

The Intel 80386, part 15: Common compiler-generated code sequences

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The Microsoft compiler employs a few patterns in its code generation that you may want to become familiar with.

As we saw earlier, a call to a `__thiscall` function passes the `this` pointer in the `ecx` register, with the remaining parameters passed on the stack. A typical calling sequence would go something like this:

```
; p->Foo(x, 42);

push    42          ; parameter 2
push    dword ptr [ebp-24h]   ; parameter 1
mov     ecx, dword ptr [ebp-20h] ; "this" for call
call    CThing::Foo        ; call the function directly
```

If the method is virtual, then there is a vtable lookup:

```
; p->Foo(x, 42);

push    42          ; parameter 2
push    dword ptr [ebp-24h]   ; parameter 1
mov     ecx, dword ptr [ebp-20h] ; "this" for call
mov     eax, dword ptr [ecx]    ; fetch vtable
call    dword ptr [eax+10h]    ; call through the vtable
```

If the method uses `__stdcall` instead of `__thiscall` (typically because it is a COM method), then the `this` parameter is passed on the stack rather than in the `ecx` register.

```

; p->Foo(x, 42);

; non-virtual call
push    42                      ; parameter 2
push    dword ptr [ebp-24h]       ; parameter 1
push    dword ptr [ebp-20h]       ; "this" for call
call    CThing::Foo              ; call the function directly

; virtual call
push    42                      ; parameter 2
push    dword ptr [ebp-24h]       ; parameter 1
mov     ecx, dword ptr [ebp-20h]  ; "this" for call
push    ecx                      ; pass as stack parameter
mov     eax, dword ptr [ecx]     ; fetch vtable
call    dword ptr [eax+10h]       ; call through the vtable

```

The Microsoft compiler uses a jump table for dense `switch` statements, but it adds a level of indirection so that multiple cases that leads to the same target share the same jump entry.

Consider the following fragment:

```

switch (value)
{
case 2:
case 3:
case 5:
case 7:
    printf("prime");
    break;

case 4:
case 9:
    printf("perfect square");
    break;

default:
    printf("I'm sure you're special");
    break;
}

```

The resulting code may look like this:

```

mov    eax, dword ptr [ebp-30h]      ; load value
sub    eax, 2                         ; table starts with "case 2"
cmp    eax, 8                         ; table has 8 entries
jae    case_default                  ; not in table, go to default case
movzx  eax, byte ptr [eax+level1]    ; get the index into the second table
jmp    dword ptr [eax+level2]        ; jump to handler

...
level2 dd  offset case_prime         ; slot 0 is for cases 2, 3, 5, and 7
      dd  offset case_square        ; slot 1 is for cases 4 and 9
      dd  offset case_default       ; slot 2 is for everybody else
level1 db  0, 0, 1, 0, 2, 0, 2, 1   ; generate the slots
;           2 3 4 5 6 7 8 9          ; corresponding cases

```

Adding a level of indirection allows the level-2 jump table to be smaller. The trade-off is that you have to pay for a level-1 jump table, but if there are a lot of duplicates (such as those created by all the missing cases that go to `default:`), the trade-off may be worth it.

If you stare at the output and do some reverse-compiling, you can imagine that the compiler internally rewrote it as

```

enum SwitchResult { Prime, PerfectSquare, Default };
static const unsigned char level1[] =
{ Prime, Prime, PerfectSquare, Prime,
  Default, Prime, Default, PerfectSquare };
unsigned int index1 = (unsigned)value - 2;
switch (index1 < 8 ? level1[index1] : Default)
{
case Prime:
    printf("prime");
    break;

case PerfectSquare:
    printf("perfect square");
    break;

case Default:
    printf("I'm sure you're special");
    break;

default:
    __assume(0); // not reached
}

```

Next time, we'll wrap up our quick tour of the 80386 by walking through a simple function.

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