Anybody who writes #pragma pack(1) may as well just wear a sign on their forehead that says "I hate RISC"

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When you use **#pragma pack(1)**, this changes the default structure packing to byte packing, removing all padding bytes normally inserted to preserve alignment.

Consider these two structures:

```
// no #pragma pack in effect.
struct S
{
    int32_t total;
    int32_t a, b;
};
#pragma pack(1)
struct P
{
    int32_t total;
    int32_t a, b;
};
```

Both structures have identical layouts because the members are already at their natural alignment, Therefore, you would expect these two structures to be equivalent.

But they're not.

Changing the default structure packing has another consequence: It changes the alignment of the structure itself. In this case, the **#pragma pack(1)** declares that the structure **P** can itself be placed at any byte boundary, instead of requiring it to be placed on a 4-byte boundary.

```
struct ExtraS
{
    char c;
    S s;
    char d;
};
struct ExtraP
{
    char c;
    P p;
    char d;
};
```

Even though the structures S and P have the same layout, the difference in alignment means that the structures ExtraS and ExtraP end up quite different.

The ExtraS structure starts with a char, then adds three bytes of padding, followed by the S structure, then another char, and three more bytes of padding to bring the entire structure back up to 4-byte alignment. This ensures that an array of ExtraS structures will properly align all of the embedded S objects.

By comparison, the ExtraP structure starts out the same way, with a single char, but this time, there is no padding before the P because the P is byte-aligned. Similarly, there is no trail padding at the end of the structure because the P is byte-aligned and therefore does not need to be kept at a particular alignment in the case of an array of ExtraP objects.

	00	01	02	03	04	05	06	07	08	09	0A	0B	0C	0D	0E	OF	10	
ExtraS	С	pac	lding	s.total		s.a			s.b			d						
ExtraP	С	p.	tota	1		p.	a			p.	b			d				

The effect is more noticeable if you have an array of these objects.

	00	01	02	03	04	05	06	07	08	09	0A	OВ	0C	0 D	0E	0F	10
ExtraS	С	pac	padding s.total		s.a			s.b				d					
ExtraP	С	p.	tota	1	p.a		p.b				d	С	p.	total			

Observe that in the array of ExtraS objects, the s.total, s.a, and s.b are always four-byte aligned. But in the array of ExtraP objects, there is no consistent alignment for the members of p.

The possibility that any **P** structure could be misaligned has significant consequences for code generation, because all accesses to members must handle the case that the address is not properly aligned.

```
void UpdateS(S* s)
{
  s->total = s->a + s->b;
}
void UpdateP(P* p)
{
  p->total = p->a + p->b;
}
```

Despite the structures **S** and **P** having exactly the same layout, the code generation is different because of the alignment.

UpdateS	UpdateP
Intel Itanium	

adds r31 = r32, 4 adds r30 = r32 8 ;; ld4 r31 = [r31] ld4 r30 = [r30] ;;	adds r31 = r32, 4 adds r30 = r32 8 ;; ld1 r29 = [r31], 1 ld1 r28 = [r30], 1 ;; ld1 r27 = [r31], 1 ld1 r26 = [r30], 1 ;; dep r29 = r27, r29, 8, 8 dep r28 = r26 r28 8
add r31 = r30, r31 ::	8 1d1 r25 = [r31], 1 1d1 r24 = [r30], 1;; dep r29 = r25, r29, 16, 8 dep r28 = r24, r28, 16, 8 1d1 r27 = [r31] 1d1 r26 = [r30];; dep r29 = r27, r29, 24, 8
st4 [r32] = r31	<pre>dep r28 = r26, r28, 24, 8 ;; add r31 = r28, r29 ;; st1 [r32] = r31 adds r30 = r32, 1 adds r29 = r32, 2 extr r28 = r31, 8, 8 extr r27 = r31, 16, 8 ;; st1 [r30] = r28</pre>
br.ret.sptk.many rp Alpha AXP	<pre>st1 [r29] = r27, 1 extr r26 = r31, 24, 8 ;; st1 [r29] = r26 br.ret.sptk.many.rp</pre>

ldl	t1, 4(a0)	ldq_u t1, 4(a0) ldq_u t3, 7(a0) extll t1, a0, t1				
ldl	t2, 8(a0)	extlh t3, a0, t3 bis t1, t3, t1 ldq_u t2, 8(a0) ldq_u t3, 11(a0) extll t2, a0, t2 extlh t3, a0, t3				
addl stl	t1, t1, t2 t1, (a0)	bis t2, t3, t2 addl t1, t1, t2 ldq_u t2, 3(a0) ldq_u t5, (a0) inslh t1, a0, t4 insll t1, a0, t3 msklh t2, a0, t2 mskll t5, a0, t5 bis t2, t4, t2 bis t5, t3, t5 stq_u t2, 3(a0) stg u t5, (a0)				
ret 1	zero, (ra),	ret zero, (ra), 1				
MIPS	R4000					
lw	t0, 4(a0)	lwl t0, 7(a0) lwr t0, 4(a0)				
lw	t1, 8(a0)	lwl t1, 11(a0)				
addu	t0, t0, t1	addu t0, t0, t1 swl t0, 3(a0)				
jr sw	ra t0, (a0)	jr ra swr t0, (a0)				
PowerPC 600						

lwz	r4, 4(r3)	lbz r4, 4(r3)
		lbz r9, 5(r3)
		rlwimi r4, r9, 8, 16, 23
		lbz r9, 6(r3)
		rlwimi r4, r9, 16, 8, 15
		lbz r9, 7(r3)
		rlwimi r4, r9, 24, 0, 7
lwz	r5, 8(r3)	lbz r5, 8(r3)
		lbz r9, 9(r3)
		rlwimi r5, r9, 8, 16, 23
		lbz r9, 10(r3)
		rlwimi r5, r9, 16, 8, 15
		lbz r9, 11(r3)
		rlwimi r5, r9, 24, 0, 7
addu	r4, r4, r5	addu r4, r4, r5
stw	r4, (r3)	stb r4, (r3)
		rlwimi r9, r4, 24, 0, 31
		stb r9, 1(r3)
		rlwimi r9, r4, 16, 0, 31
		stb r9, 2(r3)
		rlwimi r9, r4, 8, 0, 31
		stb r9, 3(r3)
blr		blr
0		1
Super	ʻH-3	

mov.l	@(4	, r4),	r2	mov.b shll8	@(7, r4), r1 r1
				mov.b	@(6, r4), r2
				extu.b	r2, r2
				or	r2, r1
				sh118	r1
					(0, 14), 12
				or	r2 r1
				shll8	r1
				mov.b	@r4, r2
				extu.b	r2, r2
_				or	r1, r2
mov.l	@(8	, r4),	r3	mov.b	@(7, r4), r1
				Sh118	r_1
				or	r3, r1
				shll8	r1
				mov.b	@(5, r4), r3
				extu.b	r3, r3
				or	r3, r1
				shll8	r1
				mov.b	@r4, r3
				or	r1 r3
add	r3.	r2		add	r3, r2
	,			mov.b	r2, @r4
				shlr8	r2
				mov.b	r2, @(1, r4)
				shlr8	r2
				mov.b	r2, @(2, r4)
rte				snirð	ΓZ
mov 1	r2	@r⊿		mov h	$r_{2} = 0(3, r_{4})$
movit	,	914			
1					

Observe that for some RISC processors, the code size explosion is quite significant. This may in turn affect inlining decisions.

Moral of the story: Don't apply **#pragma pack(1)** to structures unless absolutely necessary. It bloats your code and inhibits optimizations.

Bonus chatter: Once you make this mistake, you can't go back. You allowed the structure to be byte-aligned, and if you remove the spurious **#pragma pack(1)**, you are making the structure more strictly aligned, which will be a breaking change for any clients which used the byte-packed version.

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