic<unsigned long> weaks;
};
```
Create the first shared pointer: Create a new control block with shareds and weaks both initialized to 1.

Copy a nonempty weak pointer: Atomically increment the weaks. (We know that the weaks is at least one, because it includes the weak pointer we are copying from.)

Copy a nonempty shared pointer: Atomically increment the shareds. (We know that the shareds is at least one, because it includes the shared pointer we are copying from.)

Convert a nonempty shared pointer to a weak pointer: Atomically increment the weaks. (We know that the weaks is at least one, because it includes the weak pointer we are copying from.)

Convert a nonempty weak pointer to a shared pointer: Atomically increment the shareds, but only if it was originally nonzero.¹ If it was zero, then produce an empty shared pointer instead.

Destruct a nonempty weak pointer: Atomically decrement the weaks. If this decrements to zero, then call Delete() to delete the control block.

Destruct a nonempty shared pointer: Atomically decrement the shareds. If this decrements to zero, then there's extra work to do: Destroy the associated object and decrement the weaks. If weaks also decrements to zero, then call Delete() to delete the control block.

There's a lesser-known feature of shared and weak pointers: The *aliasing contructor*. This lets you create a shared_ptr that dereferences to an object that is different from the one managed by the control block. For example, you can do this:

```
struct S { int a = 42; int b = 99; };
auto s = std::make\_shared<S>();auto i = std::shared\_ptr<int>(s, & s--b);
```
The shared pointer i points to the b member of the shared s object, but using the lifetime of the S object controlled by the pre-existing shared pointer s.

Internally, you create an aliasing shared pointer by using the same control block as the source shared pointer (incrementing the shareds, naturally), but setting the object pointer to the second constructor parameter.

Again, the Visual Studio debugger understands all of this and gives you a nice visualization. If you're debugging at a lower level, then here's what it looks like in the debugger:

```
0:000> ?? p2
class std::shared_ptr<S>
  +0x000 _Ptr : 0x0000020a`11f084e0 S
  +0x008 _Rep : 0x0000020a`11f084d0 std::_Ref_count_base
0:000> ?? p2._Ptr
struct S * 0x0000020a`11f084e0
  +0x000 a : 0n42
0:000> ?? p2._Rep
class std::_Ref_count_base * 0x0000020a`11f084d0
  +0x000 __VFN_table : 0x00007ff7`9e788758
  +0x008 _Uses : 1
  +0x00c _Weaks : 1
```
Here's how the names we used map to the names in the Microsoft implementation of the C++ standard library:

When inspecting a shared_ptr, you can just follow the _Ptr to get to the managed object. You know that the pointer is valid because this is a shared pointer which keeps the object alive. (If the _Ptr is nullptr, then the shared pointer is empty.)

A weak pointer looks pretty much identical in the debugger:

```
0:000> ?? p3
class std::weak_ptr<S>
  +0x000 _Ptr : 0x0000020a`11f084e0 S
  +0x008 _Rep : 0x0000020a`11f084d0 std::_Ref_count_base
0:000> ?? p3._Ptr
struct S * 0x0000020a`11f084e0
  +0x000 a : 0n42
0:000> ?? p3._Rep
class std::_Ref_count_base * 0x0000020a`11f084d0
  +0x000 __VFN_table : 0x00007ff7`9e788758
  +0x008 _Uses : 1
  +0x00c _Weaks : 2
```
When inspecting a weak pointer in the debugger, you have to check the control block's shared pointer count to know whether or not the object is still valid. If the shared pointer count is zero, then the weak pointer has expired, and the object pointer points to freed memory.

Next time, we'll look at std::make_shared and std::enable_shared_from_this.

Bonus viewing: [Stephan T. Lavavej's lecture on](https://learn.microsoft.com/en-us/shows/c9-lectures-stephan-t-lavavej-standard-template-library-stl-/3-of-n) shared_ptr and unique_ptr. He also gave an advanced lecture, which I cannot find.

¹ This can be accomplished by a compare_exchange loop. [Given the scenario,](https://devblogs.microsoft.com/oldnewthing/20180330-00/?p=98395) this would be better served by a compare_exchange_weak.